

# Inverse Iteration for Calculating the Spectral Radius of a Non-Negative Irreducible Matrix

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## ABSTRACT

Noda established the superlinear convergence of an inverse iteration procedure for calculating the spectral radius and the associated positive eigenvector of a non-negative irreducible matrix. Here a new proof is given, based completely on the underlying order structure. The main tool is Hopf's inequality. It is shown that the convergence is quadratic.

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## 1. INTRODUCTION

Throughout this paper  $A$  will denote a non-negative irreducible  $N \times N$  matrix with spectral radius  $\rho$  and associated positive eigenvector  $p$ .

In [5], Noda established the convergence of an inverse iteration procedure for the determination of  $\rho$  and  $p$ . He also showed that the convergence is superlinear. Here we shall prove that it is at least quadratic.

This is an easy by-product of our proof of convergence, which uses only the underlying order structure and not (as in [5]) the Jordan form. The main tool is Hopf's inequality. As it has been used for bounding the eigenvalues  $\neq \rho$ , it is quite natural to use it for convergence proofs, too.

## 2. DEFINITIONS; TWO LEMMAS

An  $N \times N$  matrix  $B = (b_{ik})$  is called positive (non-negative) if  $b_{ik} > 0$  ( $\geq 0$ ),  $i, k = 1, \dots, N$ . We write  $B > 0$  ( $\geq 0$ ). For vectors,  $y > 0$ ,  $y \geq 0$  are defined in an analogous way.

For a pair of vectors  $x, y$  with  $y > 0$ , we define

$$\max\left(\frac{x}{y}\right) = \max_i \frac{x_i}{y_i}, \quad \min\left(\frac{x}{y}\right) = \min_i \frac{x_i}{y_i},$$

$$\text{osc}\left(\frac{x}{y}\right) = \max\left(\frac{x}{y}\right) - \min\left(\frac{x}{y}\right).$$

Hopf's inequality [1, 3, 6] states: For  $B > 0$  and any pair of vectors  $x, y$ , where  $y > 0$ ,

$$\text{osc}\left(\frac{Bx}{By}\right) \leq N(B) \text{osc}\left(\frac{x}{y}\right). \quad (1)$$

Here

$$N(B) = \frac{\sqrt{K(B)} - 1}{\sqrt{K(B)} + 1}$$

and

$$K(B) = \sup_{\substack{u > 0 \\ v > 0}} \left\{ \max\left(\frac{Bu}{Bv}\right) \max\left(\frac{Bv}{Bu}\right) \right\}. \quad (2)$$

It is obvious that

$$N(tB) = N(B), \quad t > 0, \quad (3)$$

$$N(pq^T) = 0, \quad p > 0, \quad q > 0, \quad (4)$$

$$N(D_1 B D_2) = N(B), \quad (5)$$

where  $D_i$  ( $i=1, 2$ ) are diagonal matrices with positive diagonal entries. A bound for  $N(B)$  is [3, 6]

$$N(B) \leq \frac{m_1 - m_2}{m_1 + m_2}, \quad m_1 = \max_{i,k} b_{ik}, \quad m_2 = \min_{i,k} b_{ik}. \quad (6)$$

LEMMA 1. Let  $p > 0$ ,  $q > 0$ ,  $\tilde{p} = \min p_i$ ,  $\tilde{q} = \min q_i$ , and  $B = (b_{ik})$  be a positive matrix such that

$$|b_{ik} - p_i q_k| \leq \varepsilon. \quad (7)$$

Then

$$N(B) \leq \frac{\varepsilon}{\tilde{p}\tilde{q}}. \quad (8)$$

*Proof.* Define  $\tilde{B} = (\tilde{b}_{ik})$ ,  $\tilde{b}_{ik} = b_{ik}/p_i q_k$ ,  $\tilde{\varepsilon} = \varepsilon/\tilde{p}\tilde{q}$ . Then by (7)

$$\max \tilde{b}_{ik} \leq 1 + \tilde{\varepsilon}, \quad \min \tilde{b}_{ik} \geq 1 - \tilde{\varepsilon},$$

and hence by (5) and (6)

$$N(B) = N(\tilde{B}) \leq \tilde{\varepsilon}. \quad \blacksquare$$

REMARK. By taking  $p_i = (m_1 + m_2)/2$ ,  $q_i = 1$ ,  $\varepsilon = (m_1 - m_2)/2$  in Lemma 1, (8) yields the bound (6).

LEMMA 2. For a given number  $\lambda_0 > \rho$  there is an  $M > 0$  such that

$$N((\lambda I - A)^{-1}) \leq M(\lambda - \rho), \quad \rho < \lambda \leq \lambda_0. \quad (9)$$

*Proof.* The adjoint  $\text{adj}(B)$  of a square matrix  $B$  satisfies the relation [4, p. 13]

$$B \text{adj}(B) = \text{adj}(B) B = (\det B) I$$

In particular, for  $\lambda > \rho$ ,

$$(\lambda I - A)^{-1} = \frac{1}{\det(\lambda I - A)} \text{adj}(\lambda I - A).$$

Hence by (3),

$$N((\lambda I - A)^{-1}) = N(\text{adj}(\lambda I - A)). \quad (10)$$

On the other hand,

$$\text{adj}(\rho I - A) = pq^T,$$

where  $q > 0$ ,  $A^T q = \rho q$ ,  $q$  suitably normalized. Equation (9) follows now from Lemma 1. ■

### 3. THE ITERATIVE PROCEDURE

We define

$$\|x\| = \max\left(\frac{x}{p}\right). \quad (11)$$

Let  $\{B_n\}$ ,  $n = 0, 1, \dots$  be a sequence of positive matrices commuting with  $A$ . Assume the existence of  $\gamma$  such that

$$N(B_n) \leq \gamma < 1, \quad n = 0, 1, \dots \quad (12)$$

For given  $x_0 > 0$ , define iteratively

$$\tilde{x}_{n+1} = B_n x_n, \quad (13)$$

$$x_{n+1} = \frac{\tilde{x}_{n+1}}{\|\tilde{x}_{n+1}\|}, \quad (14)$$

$$\bar{\lambda}_{n+1} = \max\left(\frac{Ax_{n+1}}{x_{n+1}}\right), \quad \underline{\lambda}_{n+1} = \min\left(\frac{Ax_{n+1}}{x_{n+1}}\right). \quad (15)$$

$\bar{\lambda}_0, \lambda_0$  are defined analogously.

We prove first some useful relations:

LEMMA 3. For  $n = 1, 2, \dots$ ,

$$\bar{\lambda}_n - \rho \leq \rho \frac{\text{osc}(x_n/p)}{1 - \text{osc}(x_n/p)} \leq C\rho(\bar{\lambda}_n - \rho), \quad (16)$$

$$\rho - \underline{\lambda}_n \leq \rho \text{osc}\left(\frac{x_n}{p}\right) \leq \tilde{C}(\rho - \underline{\lambda}_n), \quad (17)$$

where  $C, \tilde{C}$  depend on  $A$ , and  $C$  also on  $\bar{\lambda}_0$ .

*Proof.* From  $\|x_n\| = 1$ ,  $n > 0$  we get

$$1 - \text{osc}\left(\frac{x_n}{p}\right) \leq \frac{x_{n,i}}{p_i} \leq 1.$$

Hence for suitable  $s$

$$\bar{\lambda}_n - \rho = \sum_k a_{sk} \frac{p_k}{p_s} \left( \frac{x_{n,k}}{p_k} \frac{p_s}{x_{n,s}} - 1 \right) \leq \rho \left( \frac{1}{1 - \text{osc}(x_n/p)} - 1 \right),$$

showing the left inequality of (16). The left inequality of (17) follows in an analogous way. For the other inequalities we use a result in [2, Folgerung 2, p. 72]: Let  $x > 0$ ,  $z > 0$ , and  $Ax \leq \alpha x$ ,  $Az \geq \beta z$ , and choose  $i$  so that  $x_i/z_i$  is minimal. For any  $k \neq i$  there is an  $s \leq n-1$  such that  $a_{ik}^{(s)} = (A^s)_{ik} > 0$  and

$$\frac{x_i}{z_i} \leq \frac{x_k}{z_k} \leq \left( 1 + \frac{\alpha^s - \beta^s}{a_{ik}^{(s)}} \frac{z_i}{z_k} \right) \frac{x_i}{z_i} \quad (18)$$

and

$$\left( 1 - \frac{\alpha^s - \beta^s}{a_{ik}^{(s)}} \frac{x_i}{x_k} \right) \frac{z_i}{x_i} \leq \frac{z_k}{x_k} \leq \frac{z_i}{x_i}. \quad (19)$$

Taking  $x = x_n$ ,  $\alpha = \bar{\lambda}_n$ ,  $z = p$ ,  $\beta = \rho$  in (18), we get

$$\min\left(\frac{x_n}{p}\right) \leq \frac{x_{n,k}}{p_k} \leq [1 + C(\bar{\lambda}_n - \rho)] \min\left(\frac{x_n}{p}\right)$$

for suitable  $C$  depending on an upper bound for  $\bar{\lambda}_n$ ,  $n = 1, 2, \dots$ . According to Theorem 1 such a bound is provided by  $\bar{\lambda}_0$ . Thus

$$\text{osc}\left(\frac{x_n}{p}\right) \leq C(\bar{\lambda}_n - \rho) \left[ 1 - \text{osc}\left(\frac{x_n}{p}\right) \right],$$

yielding the right inequality in (16).

Taking  $x = p$ ,  $\alpha = \rho$ ,  $z = x_n$ ,  $\beta = \underline{\lambda}_n$  in (19), we get for a suitable  $\tilde{C}$

$$[1 - \tilde{C}(\rho - \underline{\lambda}_n)] \max\left(\frac{x_n}{p}\right) \leq \frac{x_{n,k}}{p_k} \leq \max\left(\frac{x_n}{p}\right)$$

or

$$\text{osc}\left(\frac{x_n}{p}\right) \leq \tilde{C}(\rho - \underline{\lambda}_n).$$

This is the second inequality in (17). ■

**THEOREM 1.** *Consider the procedure (13)–(15). For  $n = 0, 1, 2, \dots$ ,*

$$\underline{\lambda}_n \leq \underline{\lambda}_{n+1} \leq \rho \leq \bar{\lambda}_{n+1} \leq \bar{\lambda}_n, \quad (20)$$

$$\lim \underline{\lambda}_n = \lim \bar{\lambda}_n = \rho, \quad (21)$$

$$\lim x_n = p. \quad (22)$$

If  $x_n \neq p$  for all  $n$ , then the inequalities of (20) are strict.

*Proof.* If we multiply the relation

$$\underline{\lambda}_n x_n \leq Ax_n \leq \bar{\lambda}_n x_n$$

by  $B_n$  and use  $B_n A = AB_n$ , we get

$$\underline{\lambda}_n \tilde{x}_{n+1} \leq A\tilde{x}_{n+1} \leq \bar{\lambda}_n \tilde{x}_{n+1}$$

and hence  $\underline{\lambda}_n \leq \underline{\lambda}_{n+1}$ ,  $\bar{\lambda}_{n+1} \leq \bar{\lambda}_n$ . If  $x_n \neq p$ , then  $\bar{\lambda}_n x_n - Ax_n \neq 0$ ; hence  $\bar{\lambda}_n \tilde{x}_{n+1} - A\tilde{x}_{n+1} \geq 0$  and  $\underline{\lambda}_{n+1} < \bar{\lambda}_n$ . Similarly  $\underline{\lambda}_n < \underline{\lambda}_{n+1}$ . The remaining inequalities  $\underline{\lambda}_n \leq \rho \leq \bar{\lambda}_n$  follow from the quotient theorem (e.g., [4], II, 5.5.2). From (16), (17) we infer the strict inequalities for  $x_n \neq p$ . Now

$$\begin{aligned} \bar{\lambda}_{n+1} - \underline{\lambda}_{n+1} &= \text{osc}\left(\frac{Ax_{n+1}}{x_{n+1}}\right) = \text{osc}\left(\frac{B_n Ax_n}{B_n x_n}\right) \leq N(B_n) \text{osc}\left(\frac{Ax_n}{x_n}\right) \\ &= N(B_n)(\bar{\lambda}_n - \underline{\lambda}_n). \end{aligned} \quad (23)$$

From (12) and (20), we infer (21). From (17), we get

$$\lim_{n \rightarrow \infty} \operatorname{osc} \left( \frac{x_n}{p} \right) = 0$$

or

$$\lim_{n \rightarrow \infty} \frac{x_{n,i}}{p_i} = 1, \quad i = 1, \dots, N.$$

Hence, we get (22). ■

In the case of  $A > 0$ , taking  $B_n = A$ , Theorem 1 gives the convergence of the usual power method. If only  $A^m > 0$  for a suitable integer  $m$ , i.e., if  $A$  is primitive, the proof given above can be easily adapted to yield the same result. In fact,

$$\bar{\lambda}_{n+m} - \underline{\lambda}_{n+m} = \operatorname{osc} \left( \frac{A^m A x_n}{A^m x_n} \right) \leq N(A^m) (\bar{\lambda}_n - \underline{\lambda}_n).$$

More interesting is the case

$$B_n = (\bar{\lambda}_n I - A)^{-1}, \quad n = 0, 1, \dots \quad (24)$$

If we start with an  $x_0$  such that  $Ax_0 \neq \rho x_0$ , then  $x_n \neq p$ ,  $\bar{\lambda}_n > \rho$ ,  $B_n > 0$  for all  $n$ , as can be proved by induction.

Hence, Theorem 1 can be applied and gives the convergence of the inverse iteration procedure considered by Noda [5]. Additionally, we have the following statement about the rate of convergence:

**THEOREM 2.** *In the iteration procedure (13)–(15) with*

$$B_n = (\bar{\lambda}_n I - A)^{-1},$$

*the sequences  $\{\underline{\lambda}_n\}$ ,  $\{\bar{\lambda}_n\}$  converge quadratically to  $\rho$  and the  $\{x_n\}$  quadratically to the eigenvector  $p$ .*

*Proof.* From (23) and (9) we get

$$\bar{\lambda}_{n+1} - \underline{\lambda}_{n+1} \leq M(\bar{\lambda}_n - \rho)(\bar{\lambda}_n - \underline{\lambda}_n) \leq M(\bar{\lambda}_n - \underline{\lambda}_n)^2,$$

i.e.,  $\{\bar{\lambda}_n - \underline{\lambda}_n\}$  converges quadratically to zero. It is now obvious from (16) and (17) that the sequences

$$\{\bar{\lambda}_n - \rho\}, \quad \{\rho - \underline{\lambda}_n\}, \quad \text{osc}\left(\frac{x_n}{p}\right)$$

also converge quadratically. ■

*Note added in proof.*

The author learned that Theorem 2 has also been proved in Stephen M. Robinson-Karl Nickel: Computation of the Perron root and vector of a nonnegative matrix, MRC Technical Summary Report #1100, September 1970, Mathematics Research Center, University of Wisconsin-Madison, Madison, Wisconsin 53706.

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