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Integrative Simulation Framework for Modeling Dynamic Cellular Phenomena in 3D over Time

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"Writing a book is an adventure. To begin with, it is a toy and an amusement; then it becomes a mistress, and then it becomes a master, and then a tyrant. The last phase is that just as you are about to be reconciled to your servitude, you kill the monster, and fling him out to the public."

Sir Winston [Churchill](#page-145-0) [\(1949\)](#page-145-0)

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> "We shall not grow wiser before we learn that much that we have done was very foolish."

> > Friedrich August [von Hayek](#page-158-0) [\(1944\)](#page-158-0)

Contents

Introduction

1.1 Motivation

The emerging field of systems biology allows for the application and combination of experimental, theoretical, and modeling techniques. A key goal of systems biology is to understand biological processes as whole systems instead of isolated parts. With such a systems-level analysis, the study of biological phenomena at the molecular, cellular, or behavioral levels becomes more feasible [\(Kitano,](#page-151-0) [2002b;](#page-151-0) [Kell,](#page-150-0) [2004\)](#page-150-0). Improvements in the measurement of molecular interactions and rates have revolutionized our insight into the cell and its dynamics. An enormous amount of information has been gained over the past decades adding to the already-known mechanisms. This leads to a vast set of parameters making it difficult to evaluate new hypotheses on intuition alone. Thus, applying modeling techniques and computer simulations help to confirm such hypotheses [\(Kitano,](#page-151-1) [2002a;](#page-151-1) [Shapiro et al.,](#page-156-0) [2002\)](#page-156-0).

In recent years much progress in the development of cell-biological modeling and simulation tools has been achieved. A multitude of existing programs can already fulfill the major requirements of execution speed and result accuracy. However, often such programs are highly specialized in their methodology and the applications they were designed for. There is no system that allows for the use of more than one algorithm in a simulation at a time. Furthermore, the existing tools are still behind on the development of three-dimensional ([3D](#page-134-2)) visualization techniques. In comparison to existing two-dimensional ([2D](#page-134-3)) simulation applications, the [3D](#page-134-2) approaches are still at their early stages. The additional spatial information allows for a more direct referencing of the model to its real object. Methods applied to [3D](#page-134-2) share many characteristics with [2D](#page-134-3) techniques. Many existing [2D](#page-134-3) methods therefore can be either adopted or carried forward.

Along with mechanistic models it is often possible for system behavior to be formalized into solvable differential equations. On the basis of such equations the evolution of time is regarded as predictable as well as continuous and is therefore a very fast and precise method [\(Ghosh and Tomlin,](#page-147-0) [2004\)](#page-147-0). Nevertheless,

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deterministic approaches to model and simulate the dynamics of intracellular regulatory processes have their limitations. For the reason of combinatoric increase in the number of equations and the number of contained players this process becomes highly impractical. If only small numbers of particles are involved, deterministic modeling approaches and simulation might fail. Stochastic approaches, using closer knowledge of the subcellular architecture, are more appropriate in these situations [\(Kiehl et al.,](#page-150-1) [2004\)](#page-150-1).

A significant obstacle with given modeling and simulation tools is that they seem to be able to handle only a small set of applications. They were developed for simulating special models and lack the ability to integrate extra techniques for solving other problems. In fact, often systems focus their attention on handling a special problem to all its details in minimal computational time [\(Pettinen et al.,](#page-155-0) [2005\)](#page-155-0). The currently used programs therefore vary noticeably in their applicability for specific types of modeling. Moreover the integration of such programs is highly demanding. They are closed systems and often do not offer component interfaces. In practice it could be beneficial to allow for such interfaces to combine various different tools to just one tool only. A compromise to the absence of direct interfaces is the possibility of exchanging information via standardized file formats. For this, many tools use the Systems Biology Markup Language ([SBML](#page-136-0)) as a standard [\(Hucka et al.,](#page-149-0) [2004\)](#page-149-0). Often, by the use of such standardized formats, important information to the model cannot be handled correctly and might get lost. Especially with [SBML](#page-136-0), the geometrical information cannot be saved entirely and subsequently has to be reconstructed.

Applications for such a modeling and simulation tool that allows for both [3D](#page-134-2) visualization and concurrent algorithms are imperative. The applications lie in the areas of electrical excitable cells, circadian rhythm, cell cycle, cellular motility, membrane transporters, metabolic pathways, and signal transduction networks. In comparison to [2D](#page-134-3) methods, the use of a [3D](#page-134-2) geometry provides considerably more significant data. The possibility of simulating different compartments of a cell with diverse algorithms can reduce computational effort dramatically. Consequently, this allows for both improved modeling and more information output on the spatio-temporal behavior of a system. Thus it is a scientific challenge to integrate concurrently running algorithms in combination with [3D](#page-134-2) visualization.

1.2

This work addresses the problem of an in silico biology system-level analysis on the base of cellular signaling networks. Grounded on biological observations, computational simulation models for the calcium(II) (Ca^{2+}) (Ca^{2+}) (Ca^{2+}) signaling pathway were developed in an attempt to understand the non-linear dynamics of the system.

To create this kind of an all-encompassing application, several existing simulation systems were evaluated for their strengths and weaknesses. Current state-of-the-art cell simulation applications are either limited to a [2D](#page-134-3) representation of a cell or do not use any cellular geometry at all. The only exceptions to this finding are the deterministic simulator "VirtualCell", the stochastic simulators "[MCell](#page-135-0)" and "SmartCell", as well as this approach of a hybrid fourdimensional ([4D](#page-134-5)) Cell Simulator ([4DiCeS](#page-134-0)). Despite the great variety of software packages available for modeling, simulation, and analysis of data [\(Hucka et al.,](#page-149-0) [2004\)](#page-149-0) as described above, there was no application that features the complete and variable integration of different simulation methods in a [3D](#page-134-2) environment.

The key improvement, which [4DiCeS](#page-134-0) has over the other existing systems, is the ability not to be fixed on either deterministic or stochastic modeling and simulation approaches. A system of specialized interfaces therefore was designed to allow the bonding of interchangeable algorithmic modules of various types. The internal representation of the model is designed for the easy exchangeability of data. Furthermore, the integration of cell model file format standards is permitted by exchangeable plug-ins. The implemented system includes an Application Programming Interface ([API](#page-134-6)) for writing individual plug-ins to utilize different simulation algorithms. This facilitates the implementation of tailored programs and specific algorithms that can be developed for data mining as well as visualization. The resulting [4DiCeS](#page-134-0) framework presented in this work describes a concept for the integration of heterogenous data into an easy-to-use software.

1.3 Structure

This thesis is primarily concerned with the simulation of cellular phenomena. An application is provided to utilize hybrid mathematical models and to simulate [4D](#page-134-5) spatial dynamics within a cell. The work reveals that a [3D](#page-134-2) cellular geometry holds extended information of importance, and is closest to reality in its model and simulation. Also, an application framework will be shown that allows for the integration of various particle reaction and diffusion algorithms. Such algorithms can then be applied to a model either in sole or even in concurrency.

The work is divided into three parts. The first part (Chapters [1,](#page-6-0) [2,](#page-10-0) and [3\)](#page-24-0) provides this motivation and an introduction to the topic of cellular dynamics, and its modeling and simulation. Relevant current modeling and simulation application will be introduced. The second part consists of Chapter [4](#page-36-0) and focuses on the formal description, design, and the methodology of the [4D](#page-134-5) cell simulator. The third part (Chapters [5](#page-88-0) and [6\)](#page-108-0) presents a comparison of related works with the results of the project, discusses this work, and gives a brief outlook on further improvements to the system.

A condensed summary of all chapters of this work is provided with the following overview:

Chapter [2](#page-10-0) reviews biological concepts of systems biology as well as mathematical aspects of modeling, simulation, and analysis of cellular dynamics.

Chapter [3](#page-24-0) describes a number of existing systems that represent state-of-theart biochemical modeling and simulation applications.

Chapter [4](#page-36-0) focuses on the formal description as a basis for the design of the [4DiCeS](#page-134-0) framework for systems biology.

Chapter [5](#page-88-0) comprises the results that were obtained during development and testing [4DiCeS](#page-134-0).

Chapter [6](#page-108-0) evaluates and discusses the [4DiCeS](#page-134-0) platform including some ideas that will illustrate further development and directions.

Dynamics in Systems Biology

During the past century much progress was achieved in the measurement of cellular processes, molecular interactions, as well as their kinetics. Thus, a revolution was initiated in the understanding of the dynamics within biological cells. The following Figure [2.1](#page-10-1) of a pyramid composed of different molecules wants to give insight into the complexity of cellular organization.

Figure 2.1: Life's complexity pyramid. The genomic information is both stored and translated into functional units such as proteins and metabolites. These units form operational molecules, consisting of regulatory systems or metabolic pathways. On top of these rather small units, large scale organizations implement the characteristic features of an organism. Adapted from Oltvai and Barabási [\(2002\)](#page-154-0).

The fundamental units are arranged to either metabolic and signaling pathways or to motifs in genetic-regulatory networks. Motifs and pathways are linked to operational groups that are responsible for discrete cellular functions [\(Hartwell](#page-149-1) [et al.,](#page-149-1) [1999\)](#page-149-1). Such groups are then nested hierarchically and characterize the large-scale organization of a cell [\(Ravasz et al.,](#page-155-1) [2002\)](#page-155-1).

While creating models of cellular dynamics, one has to understand complex properties. Such properties include the combination of regulatory mechanisms and interlocking transport in and among cells. Electrical activity, signal transduction, or other biochemical networks are examples of intricate behaviors taking place on the cellular level. In general such dynamic phenomena refer to arbitrary processes occurring and therefore changing over time. This forms the highly dynamic basis for living cells. To maintain the typical characteristics of a cell's life such as growth, movement, responsiveness, cell division, and intercellular communication, cells must continuously obtain energy from their direct neighborhood. Therefore cells need to act far from static thermal equilibrium as thermodynamically open systems. Hence, cells require a huge amount of energy to sustain the gradients of metabolites and ions in order to function properly [\(Fall and Keizer,](#page-146-0) [2002\)](#page-146-0).

The following two Sections [2.1](#page-11-0) and [2.2](#page-17-0) give an overview on the scope of modeling and simulation of such dynamic systems within cells. By doing so, the need for simulation tools, such as the [4DiCeS](#page-134-0) software presented here, handling such models becomes apparent.

2.1 __________________ Modeling and Simulation

Theoretical methods joined with experimental measuring have offered comprehensive perception of dynamics for many years [\(Ortoleva et al.,](#page-155-2) [2003\)](#page-155-2). Computers have demonstrated to be a necessity in assisting the dissection of molecular processes. Yet, the bare amount of quantitative experimental cellular information allows for the cooperation of computer science and biology [\(Chong and](#page-145-1) [Ray,](#page-145-1) [2002\)](#page-145-1). The interaction of theory, experiment, and computation succeeds a conceptual formulation analog to successfully proven physical models (see Table [2.1\)](#page-12-0).

Here all modeling is an abstraction of reality. The only exact model of any system is the system itself. So when a model of a system is designed, a choice must be made regarding the level of detail and feature types to be included into that model. To a large extent, this is prescribed by the characteristics of the examined system, the type of experimental data available, and the type of questions that are addressed to modeling [\(Bolouri and Davidson,](#page-144-0) [2002\)](#page-144-0).

Table 2.1: Conceptual formulation of models. Here, the interdependency of experiment, theory, and computation describe the production of a conceptual formulation of models. All given steps depend highly on a close collaboration with experimentalists working at the same problem. Adapted from [Fall and Keizer](#page-146-0) [\(2002\)](#page-146-0).

The problems theorists encounter in biology are therefore very alike to that in physical sciences. At this level equations are analyzed, if possible simplified, solved, and, most importantly predictions can be made. These predictions are checked by further experiments. Moreover, such experiments may disclose discrepancies that in turn will require changes to the model [\(Alvarez-Vasquez](#page-142-1) [et al.,](#page-142-1) [2005\)](#page-142-1). The procedure addressed here is an improving cycle of approximations where the theoretical model acts as a quantitative hypothesis (see Figure [2.2\)](#page-13-1).

The history of simulation and in silico analysis of biological systems dates back to the earliest mechanical and analogue computers in 1940 [\(Chance,](#page-145-2) [2004\)](#page-145-2). The recent progress in quantitative simulation and modeling are due to advances in modern information technology and was enhanced by the recent burst of molecular data. It becomes obvious that future progress in the understanding of biological functions will rely on the development and the use of computational methods [\(Arkin,](#page-142-2) [2001\)](#page-142-2). Therefore, the following two Sections [2.1.1](#page-13-0) and [2.1.2](#page-15-0) give a closer look at both, the roles of biology as well as computation to the modeling of biological cellular systems.

Figure 2.2: Systems biology triad. At the center are the complex biological phenomena. Interpretation of observations and data is supported by computational algorithms. The translation of interpretation to understanding is supported by systems science. In turn, systems science provides a framework for understanding. It indicates hypotheses to be tested and modified in an iterative cycle of experimentation. Adapted from [Mesarovic et al.](#page-152-0) [\(2004\)](#page-152-0).

2.1.1 Biology as a Model Featurer

Cells show very complex and different behaviors. Single cells are able to contract, excrete, move, reproduce, send signals, or even respond to them. Furthermore, cells accomplish the energy handling required for such activities. In cooperation cells manage all of the various processes required to perpetuate life as is [\(Hofmeyr,](#page-149-2) [1986\)](#page-149-2). However, everything that cells do can be represented in the form of basic natural laws. Although the rules of behavior are rather basic, cells comprise huge and complex networks of interacting substrates. An enormous amount of work was used disentangling only very few of such reaction schemes, and it is quite obvious that there are many more such interaction networks yet to be revealed [\(Keener and Sneyd,](#page-150-2) [2001\)](#page-150-2). Table [2.2](#page-14-0) gives a brief overview of dynamic behaviors happening on a cellular level. Accordingly, the need for powerful modeling and simulation tools, as [4DiCeS](#page-134-0), arises very quickly from the study of such phenomena.

In the subsequent sections the dynamical behavior of Ca^{2+} Ca^{2+} signaling for its spatio-temporal patterns (see Figure [2.3\)](#page-14-1) will be discussed in more detail. The modeling and simulation of such a phenomenon were used as examples for testing the [4DiCeS](#page-134-0) software. The role of computational techniques handling such models is going to be presented in the following subsection.

Table 2.2: Phenomena of cellular dynamics are here depicted, which have trigger events in common. There exist theoretical models and simulations for each incident. Extracted from [Fall et al.](#page-146-1) [\(2002\)](#page-146-1).

Figure 2.3: A spiral wave of Ca^{2+} Ca^{2+} ions detected from a dye with microinjection of [IP](#page-135-1)³ into an Xenopus laevis oocyte after 30 and 60 seconds. By courtesy of James D. Lechleiter, University of Texas, USA.

2.1.2 Computation as a Model Solver

Techniques applied to research problems in systems biology comprise very much of applied mathematics and computer science. A majority of features in computational modeling of cell biology play an important role. One of these features is the development of algorithms and techniques that give tools for numerical analysis [\(Takahashi et al.,](#page-157-0) [2002\)](#page-157-0). The computation of mathematical problems on computers is basically an estimation process. The efficiency and accuracy of these methods of estimation are the subjects of intensive study. The work on constructing models is also basically an approximation process. This is due to necessary simplifications that must be used to produce a helpful model. These simplifications must both be valid in terms of the physical process being studied and from a mathematical point of view [\(Eungdamrong and Iyengar,](#page-146-2) [2004\)](#page-146-2).

Computer models allow for testing conditions that may be difficult to obtain or that have not been examined in the laboratory yet. Therefore, every solution of a mathematical model may offer a simulation of a potential or real experiment. Such simulations can help to estimate parameters, i.e. diffusion or kinetic constants, that are challenging to collect in experiments. Simulations can verify how pharmacological agents may affect a biological process [\(Mendes and Kell,](#page-152-1) [1998\)](#page-152-1). Hypothesis about the role of individual mechanistic components can be checked by simulations. Accordingly, predictions made by such simulations then can be tested in the laboratory. It has to be noted that often the most important result of a simulation is a negative one. Therefore, a well-crafted model has to be redefined and tested again [\(Fall and Keizer,](#page-146-0) [2002\)](#page-146-0).

Improvements in numerical analysis as well as computer hardware have made the solving of complex deterministic systems accurate and fast. However, models of biological processes almost always comprise nonlinear components in their control mechanisms. By using traditional mathematical methods, as coupled Ordinary Differential Equations ([ODE](#page-136-1)s), often such problems can be solved exactly. But nonlinearities often create difficulties in getting any exact solution. Admittedly, good estimates of nonlinearities can be obtained by using computer implemented numerical methods. Often a major property in cellular mechanisms is spatial variation. Therefore, the analyzing and the solving of spatially explicit Partial Differential Equations ([PDE](#page-136-2)s) is very important. Such [PDE](#page-136-2)s can be more complex and less analytically tractable than [ODE](#page-136-1)s.

Some models need to handle noise and the tracking of single particles instead of particle concentrations. Here stochastic methods such as Monte Carlo ([MC](#page-135-2)) simulations [\(Metropolis and Ulam,](#page-152-2) [1949\)](#page-152-2) come into play. Such discrete models facilitate the qualitative modeling and are based on various different computa-tional approaches (Hofestädt and Meineke, [1995\)](#page-149-3). Unlike the different deterministic methods, a stochastic approach does not approximate the model as a continuous macroscopic system. In contrast, it treats a model as a discrete and microscopic process. However, this increase in accuracy comes at high costs as each individual chemical entity has to be modeled as a stochastic process. Hence, stochastic simulations are computationally more demanding (Puchalka [and Kierzek,](#page-155-3) [2004\)](#page-155-3).

Table 2.3: Important modeling and simulation techniques. The given methods represent the main approaches by which modeling and simulation are handled in systems biology. Applications vary in their support for either a pure method's implementation or hybrid attempts mixing different techniques. Extracted from [Hucka et al.](#page-149-0) [\(2004\)](#page-149-0).

Various attempts have been made to construct and simulate biochemical behaviors. The majority of approaches depend on the use of the before mentioned deterministic and stochastic techniques [\(Hucka et al.,](#page-149-0) [2004\)](#page-149-0). Recently other well-established techniques have been applied including boolean networks [\(Kauffman,](#page-150-3) [1969\)](#page-150-3), Cellular Automata ([CA](#page-134-7)) [\(von Neumann,](#page-158-1) [1966\)](#page-158-1), Bayesian networks [\(Pearl,](#page-155-4) [1988\)](#page-155-4), and petri nets [\(Petri,](#page-155-5) [1962\)](#page-155-5), to biological applications. Hybrid methods that combine the best features of all approaches exist as well [\(Lu et al.,](#page-152-3) [2004\)](#page-152-3). Table [2.3](#page-16-0) summarizes some highly recognized basic techniques for modeling and simulation in systems biology.

To solve equations that result from a model is only one part of the work. The other side needs to comprehend the model's behavior. Mathematical methods were developed for the system analysis of models that characterize complex processes. Such methods disclose the dynamical behavior, properties, and structure of the system. This is very much as molecular biological, physiological, and anatomical techniques uncover the physical basis of the model [\(Hartwell et al.,](#page-149-1) [1999\)](#page-149-1). Hence, the analysis exposing complicated behaviors within a model may result in further study of these biological phenomena. Admittedly, decisive analysis of complex equations demands skill. This is due to the fact that there are many intricacies that can only be comprehended through intensive training [\(Mishra et al.,](#page-152-4) [2003\)](#page-152-4). The construction and alteration of simple models is within the reach of cell biologists. Therefore it is necessary for scientists to seek association with mathematicians and computer scientists for the effectual simplification of complex models [\(Sontag,](#page-156-1) [2004\)](#page-156-1).

The next section will go into the detail of a well-studied Ca^{2+} Ca^{2+} model systems. This model was chosen to test the [4DiCeS](#page-134-0) software in its functionality and accuracy. The next chapter on the other hand will give insight into the design of the [4DiCeS](#page-134-0) system itself. After reading these sections, it will become obvious that the application should be capable of applying any of the before mentioned mathematical techniques. As for now stochastic and deterministic methods are implemented to be used by [4DiCeS](#page-134-0) [\(Oleson et al.,](#page-154-1) [2006\)](#page-154-1).

2.2 Cellular Calcium Models

Cellular Ca^{2+} Ca^{2+} has an overall very low concentration. At rest it is approximately 0.1μ M, and only about $1-10\mu$ M at its peak. On the other hand, potassium (K^+) (K^+) (K^+) and sodium (Na^+) (Na^+) (Na^+) show millimolar concentrations. Cells require to keep cytoplasmic [Ca](#page-134-4)²⁺ concentration ($[Ca^{2+}]_i$ $[Ca^{2+}]_i$) at low levels due to the fact that Ca^{2+} can modify the enzymatic properties of binding proteins. Hence increases in the cellular Ca^{2+} Ca^{2+} level are locally defined and quick to circumvent the runaway activation of enzymatic cascades. Two basic mechanisms hold responsibility for this impoundment and buffering. The buffers are highly specialized Ca^{2+} Ca^{2+} binding proteins that absorb 95–99% of the cytosol's Ca^{2+} Ca^{2+} . Ca^{2+} is impounded to either the sarcoplasmic reticulum ([SR](#page-136-5)) in muscle cells or the endoplasmic reticulum ([ER](#page-135-5)) in all other cell types. Proteins hydrolyze adenosine–3',5'– triphosphate ([ATP](#page-134-9)) to transport Ca^{2+} Ca^{2+} against rampant concentration gradients. Such proteins are [ATP](#page-134-9) hydrolases ([ATPase](#page-134-10)s) that are classified as $SR/ER Ca²⁺$ $SR/ER Ca²⁺$ $SR/ER Ca²⁺$ $SR/ER Ca²⁺$ $SR/ER Ca²⁺$ $SR/ER Ca²⁺$ transport [ATPase](#page-134-10) ([SERCA](#page-136-6)) pumps. On the other side, plasma membrane ([PM](#page-136-7)) $Ca^{2+}-ATPase$ $Ca^{2+}-ATPase$ $Ca^{2+}-ATPase$ ([PMCA](#page-136-8)) pumps dispose Ca^{2+} of the cell. [SR](#page-136-5) and [ER](#page-135-5) membranes have ion channels, which are different from [PMCA](#page-136-8)s, that transports Ca^{2+} Ca^{2+} back into the cytoplasm. Therefore, every cell has Ca^{2+} Ca^{2+} pumps for homeostasis as well as negative feedback. Some cells have developed ion channels such as the inositol–1,4,5–trisphosphate (IP_3) (IP_3) (IP_3) receptor (IP_3R) , which is activated and inhibited by Ca^{2+} Ca^{2+} . The [IP](#page-135-6)₃R is able to give both positive and negative feedback. Thus, brief channel openings may enable oscillations of free cytoplasmic Ca^{2+} Ca^{2+} that are utilized for signaling. Interestingly, Ca^{2+} Ca^{2+} oscillations were discovered in vitro after they were predicted by a model [\(Chay and Keizer,](#page-145-3) [1983\)](#page-145-3).

2.2.1 Calcium Signaling

 $Ca²⁺$ $Ca²⁺$ is the most common cellular signals' carrier. Due to its special adaptability as a ligand, it regulates very many important aspects of cellular activity. This goes from the creation of new life at fertilization to the radical incident of cellular apoptotic suicide. Signaling by Ca^{2+} Ca^{2+} shows a number of properties that make it unique among all other carriers of signals. An important example is its ability to function both as a first and as a second messenger [\(Carafoli,](#page-144-1) [2005\)](#page-144-1). Then Ca^{2+} Ca^{2+} cannot be metabolized like other second-messenger molecules. Hence cells tightly regulate intracellular levels through numerous bindings as well as specialized extrusion proteins [\(Clapham,](#page-145-4) [1995\)](#page-145-4).

Almost all eucaryotic cell types use both intracellular as well as extracellular resources of Ca^{2+} Ca^{2+} . The responsible regulating mechanisms for the influx of external Ca^{2+} Ca^{2+} are already well known to scientists [\(Carafoli,](#page-144-2) [2002\)](#page-144-2). As an example, voltage-operated channels in neurons assist action potentials by triggering the release of neurotransmitters at synaptic junctions. Neurotransmitters can establish an influx of Ca^{2+} Ca^{2+} utilizing receptor-operated channels primarily localized postsynaptically. Despite such well-established influx pathways, there is not much known on the mechanism of the intracellular Ca^{2+} Ca^{2+} supply of neurons although [IP](#page-135-6)3Rs allocated all over the [ER](#page-135-5) are accountable for the dispense of Ca^{2+} Ca^{2+} [\(Berridge,](#page-143-0) [2005\)](#page-143-0).

The [ER](#page-135-5) network in cells accounts for the dynamics of Ca^{2+} Ca^{2+} signaling by operating both as a sink as well as a source of Ca^{2+} Ca^{2+} . Such an internal storage of Ca^{2+} is possible to have a far reaching impact on Ca^{2+} Ca^{2+} signals in cells. Ca^{2+} can be localized within compartments in high levels or can spread across cells as widespread Ca^{2+} Ca^{2+} waves [\(Berridge,](#page-143-1) [1998\)](#page-143-1).

The [IP](#page-135-1)₃ signaling network (see Figure [2.4\)](#page-19-0) is highly sophisticated within cell tissues. A manifold amount of diverse receptors excite the hydrolysis of phosphat-idylinositol-4,5-bisphosphate ([PIP](#page-136-9)₂) into [IP](#page-135-1)₃ and sn–1,2–diacylglycerol ([DAG](#page-135-7)). Both are well known second messengers. [IP](#page-135-1)₃ then releases Ca^{2+} Ca^{2+} from the [ER](#page-135-5)'s $IP₃Rs$ $IP₃Rs$ [\(Berridge,](#page-143-2) [1993\)](#page-143-2).

Figure 2.4: The Ca^{2+} Ca^{2+} signaling pathway. When a hormone or neurotransmitter ("first messenger") interacts with a receptor on the cell membrane, $IP₃$ $IP₃$ is released within the cell, causing a calcium response. The [ER](#page-135-5) plays a decisive role in calcium regulation next to the extracellular space as a storage for the cytotoxic Ca^{2+} Ca^{2+} . Adapted from [Alberts](#page-142-3) [et al.](#page-142-3) [\(2003\)](#page-142-3).

By the fact that the intracellular Ca^{2+} Ca^{2+} release by IP_3Rs IP_3Rs is sensitive to a diversity of different factors makes it very complex. IP_3Rs IP_3Rs are most effectively activated when Ca^{2+} Ca^{2+} and [IP](#page-135-1)₃ are both presented at the same time. This dual activation has noteworthy effects for signal transduction mechanisms. Foremost of these is that the IP_3 IP_3 receptor may act as a cooccurrence detector due to the requirement for two separate messengers. Small levels of IP_3 IP_3 , unable to excite an immediate Ca^{2+} Ca^{2+} release, may enhance the [IP](#page-135-6)₃R's Ca^{2+} responsiveness. Thereby they change the cytoplasm into a medium capable of producing regenerative Ca^{2+} Ca^{2+} waves. An increasing Ca^{2+} Ca^{2+} gradient within the [ER](#page-135-5) can have a positive feedback by sensitizing the [IP](#page-135-6)₃Rs' [Ca](#page-134-4)²⁺ uptake. The increase of $|Ca^{2+}|\$ i represses IP₃Rs [\(Berridge,](#page-143-3) [2004\)](#page-143-3).

2.2.2 Two Pool Model

The two pool model presented in this section is based on the works of [Goldbeter](#page-148-0) [et al.](#page-148-0) [\(1990\)](#page-148-0); [Berridge](#page-143-4) [\(1991\)](#page-143-4). This model shows an oscillating behavior of Ca^{2+} Ca^{2+} concentrations very similar to phenomena discovered in living cells. If a diffusion term is applied to the system even [3D](#page-134-2) wave-fronts can be simulated.

The basis of the two pool model is that oscillations are set up through an interaction between two releasable pools of Ca^{2+} Ca^{2+} . Here it is assumed that the external signal triggers the synthesis of IP_3 IP_3 . The effect is simply a discharge of an intracellular pool of Ca^{2+} Ca^{2+} leading to a rise of cytosolic Ca^{2+} . A simple assumption is made that a constant influx of Ca^{2+} Ca^{2+} from the [IP](#page-135-1)₃-sensitive pool occurs as long as the stimulus is present. This affects the probability of the occurrence of oscillations exclusively by the cycling of Ca^{2+} Ca^{2+} between the cytosol and the IP_3 IP_3 -insensitive pool. The IP_3 -insensitive pool is therefore considered to remain constant as a result of a fast backfill by the influx of extracellular $Ca²⁺$ $Ca²⁺$ [\(Goldbeter et al.,](#page-148-0) [1990\)](#page-148-0). This autoregulatory mechanism is controlled by the content of Ca^{2+} Ca^{2+} in the [IP](#page-135-1)₃-sensitive pool and by the uptake of cytosolic Ca^{2+} Ca^{2+} after a spike. The magnitude of the influx, $v_1\beta$, from this [IP](#page-135-1)₃-sensitive pool is assumed to be proportional to the saturation function β of the [IP](#page-135-6)₃R. The cooperative nature of this saturation function is expressed implicitly in β. The level of [IP](#page-135-1)₃ is proposed to be caused by stimulation increases with the magnitude of the external signal. [IP](#page-135-1)₃ thus controls the flow of Ca^{2+} Ca^{2+} into the cytosol. This again assists the [IP](#page-135-1)₃-insensitive pool for releasing Ca^{2+} Ca^{2+} in oscillatory cycles [\(Berridge,](#page-143-4) [1991\)](#page-143-4).

The two variables of this model are the concentration of free Ca^{2+} Ca^{2+} in the [IP](#page-135-1)₃insensitive pool (e.g., the [ER](#page-135-5) or [SR](#page-136-5)) and in the cytosol. These two variables are denoted by Z and Y , respectively. If assumed that buffering is linear with respect to the Ca^{2+} Ca^{2+} concentration, then the time evolution of the system is driven by the following two kinetic equations [\(Goldbeter et al.,](#page-148-0) [1990\)](#page-148-0):

$$
\frac{dZ}{dt} = v_0 + v_1\beta - v_2 + v_3 + k_f Y - kZ, \qquad (2.1)
$$

$$
\frac{dY}{dt} = v_2 - v_3 - k_f Y. \tag{2.2}
$$

In the above equations, all rates and concentrations are defined with respect to the total cell volume. Here, v_0 and kZ relate to the influx and efflux of $Ca²⁺$ $Ca²⁺$ into and out of the cell. This occurs even in absence of external stimuli. These terms are assumed to be constant and linear for simplicity. The rate of [ATP](#page-134-9)-driven pumping of Ca^{2+} Ca^{2+} from the cytosol into the [IP](#page-135-1)₃-insensitive store is denoted with v_2 . In contrast v_3 represents the rate of transport from this pool into the cytosol. The term $k_f Y$ refers to a leaky, non-activated transport from Y into Z. This process was found to stabilize the amplitude of Ca^{2+} Ca^{2+} transients at different levels of stimulation.

An increase in IP_3 IP_3 is triggered the reception of an external signal. This in turn leads to a rise in the saturation function β and to a subsequent increase of cytosolic Ca^{2+} Ca^{2+} . The conditions in which this initial rise triggers Ca^{2+} oscillation can be determined by resorting to phase plane analysis. This is especially possible because the system of equations comprises only two variables. Here, it was indicated that the activation of v_3 by Z is most appropriate for inducing sustained oscillations upon external stimulation. This condition directly corresponds to an activation by cytosolic Ca^{2+} Ca^{2+} where Ca^{2+} is transported from the intracellular store into the cytosol. The two pool model predicts that at least in the absence of time delays, such a process does not satisfy the triggering of a sustained oscillatory response. When taking into account the cooperative nature of the pumping process, the Ca^{2+} Ca^{2+} release from the intracellular store, and the positive feedback performed by the latter transportation of cytosolic $Ca²⁺$ $Ca²⁺$, the rates v_2 and v_3 in the equations [2.1](#page-21-0) and [2.2](#page-21-0) take the following form [\(Goldbeter et al.,](#page-148-0) [1990\)](#page-148-0):

$$
v_2 = V_{M2} \frac{Z^n}{K_2^n + Z^n},\tag{2.3}
$$

$$
v_3 = V_{M3} \frac{Y^m}{K_R^m + Y^m} \times \frac{Z^p}{K_A^p + Z^p}.
$$
 (2.4)

In these equations, V_{M2} and V_{M3} denote respectively the maximum rates of $Ca²⁺$ $Ca²⁺$ pumping into and the release from the intracellular store. These processes are described by Hill functions whose cooperativity coefficients are taken as n and m. Here, p denotes the degree of cooperativity of the activation process. K_2 , K_R , as well as K_A are threshold constants for activation, pumping, and releasing.

The equations [2.1,](#page-21-0) [2.2,](#page-21-0) [2.3,](#page-21-1) and [2.4](#page-21-1) admit a unique steady-state solution. Linear stability analysis of these equations indicated that the steady state is not always stable. In the absence of stimulation, a situation is considered in which the system is initially in a stable steady state characterized by a low cytosolic Ca^{2+} Ca^{2+} level close to 0.1 μ M. The system reacts to an increase in β up to 30%, due to a rise in IP_3 IP_3 triggered externally. Here, an oscillation of cytosolic Ca^{2+} Ca^{2+} occurs. Such repeating spikes are accompanied by a sawtooth variation of the intracellular store's Ca^{2+} Ca^{2+} content. The period of the oscillations is of the order of 1s, as in some experimental systems. Periods of $1min$ or more are readily obtained if the kinetic parameters are divided by a factor of 10–100 [\(Kraus and](#page-151-2) [Wolf,](#page-151-2) [1992\)](#page-151-2). A spatio-temporal extension of the two pool model allows for the modeling of intercellular Ca^{2+} Ca^{2+} -waves. To make this happen a diffusion term has to be added to the system. In the model it is assumed that the IP_3 IP_3 -sensitive $Ca²⁺$ $Ca²⁺$ pools are only located near the membrane and the [IP](#page-135-1)₃-insensitive pools are spread all over the intracellular space.

The mathematical description of the deterministic methods account the system with a coupled set of nonlinear [ODE](#page-136-1)s of first order. [Kraus et al.](#page-151-3) [\(1992\)](#page-151-3); [Kraus](#page-151-2) [and Wolf](#page-151-2) [\(1992\)](#page-151-2) derived a stochastic model from this system, which is numerically traceable by means of a stochastic simulation. The stochastic method models the system through a master equation. With the first method the pools correspond to particle concentrations and the processes correlate with mathematical functions, which move into the pools via fluxes connected with concentrations. With differentiation of the master equation the pools contain a certain number of particles. Now the process describe transition probabilities between state transitions, which are connected by fluxes. The states are characterized by the overall occupation number of the pools. External entities comply in both cases with externally defined system parameters. While the deterministic method inspects the temporal trend of concentrations, the stochastic method describes changes in particle numbers of every particle species in contrast to that. Each reaction, where the number of particles changes, is simulated directly. The stochastic method therefore constitutes the microscopic view of the system, which is in diametric opposition to the deterministic - and thus

macroscopic - depiction.

The stochastic model was applied to the simulator described in this work for final a testing and comparison purpose of state-of-the-art modeling and simulation tools in Section [5.2.2.3.](#page-96-0) The mathematical formulation of the stochastic two pool model is displayed in Table [2.4.](#page-23-0)

Table 2.4: Stochastic two pool model. This table specifies the mathematical process of the stochastic two pool model. X and Y describe the number of Ca^{2+} Ca^{2+} ions in the cytosol or the [IP](#page-135-1)₃-insensitive pool. There is the assumption that the number of Ca^{2+} Ca^{2+} ions in the extracellular space and within the [IP](#page-135-1)₃-sensitive pools are kept constant. Adapted from [Kraus et al.](#page-151-3) [\(1992\)](#page-151-3).

Chapter [3](#page-24-0) now gives a closer look on currently used simulation software tools. Their architecture and utilized algorithms will be continued and compared in Chapter [5](#page-88-0) with a preceding formal description on [4DiCeS](#page-134-0) (see Chapter [4\)](#page-36-0).

Related Approaches and Tools

The relevance to model as well as to simulate biological systems was discussed in the previous chapter. The following chapters give an overview of the stateof-the-art in modeling and simulation in comparison to the new cell biology framework [4DiCeS](#page-134-0). Therefore, some of the inherent problems in characterizing the different facets of biological function are stated. This includes a brief overview of how functional information is currently represented in databases. And also prevailing applications for modeling and simulation of biochemical networks are introduced.

3.1 Database and Information Retrieval

The tremendous but valuable information gathered together in recent years has to be organized and pooled in databases. In this respect databases are widely deployed to store the relationships of biochemical systems. Currently there exist over 1000 biological databases [\(Galperin,](#page-147-1) [2008\)](#page-147-1) and about 45 databases supplying cellular signaling pathways at different levels of detail and complexity.

Public bio-molecular interaction databases are resources to basic building blocks of biological signaling pathways. Huge clusters of molecular interactions can be generated based only on this information. However, a molecular interaction cluster does not represent a signaling pathway per se. In effect, more information about each interaction, such as its outcome (e.g. activation as well as inhibition), is required for it to become a trustful component of a signaling pathway [\(Cary et al.,](#page-144-3) [2005\)](#page-144-3). Both public as well as private database initiatives have taken up the effort of creating biological pathway databases and providing computational biology tools for their analysis. Some of the databases focus on static (manually drawn) representations [\(Bhalla and Iyengar,](#page-144-4) [1999;](#page-144-4) [Sivakumaran et al.,](#page-156-2) [2003;](#page-156-2) [Trost,](#page-158-2) [2002\)](#page-158-2) whereas other systems support dynamic visualizations based on graph drawing algorithms [\(Fukuda and Takagi,](#page-147-2) [2001;](#page-147-2) [Fukuda et al.,](#page-147-3) [2004\)](#page-147-3)). There is also a variety of databases specialized in molecular pathways with physical parameters as rate constants and concentrations [\(Igarashi and Kaminuma,](#page-150-4) [1997\)](#page-150-4). In addition to the previously described pathway databases, there exist databases containing detailed information regarding characterized enzymatic reactions. Additional links to other databases provide useful information on involved enzymes and biochemical reactions [\(Gough and](#page-148-1) [Ray,](#page-148-1) [2002;](#page-148-1) [Gough,](#page-148-2) [2002\)](#page-148-2).

Currently three simulation model repositories serve actively in the internet – namely the Cell Markup Language ([CellML](#page-134-11)) [\(Lloyd et al.,](#page-152-5) [2004\)](#page-152-5), the JWS Online [\(Olivier and Snoep,](#page-154-2) [2004\)](#page-154-2), and the BioModels (Novère et al., [2006\)](#page-154-3) repository (see Table [3.1\)](#page-25-0).

Table 3.1: Model repositories. The simulation model repositories of the two most prominent modeling languages [SBML](#page-136-0) and [CellML](#page-134-11) contain models on metabolic networks, cell cycle and cellular signaling. The (x)Cellerator sites provide example models as Mathematica (*.nb) files.

The [SBML](#page-136-0) [\(Hucka et al.,](#page-149-0) [2004\)](#page-149-0) repository ceased work at the end of 2005. The E-Cell project [\(Takahashi et al.,](#page-157-1) [2003\)](#page-157-1) has plans for its own model repository, however, there is no concrete data available on the internet at present.

Further information on biological pathway databases can be retrieved from the Nucleic Acids Research database issues [\(Baxevanis,](#page-143-5) [2000,](#page-143-5) [2001,](#page-143-6) [2002,](#page-143-7) [2003;](#page-143-8) [Galperin,](#page-147-4) [2004,](#page-147-4) [2005,](#page-147-5) [2006,](#page-147-6) [2007,](#page-147-7) [2008\)](#page-147-1) and by the Pathway Resource List $(PRL)^1$ $(PRL)^1$ $(PRL)^1$ $(PRL)^1$ – a database that contains information on over 240 internet pathway resources. Most of these resources are databases themselves containing protein–protein interactions, metabolic reactions, or cellular signaling. The [PRL](#page-136-10) provides resource links and is building up additional information such as the amount of data and the organism coverage within each pathway resource [\(Bader et al.,](#page-142-4) [2006\)](#page-142-4).

¹Pathway Resource List: <http://www.cbio.mskcc.org/prl/>

3.2 ________ Modeling and Simulation Software

The preceding section gave a brief introduction to an overwhelming amount of data repositories present for use in systems biology. The amount of simulation tools dealing with biochemical reaction and diffusion systems is not quite as huge, but is still plentiful. Therefore, this section deals with the description of only the most important applications of this category (see Table [3.2\)](#page-27-0).

A simulation tool is defined as an application performing time series simulation of predefined mathematical models. In contrast a design or modeling tool is applied for building a model graphically. Often simulation tools bring an attached design tool along. If not then models can either be described by markup or scripting languages [\(Pettinen et al.,](#page-155-0) [2005\)](#page-155-0).

The simulation software tools can be categorized into either deterministic, stochastic, or hybrid (deterministic and stochastic) programs. Other categories apart from the algorithmic approaches are the modeling of either [2D](#page-134-3) or [3D](#page-134-2) geometry and the separation of programs into either stand-alone tools or frameworks.

One of the very first programs available for reaction simulations was the GEneral PAthway SImulator ([Gepasi](#page-135-8)) [\(Mendes,](#page-152-6) [1993\)](#page-152-6). It translates biochemical reaction equations into coupled [ODE](#page-136-1)s which in turn are then solved numerically. Thus the [Gepasi](#page-135-8) system represents a purely deterministic approach such as BIOCHemical Abstract Machine ([BIOCHAM](#page-134-12)) [\(Calzone et al.,](#page-144-5) [2006\)](#page-144-5), Cellerator [\(Shapiro et al.,](#page-156-3) [2003\)](#page-156-3), E-Cell, the Python Simulator for Cellular Systems ([PySCeS](#page-136-11)), and VirtualCell do as well (see Section [3.2.1\)](#page-28-0). The deterministic simulators Genesis and Neuron were originally designed to model neurons and neuronal networks but have shown that cell signaling simulations are just as adequate [\(Bhalla,](#page-143-9) [2002\)](#page-143-9). Xyce is actually a deterministic massively parallel simulator for electronic circuits that was used for solving biochemical problems [\(Schiek and May,](#page-156-4) [2003\)](#page-156-4). Simulators such as the Stochastic Simulator ([StochSim](#page-136-12)), the [MC](#page-135-2) Simulator of Cellular Microphysiology ([MCell](#page-135-0)) (see Section [3.2.2\)](#page-30-0), and Mesoscopic Reaction Diffusion (simulator) ([MesoRD](#page-135-9)) [\(Hattne et al.,](#page-149-4) [2005\)](#page-149-4) implement stochastic algorithms only. Examples for hybrid simulators are the Bio-chemical Network (stochastic) Simulator ([BioNetS](#page-134-13)) [\(Adalsteinsson et al.,](#page-142-5) [2004\)](#page-142-5), the [PySCeS](#page-136-11) [\(Olivier et al.,](#page-154-4) [2005\)](#page-154-4), WebCell [\(Lee et al.,](#page-151-4) [2006\)](#page-151-4), and xCellerator.

Very recently, efforts have been made to mix various approaches in order to obtain either the combination of many tools in one software package [\(Rost and](#page-155-6)

Table 3.2: Simulation and modeling environments. This table presents a subset of software tools available today for cellular modeling and simulation. Shown here are 20 out of more than 80 applications. The simulators [BioNetS](#page-134-13), [Copasi](#page-134-15), [MCell](#page-135-0), SmartCell, and [StochSim](#page-136-12) are stochastic applications. The [SBW](#page-136-13) and the [Bio-SPICE](#page-134-14) are actually simulation frameworks rather than programs. The remaining applications within this table are based on deterministic modeling and simulation. A special position take [XmdS](#page-136-14) and [BioNetS](#page-134-13) as they are code generators. The designations superscripted with an asterisk ('∗') are going to be dis[cussed in more detail in the following sections.](#page-155-6)

[Kummer,](#page-155-6) [2004\)](#page-155-6), e.g. the Complex Pathway Simulator ([Copasi](#page-134-15)) (see Section [3.2.2\)](#page-30-0) or a tool offering access to many different software packages, e.g. the Systems Biology Workbench ([SBW](#page-136-13)) (see Section [3.2.3\)](#page-32-0). A very special position takes the eXtensible multi-dimensional Simulator ([XmdS](#page-136-14)) and [BioNetS](#page-134-13) as they actually are C++ code generators [\(Collecutt and Drummond,](#page-145-5) [2001\)](#page-145-5). If their code is compiled then the resulting programs are simulation applications by their own again. Further information on existing modeling and simulation applications can be retrieved by the [SBML](#page-136-0) Software Guide^{[2](#page-28-1)} – a matrix that contains information on software providing support for [SBML](#page-136-0).

The following three sections give a closer look to ten well reputed applications and two frameworks. The last Section [3.2.4](#page-33-0) then defines comparison criteria for the comparison of the ten programs. The [4DiCeS](#page-134-0) approach, which will be presented in the subsequent Chapter [4,](#page-36-0) is going to be brought into context with Section [5.3.](#page-98-0)

3.2.1 Deterministic

This section presents five well known and used deterministic simulation applications. Although [Gepasi](#page-135-8) ceased further development, it is still in use to date and has played a major role in the development of all the other tools discussed here. There are plans to completely replace [Gepasi](#page-135-8) by the newer [Copasi](#page-134-15) (see Section [3.2.2.2\)](#page-31-0). Excepting [Gepasi](#page-135-8) all other described simulators are still under development and have a user community of their own.

3.2.1.1 [BIOCHAM](#page-134-12)

[BIOCHAM](#page-134-12) is a programming environment for modeling biochemical systems, making simulations, and querying the model in temporal logic. It provides a rule-based language for modeling biochemical systems, a simulation engine, and a query language based on temporal logic, Computational Tree Logic ([CTL](#page-135-11)), or Linear Temporal Logic ([LTL](#page-135-12)). A machine learning system is provided for correcting and completing models either by changing rules with respect to a [CTL](#page-135-11) specification or by estimating parameters of an [LTL](#page-135-12) specification. An interface to the symbolic model checker (NuSMV) is provided also. [BIOCHAM](#page-134-12) was initiated by the Constraint Programming group of The National Institute for Research in Computer Science and Control (INRIA) at Rocquencourt, France.

² [SBML](#page-136-0) Software Guide: [http://sbml.org/SBML](http://sbml.org/SBML_Software_Guide/) Software Guide/

3.2.1.2 E-Cell

The E-Cell project is based on international research aiming to model and reconstruct biological phenomena in silico and to develop necessary theoretical supports, technologies and software platforms to allow precise whole cell simulation [\(Tomita et al.,](#page-157-2) [1997\)](#page-157-2). The E-Cell Model Language ([EML](#page-135-13)), a subset of the eXtensible Markup Language ([XML](#page-136-15)), is used for describing the models. The [SBML](#page-136-0) support was also included to enable a wide cross-platform model exchange [\(Tomita et al.,](#page-157-3) [1999\)](#page-157-3). The E-Cell project is managed by the Institute for Advanced Biosciences, Laboratory for Bioinformatics, Fujisawa and the Mitsui Knowledge Industry, Bioscience Division, in Tokyo, Japan.

3.2.1.3 General Pathway Simulator ([Gepasi](#page-135-8))

The GEneral PAthway SImulator is one of the first software packages for modeling biochemical systems. [Gepasi](#page-135-8) simulates the biochemical reaction kinetics, provides a number of tools to fit models to existing data, optimizes the functions of the models, and performs a metabolic control analysis and a linear stability analysis [\(Mendes,](#page-152-6) [1993\)](#page-152-6). The application simplifies the task of model-building by assisting the user in automatically translating given reactions into matrices and differential equations transparently [\(Mendes,](#page-152-7) [1997\)](#page-152-7). This is combined with a set of numerical algorithms that ensure fast and accurate results [\(Mendes and](#page-152-1) [Kell,](#page-152-1) [1998\)](#page-152-1). It was developed at the Virginia Bioinformatics Institute, USA, as a pure deterministic modeling and simulation environment.

3.2.1.4 VirtualCell

The National Resource for Cell Analysis and Modeling at the University of Conneticut Health Center, in Conneticut, USA, created a remote user simulation and modeling application. A general purpose differential equation solver is used to translate the initial biological description into a set of differential equations [\(Loew and Schaff,](#page-152-8) [2001\)](#page-152-8). The generated results are stored on a remote server and can be reviewed and exported into various formats. The compartments represent [3D](#page-134-2) volumetric regions, while the membranes represent [2D](#page-134-3) surfaces separating the compartments [\(Schaff et al.,](#page-156-5) [1997;](#page-156-5) [Schaff and Loew,](#page-156-6) [1999\)](#page-156-6). The geometry may be captured by various imaging modalities, such as wide field, confocal, or electron microscopy [\(Slepchenko et al.,](#page-156-7) [2002\)](#page-156-7).

3.2.1.5 xCellerator

The xCellerator is the successor to the Cellerator package, which was designed as an interface to Wolfram Mathematica for facilitating biological modeling by automated equation generation. It provides tools for generating, translating, and numerically solving a potentially unlimited number of biochemical interactions [\(Shapiro et al.,](#page-156-3) [2003\)](#page-156-3). xCellerator solves the complete set of equations predicted by the law of mass action. The package also contains a number of transcriptional regulation models that are not Michaelis-Menten equations. xCellerator may write its results in [SBML](#page-136-0). The Cellerator package was formerly developed at the National Aeronautics and Space Administration's (NASA) Jet Propulsion Laboratory, California, USA, and is now privately continued.

3.2.2 Stochastic

The simulation applications shown in this section are stochastic approaches [\(Kibby,](#page-150-5) [1969\)](#page-150-5). They either implement proprietary algorithms of their own ([MCell](#page-135-0) and [StochSim](#page-136-12)) or make use of simulation algorithms from literature ([BioNetS](#page-134-13), [Copasi](#page-134-15), and SmartCell). [BioNetS](#page-134-13) and [Copasi](#page-134-15) are also utilizing numerical solvers in combination to their stochastic algorithms.

3.2.2.1 Biochemical Network Stochastic Simulator ([BioNetS](#page-134-13))

[BioNetS](#page-134-13) was developed at the University of North Carolina at Chapel Hill, USA. It was designed to simulate biochemical network models in a hybrid, stochastic and deterministic manner. The type of used discrete or continuous random variable for each chemical species in the network can be specified individually. The package was implemented to efficiently scale with any network size to allow the study of large systems. [BioNetS](#page-134-13) is available as a stand alone package but runs also as a Bio–Simulation Program for Intra- and Inter-Cell Evaluation ([Bio-](#page-134-14)[SPICE](#page-134-14)) (see Section [3.2.3.1\)](#page-32-1) agent. The output of the software is portable $C/C++$ code that may be compiled and run on any system with the appropriate compiler [\(Adalsteinsson et al.,](#page-142-5) [2004\)](#page-142-5).

3.2.2.2 Complex Pathway Simulator ([Copasi](#page-134-15))

The [Copasi](#page-134-15) project is based on [Gepasi](#page-135-8) and a program for the automatic parsing and STochastic simulation of [ODE](#page-136-1)s ([STODE](#page-136-16)) [3](#page-31-1) [\(van Gend and Kummer,](#page-158-3) [2001\)](#page-158-3). [Copasi](#page-134-15) incorporates a model generator, user-friendly visualization platforms, optimization routines, methods from non-linear dynamics, and different simulation techniques. [Copasi](#page-134-15) is supervised by Pedro Mendes (Bioinformatics Institute, Virginia, USA) along with Ursula Kummer of the European Media Laboratory, Heidelberg, Germany. This application is planned to enable the simulation of complex metabolic processes in cells without having to master complex mathematical or computer skills [\(Rost and Kummer,](#page-155-6) [2004\)](#page-155-6).

3.2.2.3 [MC](#page-135-2) Simulator of Cellular Microphysiology ([MCell](#page-135-0))

[MCell](#page-135-0) is a modeling application for simulations of cellular signaling in complex [3D](#page-134-2) subcellular micro-environments. Optimized [MC](#page-135-2) algorithms are used to track the stochastic behavior of discrete molecules in space and time. These particles diffuse and interact with other heterogeneously distributed molecules within the [3D](#page-134-2) geometry [\(Bartol Jr. et al.,](#page-143-10) [1996\)](#page-143-10). All simulation components are defined by using a specific programming language called Model Description Language ([MDL](#page-135-14)) [\(DeSchutter and Cannon,](#page-145-6) [2000\)](#page-145-6). The project was initiated by the Computational Neurobiology Laboratory at the Salk Institute for Biological Studies, in California, and by the Pittsburgh Supercomputing Center's working group for Biomedical Applications, Pennsylvania, USA.

3.2.2.4 SmartCell

SmartCell is a general tool for modeling and simulation of reaction and diffusion pathways within cells. It supports diffusion and localization by using a mesoscopic stochastic reaction model. The SmartCell package should handle various cell geometries, allows the localization of species, supports desoxyribonucleic acid ([DNA](#page-135-15)) transcription and translation, membrane diffusion and multi-step reactions, as well as cellular growth. Moreover, different temporal and spatial constraints can be applied to the model [\(Ander et al.,](#page-142-6) [2004\)](#page-142-6). It is expected to provide a suitable model description format [\(Nasi,](#page-153-0) [2004\)](#page-153-0). SmartCell is a project of the European Molecular Biology Laboratory, Heidelberg, Germany.

³ [STODE](#page-136-16): <http://atlas.villa-bosch.de/bcb/software/Carel/>

3.2.2.5 Stochastic Simulator ([StochSim](#page-136-12))

[StochSim](#page-136-12) was written by Carl Firth as part of his PhD work at the University of Cambridge [\(Morton-Firth,](#page-153-1) [1998\)](#page-153-1). It was developed as part of a study of bacterial chemotaxis as a more realistic way to represent the stochastic features of this signaling pathway. It is able to handle large numbers of individual reactions encountered [\(Morton-Firth and Bray,](#page-153-2) [1998;](#page-153-2) [Morton-Firth et al.,](#page-153-3) [1999\)](#page-153-3). The application consists of a platform-independent core simulation engine. The program encapsulates the algorithm described above as well as a separate graphical user interface. [StochSim](#page-136-12) represents individual molecules or molecular complexes as individual software objects (Novère and Shimizu, [2001\)](#page-154-5).

3.2.3 Frameworks

This section introduces frameworks rather than applications in comparison to the two previous sections. Such frameworks are developed to accomplish very particular problems. They provide entire workbenches of interfaces that allow for the integration of many tools and methods to interact with each other.

3.2.3.1 Bio–Simulation Program for Intra- and Inter-Cell Evaluation ([Bio-SPICE](#page-134-14))

The [Bio-SPICE](#page-134-14) toolkit was developed to model and simulate cellular dynamic networks. Contributed modules are organized in the [Bio-SPICE](#page-134-14) dashboard, which is a graphical user environment [\(Sauro et al.,](#page-156-8) [2003\)](#page-156-8). It permits data sources, models, simulation engines, and output displays provided by different investigators, and running on different machines, to work together across a distributed, heterogeneous network [\(Garvey et al.,](#page-147-8) [2003\)](#page-147-8). Among several other features, the environment enables users to create a graphical workflow by configuring and connecting available [Bio-SPICE](#page-134-14) components [\(Kumar and Feidler,](#page-151-5) [2003a,](#page-151-5)[b\)](#page-151-6). The project was initiated by the Defense Advanced Research Projects Agency Information Processing Technology Office, Virginia, USA.

3.2.3.2 System Biology Workbench ([SBW](#page-136-13))

The [SBW](#page-136-13) enables different tools to interact with each other. The framework supports tools written in different programming languages, which may run on different platforms and physical machines [\(Sauro et al.,](#page-156-8) [2003\)](#page-156-8). The aim is to facilitate collaboration among developers of systems biology software. Developers should find it easier to build an [SBW](#page-136-13) interface than to recreate functionality. They can then concentrate on developing best-of-breed solutions in the areas where they have special expertise [\(Hucka et al.,](#page-149-5) [2002\)](#page-149-5). Both [SBW](#page-136-13) and [SBML](#page-136-0) are being developed in collaboration with several groups developing simulation packages as described in the last two sections.

3.2.4 Comparison

While the described simulation tools thus have their benefits, none have so far addressed all the currently emerging research problems. The efforts in the field of cellular simulation can be roughly categorized as stand-alone modeling and simulation tools or extendable frameworks. The first can then be divided again into either more or less strict deterministic or stochastic methods. Even though the stand-alone tools often provide software interfaces, frameworks broaden this ability to include new technologies to the system.

Specific interests of research groups often have great influence on the development of simulation applications. The usability of a tool is highly affected by user requirements [\(Schwehm,](#page-156-9) [2001\)](#page-156-9). Here the chosen operating system and the selection of a Graphical User Interface ([GUI](#page-135-16)) versus a scripting or batch mechanism are key features. Programs currently used vary noticeably in their applicability for specific types of modeling [\(Pettinen et al.,](#page-155-0) [2005\)](#page-155-0). The general usability indicated by the user's learning curve and application-provided model designers have here a great impact. Well featured textual or graphical utilities often aid the modeling process significantly. The number and quality of the utilizable algorithms, supported model-exchange formats, and additional features round up such user requirements dramatically. Next to this, it is of major importance to have a tool with reliable and precise results at optimal performance. An extension mechanism for the integration of new features and functionality is then crucial. The support of spacial modeling information is necessary for close to reality simulations. And last but not least it is important to segregate parts of the model and handle such parts differently.

The following Table [3.3](#page-34-0) will therefore define 12 comparison criteria extracted form the preceding paragraphs. These criteria are applied to the previously introduced related applications in a three-state manner. The three states range from unsatisfactory to sufficient. This will allow for an easier qualification of the differences among the tools. Table [3.4](#page-35-0) is then going to summarize all

comparison findings from that definition.

Table 3.3: Comparison criteria. The defined comparison criteria were extracted from an objective analysis of state-of-the-art simulation applications and common software quality assurance considerations. They are further on used in a three-state manner (unsatisfactory, average, and sufficient). Table [3.4](#page-35-0) will display the three states as circles from empty to filled.

As can be seen there is not one application making up for all defined comparison criteria. The concurrency feature is left out, because it is not supported by any of the comparison candidates. When weighting the three-state [unsatisfactory (\circ) , average (\circ and \circ), and sufficient (\circ) with zero (0), a half (0.5), and one (1) then the given applications range from four to eight criteria points. A new application as of [4DiCeS](#page-134-0) (see Section [5.3.3](#page-105-0) for a comparison) should at least have an equal or even higher level to the best to keep up with or even outperform the state-of-the-art.

The following Chapter [4](#page-36-0) will now formally describe the [4DiCeS](#page-134-0) approach.

Table 3.4: Comparison of simulators. Filled $\left(\bullet \right)$ circles are generally superior to their half-full (\bullet) counterparts. Empty (\circ) circles indicate the need for improvement or a total absence. In methodology hybrid approaches are considered most sufficient. Designers are either defined as textual (Q) , graphical (Q) , both (Q) , or none (Q) at all. Concurrency was omitted, due to missing support by any of the comparison candidates. 30
Definitions and Implementation

Based on the criteria described in the previous chapter with respect to already existing solutions, a general concept has been designed for the implementation of a [4D](#page-134-0) Cell Simulator ([4DiCeS](#page-134-1)). The system is developed as a common and extensible framework flexible enough and well suited to serve as a platform in systems biology. The four dimensions describe here the utilization of the three space coordinates and a time axis. Cellular phenomena can therefore be simulated in a [3D](#page-134-2) geometric environment over time. The current project logo is presented in Figure [4.1.](#page-36-0) It consists of four rolling dice that represent the ability to model and simulate stochastically in [4D](#page-134-0) space.

Figure 4.1: The [4DiCeS](#page-134-1) logo. The four rolling dice stand for modeling and simulating by the use of [4D](#page-134-0) stochastic methods. Also, the dice could be compared to the cubes of the [3D](#page-134-2) geometry used with [4DiCeS](#page-134-1).

The aim of this project is to provide a modeling and simulation application to the user that is most adaptive in its functionality. The main features can be roughly summarized as seven main design characteristics (see Table [4.1\)](#page-37-0).

The first four features of Table [4.1](#page-37-0) have a great impact on the modeling and simulation core itself. First of all [4DiCeS](#page-134-1) provides a [3D](#page-134-2) simulation geometry that makes use of either imported or user-defined cellular compartments [\(Oleson](#page-154-0) [et al.,](#page-154-0) [2003\)](#page-154-0). Thus modeling and simulation will be based on such a geometry that in turn can then be easily reduced to a [2D](#page-134-3) or one-dimensional ([1D](#page-134-4))

Table 4.1: [4DiCeS](#page-134-1) feature list. These seven features describe the main goal of the design and implementation of [4DiCeS](#page-134-1). An eighth feature is the definition of a strong and well-defined software interface system (see Section [4.2\)](#page-50-0). This then should give the application users a modeling and simulation utility that can be easily adopted or even extended to their needs.

model if needed (Möller et al., [2002\)](#page-153-0). Possible sources for importing geometrical data could be either cellular image stacks such as from Confocal Laser Scanning Microcopy ([CLSM](#page-134-5)) and Multi-Photon Fluorescence Microscopy ([MPFM](#page-135-1)), or topological information such as from Scanning Electron Microscopy ([SEM](#page-136-0)) and Atomic Force Microscopy ([AFM](#page-134-6)). Such data has to be adapted to allow for entirely sealed compartments on which algorithms can work on within simulations. This already implies other features as well. Both reaction and diffusion algorithms should be applicable to the system (Möller et al., 2003). Diffusion must work for either freely diffusible particles or for membrane-bound particles. This means that it is either possible to apply a free [3D](#page-134-2), or a lateral [2D](#page-134-3) membrane diffusion to simulated particles [\(Oleson et al.,](#page-154-1) [2002\)](#page-154-1). Additionally definable boundary conditions are essential to reaction and diffusion by other

means as well. It must be possible to have different reaction-diffusion systems in unequal compartments of the same simulation environment. There is also the need for particles to undergo transmembrane diffusion under definable conditions. As another feature all reaction and diffusion algorithms can be of diverse types including either deterministic or stochastic methods. In doing so it should then be possible to concurrently run different algorithms on varying compartments. Single compartments can be handled deterministically whilst others can use stochastic methods in the same simulation environment [\(Oleson](#page-154-2) [et al.,](#page-154-2) [2006\)](#page-154-2).

The remaining three requirements of Table [4.1](#page-37-0) give more attention to the usability of the application as is. In this case it is important to provide diverse import and export filters of well-established storage formats such as [CellML](#page-134-7) and [SBML](#page-136-1) (see Section [4.2.2\)](#page-60-0). This should enable the exchange of existing models with the [4DiCeS](#page-134-1) system and then back to model repositories again. Thus simulations should run on either [GUI](#page-135-0)-based graphical frameworks or as Command-Line Interface ([CLI](#page-134-8))-based (i.e. textual) batch jobs. This allows for both high-throughput simulation applications and experimental model evaluation simulations [\(Oleson et al.,](#page-154-3) [2004\)](#page-154-3). The whole architecture must also enable other programming or scripting languages to effortlessly extend the [4DiCeS](#page-134-1) system. All these requirements depend on a well-defined plug-in mechanism that can be easily maintained and used by the system's core. This again is beneficial for third party programmers because they may also use the same plug-in concept to integrate their own ideas readily [\(Oleson et al.,](#page-154-2) [2006\)](#page-154-2).

It has to be noted that within this work the word 'interface' is mainly used for connection sockets between different parts of the application. In the case of human-machine interfaces the term 'user interface' is going to be applied.

The following sections will now subsequently describe the [4DiCeS](#page-134-1) application in a top-down manner. In connection with the general overview (see Section [4.1\)](#page-39-0) the implemented system will be eventually broken down into its functional parts. These functional units are then discussed at more detail later on. Anyway, each level of detail is going to comprise an overview, its dynamics, and dependencies (see Section [4.2\)](#page-50-0); therefore every subsection will define the environment of the level of detail, its components, the possible states, and the error-case handling, if applicable. This chapter finally closes with remarks to the deployment, the applicability, and the availability of the [4DiCeS](#page-134-1) system to the public. This includes specific programming design decisions made for applied programming languages, third party [API](#page-134-9)s, and further technologies. Noteworthy technicalities to concrete implementations and a general complexity consideration are then also referenced to Section [4.3.](#page-79-0)

The formal description of the [4DiCeS](#page-134-1) application presented in this chapter makes extensive use of the Unified Modeling Language ([UML](#page-136-2)). The [UML](#page-136-2) diagrams and types used in this work are therefore described in Appendix [C](#page-128-0) in further detail.

4.1 General Overview

This section will examine the general structure of the [4DiCeS](#page-134-1) framework. Both [UML](#page-136-2) use case studies (see Appendix [C.1\)](#page-128-1) and [UML](#page-136-2) static program structures (see Appendix [C.2\)](#page-130-0) provide an introduction of the overall architecture and features of the system. In doing so it is thereafter easier to comprehend with the system's dynamics.

4.1.1 Survey of Integral Parts

The primary goal for [4DiCeS](#page-134-1) is to provide a modeling and simulation environment capable of interpreting a given data model and running simulation iterations therewith. For this purpose the application framework was divided into functional pieces. Whereas the central part of the program is the [4DiCeS](#page-134-1) kernel, which manages incoming as well as outgoing events and monitors the stability of the system and all its units.

To allow for reusability, portability, and thus effectiveness of the [4DiCeS](#page-134-1) application, a consistent system of kernel interfaces (see Section [4.2\)](#page-50-0) was established. These kernel interfaces offer a centralized gateway to and from simulation algorithms, model data, and the actual user interface. The general concept is displayed in Figure [4.2](#page-40-0) for clarification, which displays a static [UML](#page-136-2) class diagram (see Appendix [C.2\)](#page-130-0).

4.1.1.1 Application Kernel

The [4DiCeS](#page-134-1) kernel provides functions for creating a simulation model independent of a specific application. Through the user interface, either the user or another program takes access to these functions. Input and output devices form a module, which is linked by the module interface with the kernel. The

Figure 4.2: Established interface system. The kernel, as the central part of the simulator, is surrounded by interfaces for simulation algorithms, model data input and the user interface integration. The extending modules to all interfaces are programming language independent. First implementations include the use of Java modules as can be seen from the Java typical names ending in '-let'.

kernel maps its functions at the user interface to the various capabilities of the different modules. Then the realization of the device independent from the requirement of single modules is accomplished by individual module drivers. The kernel interfaces are defined independent of special programming languages. The utilization of the underlying kernel therefore requires an adaptor interface instance for particular programming or scripting languages. This can be easily illustrated by the separation of the kernel into layers (compare to Figure [4.3\)](#page-41-0). Application-dependent layers may then build upon any individual programming language layer.

Figure 4.3: Kernel layer model. Very similar to an operating system layered model, the [4DiCeS](#page-134-1) kernel is separated into layers of different disposability. This structure allows for connecting various programming and scripting languages to the functions and interfaces provided by the kernel.

4.1.1.2 Data Model

All utilized data structures of [4DiCeS](#page-134-1) are integrated as the [4DiCeS](#page-134-1) application data model. The data model itself is not used for a simulation model representation directly but to a much bigger extent for testing and simulation purposes. This scheme follows the overall aim to understand, visualize, and experiment with the underlying workflow of the application. The data model does not necessarily include implicit simulation model information. Such can rather be extracted out of the application's data model if the components from Table [4.2](#page-42-0) are given.

4.1.2 Function Classes

To allow for any state-changing functionality within the kernel functions are required that address the design elements, attributes, modules, data structures as well as formats, and the human-machine interaction. All error conditions are processed by the kernel with its own exception handling and logging mechanism. The description of kernel functions includes the function identifiers, names, numbers, data-types, and the meaning of every parameter. Additionally the functions' effect(s) and possible error conditions are contained also. An effect of a specific kernel function could be either to change the underlying model, to return values to the user, to change the kernel's state, or a mixture of all three.

The state of the kernel is given by values from a set of state lists that are associated to the kernel itself, to modules, and to segments. The data structure of the state lists with its data types, its intent, and the preallocation of its

Table 4.2: Extraction of simulation model information. Since the application's data model does not necessarily include the simulation model itself there must be a way to extract such information. If the basic data is given with its dependencies, identifiers, and data extraction rules the simulation data can be extracted explicitly from the application's data model.

contained values are defined by kernel specifications. To avoid the inclusion of the full set of functions provided by the kernel, all functions are organized into levels of duty according to their increasing requirements. They are built as matrices of three levels for the input, the output, as well as the segmentation.

A brief overview of the kernel functions' classes is presented in the following. In subsequent sections, significant functions and interfaces will be considered in greater detail. The realizations of concrete modules will then be described in Section [4.3.](#page-79-0)

Design Elements: Design elements are the elementary components from what a simulation model is constituted of. In [4DiCeS](#page-134-1) there are particles, reactions, kinetics, and compartments. Particularly for [3D](#page-134-2) modeling, the design elements voxel and grid matrix are of special importance. A voxel is an indivisible and therefore distinct volume element of the overall grid matrix. The grid matrix is the simulation space. For addressing platform or language dependent (i.e. non-standardized) functions, the "general modeling primitive" is defined.

Design Attributes: A design element can be characterized by its geometrical form but also by its appearance. This is described by a set of design attributes. For voxels these can be position, size, volume, and their particle incorporation.

The grid matrix offers attributes that depict the grid type, number of voxels, and compartment details. Then the compartments are segmentations of the grid matrix providing details on the voxels owned and boundary conditions.

Modules: A significant component of the [4DiCeS](#page-134-1) architecture is the concept of interface modules. These allow for language-independent programming of simulation applications and plug-ins. A module is an abstraction of interfaces to the [4DiCeS](#page-134-1) kernel implementation. Also modules can be processed in sequence or in parallel. A design element instance may vary in display and in execution according to the special abilities of a connected module. The application is able to request features of a connected module through an underlying event-handling mechanism.

Model Structure: Interactive applications also require the capability for manipulating parts of the given model structure. There must exist opportunities to handle all design elements as separate entities when constructing a model. Therefore, a model can be composed of sub-models, "segments". Segments may be individually manipulated, deleted, and processed. [4DiCeS](#page-134-1) manages segments from within its segment repository.

Human-Machine Interaction: An application can request input values from a user through a module. [4DiCeS](#page-134-1) is currently aware of five classes that can be connected dynamically to the system (see Table [4.3\)](#page-43-0).

Table 4.3: Module classes. The input description for both geometry as well as activity, the algorithm description for both reactions as well as diffusion, and the programmatic interaction interfaces provide the flexible base for designing models in [4DiCeS](#page-134-1). They are designed to be accessed independently of programming language, and to run concurrently when appropriate.

The flexible control of these module classes by the kernel in combination with the segmentation provide the base of designing models without a loss of programming language dependency. The concurrency of algorithms gives then the basis for the segmentation of models and their simulation in parallel.

Model Format: [4DiCeS](#page-134-1) contains an interface to the model, which allows for both the long-term storage as well as the reload of models. The model file can either be used to persist models, for their transmission, or for a resume of a previously interrupted simulation session. Models are stored including their segmentation structure.

Event-Handling: All values in the tables, which describe the system's state, can be retrieved by calling functions. There exist both mechanisms for actively calling for information (synchronous events) and obtaining data from a model automatically (asynchronous events) when a certain trigger condition is reached.

Error-Handling: If an error occurs during the execution of a [4DiCeS](#page-134-1) function, then the kernel is set to a special error-handling state. Depending on the severity of the error, an error-handling routine of the used module can be invoked, or the kernel can try to solve the fault by itself. In either case an error-log file will be produced for further maintenance purposes.

4.1.3 Use-Case Descriptions

In the following section a use-case scenario is presented for the general use of the application. The input as well as the output, the simulation handling, and the algorithm management will be emphasized. All these components will then be described in more detail later on.

4.1.3.1 User Perspective

A general approach to [4DiCeS](#page-134-1) is the description of its main features by a use-case scenario that shows the user's perspective of the system. Figure [4.4](#page-45-0) presents a general overview [UML](#page-136-2) use-case diagram of the simulator, while Table [4.4](#page-45-1) provides further details to the specific use-case.

Here the [UML](#page-136-2) use-case actor is a potential user of the simulation program. This user has a direct interface to three essential use-cases to the system: the simulation engine, the input data, and the output data. In the center of the diagram

Figure 4.4: Use-case: User perspective.

Table 4.4: Use-case: User perspective. Both this table and Figure [4.4](#page-45-0) show the user's perspective of the [4DiCeS](#page-134-1) system. The figure displays a [UML](#page-136-2) diagram, while this table provides further information about the requirements, assumptions, and possible error-cases on the usecase.

there is the application data model, which is fed with information from both the descriptions of the geometries as well as the activities that were uploaded by the actor of the system. A geometry description is a [3D](#page-134-2) mesh of compartments subdividing the simulation space into distinct vessels for both reaction and diffusion events. These events in turn are defined by the activity description. The upper part of the use-case diagram represents the model description and definition part of the application. The lower part on the contrary shows use-cases utilizing these model boundaries for further simulation. Both reaction as well as diffusion algorithms perform using the given rules and produce output data in turn.

The user is required to feed the simulator with model definitions and algorithms before a simulation can be started. Then the production of output depends on at least one successful iteration of the simulation. Another point of possible conflict is the dependencies between model description and simulation description. The definition files must either be general enough to cope with diverse types of algorithms or preferred algorithms must also be defined. The precise location of every algorithm used has to be defined before the simulation is attempted, as is stipulated with the criteria that algorithms may run concurrently.

4.1.4 Program Dynamics

This section provides a more in-depth examination of the overall program's dynamics by highlighting the previously introduced user's perspective through a [UML](#page-136-2) sequence diagram (see Figure [4.5](#page-47-0) and Appendix [C.3\)](#page-131-0).

From the user's perspective, the application data model first must be constructed and initialized. The construction (see Figure [4.6\)](#page-48-0) consists of the loading of data files required in the subsequent simulation process. The data that is added to the model is optional but highly dependant on the simulation's purpose. Here it is defined which model segments are used as modules for the description of the geometry and the activity.

The initialization (see Figure [4.7\)](#page-49-0) of the model then links the previously loaded data together. This makes sure that every segment of the model knows which data has to be applied and how it is going to be used. After the linking of data has terminated, the simulation engine is instantiated and further connected to the data model. This accounts for possible errors that could occur, if the data was either incompletely or improperly linked together. Therefore, the initialization is also a data-verification and data-linkage verification mechanism.

Figure 4.5: Sequence diagram: User perspective. The given figure is a [UML](#page-136-2) sequence diagram representation of the user's perspective use-case as described in the previous Section [\(4.1.3.1\)](#page-44-0). Here, the function calls for the creation of data objects and the triggering of a simulation are outlined. It can be seen that in addition to ordinary function calls with a synchronous return of approval, there is an additional incidence of an asynchronous return event: This repetitive event call states the availability of simulation output data that can be collected by the user. The model construction (Figure [4.6\)](#page-48-0), initialization (Figure [4.7\)](#page-49-0), and destruction (Figure [4.8\)](#page-49-1) are out-sourced for better clarity and readability.

Hence a simulation can only occur if the initial data is complete and properly linked.

If the model was loaded and found acceptable by the initialization routines, the actual simulation can begin. By doing so the simulation engine is invoked to iterate over the given data through the use of the given algorithms in defined time steps (see Figure [4.5\)](#page-47-0). After each iteration output data is produced which can be further used as a base for the next iteration cycle and is returned to the

Figure 4.6: Sequence diagram reference: Model construction. The construction of the application's data model consists of the addition of data segments as of the geometry description, the activity description, reaction, as well as diffusion algorithms. The adding of such segments is optional and highly depends on the purpose of the simulation. At this stage the data is loaded but not further connected to each other. This connection process will be executed during data initialization (see Figure [4.7\)](#page-49-0).

user for subsequent analysis. To ensure that the running simulation does not block any user interactions, it is set up as a separate thread. Hence the user is informed of newly created output data via asynchronous calls from that thread. The output data itself is continuously saved to the connected permanent storage of the systems hardware.

Next either the user is able to stop the simulation manually or simply waits until the previously defined number of iterations have been processed and the simulation engine auto-terminates. With either suspension mode, the user is able to continue simulation since the output data can be reused as new input

Figure 4.7: Sequence diagram reference: Model initialization. After the application model data was loaded and set to the program's internal data sets it has to be connected to each other. This is done iteratively by searching the geometry descriptions for compartments that give further clues on what other definitions have to be linked.

Figure 4.8: Sequence diagram reference: Model destruction. If the currently loaded model is no longer needed, it has to be destroyed before other (new) models can be loaded. The process of destruction is highly dependant on the chronology of what object has to be deleted first. It is of major importance to remove the simulation engine instance first and then destroy the input description and the algorithms in this order. Doing so ensures that the data is removed from its linkage dependencies safely without creating broken connections that may trigger undesirable memory leaks.

data. Here no reinitialization is necessary because all data remains complete and linked until model destruction.

The destruction of data (see Figure [4.8\)](#page-49-1) has to be then performed to avoid possible memory leakage or the overcrowding of physical memory. The destruction is an ordered process of chronological data model segment instance removal. First the simulation engine is destroyed since the destruction will finally invalidate the given data thereby making the simulation engine obsolete. At this point the data is removed from top to bottom by first deleting the input descriptions and then also the loaded algorithms from memory. It should be noted that no data is finally lost at this point due to the automatic saving of output data by the simulation engine. If the user decides to remove such stored files as well, then the data is definitely lost and can only be reproduced by a repeated simulation.

4.2 Considerations and Definitions

As the general concept of the [4DiCeS](#page-134-1) framework was displayed already, this section is concerned with the design and definition of the interfaces. Details regarding how [4DiCeS](#page-134-1) offers its implemented modularity as combined with cross platform compatibility will then be further described in Section [4.3.](#page-79-0)

The design of the framework specifies that a plug-in module must be chosen and connected to the kernel for the system to run successfully. All interfaces are defined in a modular way in order to provide greater exchangeability of algorithms and to facilitate hassle-free software updates.

Currently there are four different interfaces to the [4DiCeS](#page-134-1) kernel (see Table [4.3\)](#page-43-0). They allow for connections of modules that handle diffusion and reaction algorithms, the parsing of description files, and describe the connection to user interfaces. The dynamics of the interfaces facilitate the integration of new methods and algorithms not yet implemented. Additionally it is thus more easily possible to exchange old modules for new versions. The potential for extending the interfaces to other programming languages as well as other applications is significant. The plug-ins could also integrate sockets to allow for network distributed computing or the remote diagnostics of running processes.

The following subsections characterize the current implementations for modeling of [3D](#page-134-2) geometry (Section [4.2.1\)](#page-51-0), for handling activity descriptions (Section

Figure 4.9: [4DiCeS](#page-134-1) modules. The kernel of [4DiCeS](#page-134-1) comprises four interfaces to modules that have to be engaged before a simulation is able to run. These ports allow for the integration of various reactions as well as diffusion algorithm modules, model (activity and geometry) parsers, and for user interactions.

[4.2.2\)](#page-60-0), and for utilizing algorithms (Section [4.2.3\)](#page-70-0). Thereafter a subsection describing the user interfaces will consider the momentary connections between [4DiCeS](#page-134-1) and the user (see Section [4.2.4\)](#page-78-0). The underlying plug-in design details are then discussed in Section [4.3.](#page-79-0)

4.2.1 Geometry Model

Before considering the actual geometry implementation of [4DiCeS](#page-134-1) the following section briefly introduces fundamental [3D](#page-134-2) geometry characterization descriptions. Actually, there exist several techniques to present [3D](#page-134-2) geometries. But common representation schemas for modeling of [3D](#page-134-2) objects can be divided into three main groups (see Figure [4.10\)](#page-52-0): 'wire-frame models', 'surface models', and 'body models' (Encarnação et al., [1997b\)](#page-146-0).

Wire-Frame Models: Among the [3D](#page-134-2) model types described above the wireframe model requires the least information for representing [3D](#page-134-2) bodies. At large, its structure elements are limited to outline elements as of 'straight edges', 'circular arcs', or 'spline curves'. These elements are not related to each other within this model type. Hence, an assignment of such elements to surfaces is not defined. Therefore, wire-frame models can be considered simple in that data on the surface and interiors are lost in the sense that the wire-frame model no longer describes that information. Hence they do not qualify for a representation of body models due to their incomplete and ambiguous presentation. In many cases algorithms on wire-frame models demand for an intensive interac-

Figure 4.10: Representation schemas for [3D](#page-134-2) models. Both the wire-frame and the surface models allow for only an visually-interactive user interpretation. They are not complete and distinct in their representation of [3D](#page-134-2) objects. On the contrary body models provide a comprehensive description of a [3D](#page-134-2) object, and can be interpreted algorithmically. There are many representations for body models. The most commonly used models are the cell model, [CSG](#page-134-10), [BRep](#page-134-11), and hybrids between the [CSG](#page-134-10) and [BRep](#page-134-11) representations. Adapted from Encarnação et al. [\(1997b\)](#page-146-0).

tion with the user, who then has to specify corresponding surfaces interactively.

Surface Models: Surface models comprise the creation of [3D](#page-134-2) objects where the facing and its attributes as for bend, torsion, and smoothness are of major importance. Therefore the main content of information lies in single face descriptions. Again there is no correspondence between such faces. In particular relationships to neighbors are not included to its data representation. Such a surface modeler is not qualified for presenting proper body models. A classification of one point in space according to an object cannot be dealt with. That is why it is hard to decide if a point lies inside or outside of a displayed [3D](#page-134-2) object. To come to such a decision, normal vectors have to be applied facing to, and therefore defining, the outside. Surface models without such orientated faces are neither distinct nor complete. Likewise surface models are better than wire-frame because more data is included. This data is however not a comprehensive representation of the information it is meant to embody.

Body Models: The term body models stands for many representation schemas (see Figure [4.10\)](#page-52-0). Body models form a complete description of [3D](#page-134-2) objects and can be interpreted automatically by a program. Due to complete storage of the body's [3D](#page-134-2) representation geometry, the question can be addressed algorithmically. As another positive feature the consistence of objects can be ensured with algorithmic manipulation, as the result of an operation (e.g. the combination of two objects) gives a valid object again. The following geometry considerations will deal with body models only.

As for the current implementation of [4DiCeS](#page-134-1), body models were applied as the geometrical description of a simulation space. Therefore the following Section [4.2.1.1](#page-53-0) will go further into illustrating the general simulation space definition in [4DiCeS](#page-134-1). The proximate Section [4.2.1.2](#page-55-0) will then provide details on the reference implementation of this body model, and how it is handled by the [4DiCeS](#page-134-1) kernel. The activity description (see Section [4.2.2\)](#page-60-0) then formally describes the interaction of both the currently supported geometry data structure and activity models in detail.

4.2.1.1 Simulation Space

As can be seen from Figure [4.10,](#page-52-0) there are several different approaches of modeling a [3D](#page-134-2) geometry. As the most commonly used representatives of such techniques are [BRep](#page-134-11), the [CSG](#page-134-10), and the cell model, they are going to be briefly described. A specialization of the cell model was used for the underlying simulation space geometry of [4DiCeS](#page-134-1) and is going to be laid open in detail with the upcoming section.

Boundary Representation: A body model can be described explicitly by its surface and an associated topological orientation. Because of this topological orientation, every point is well-defined with regard to whether it lies inside or outside of a plain. Boundary Representations ([BRep](#page-134-11)s) use this fact, and describe [3D](#page-134-2) objects by their surface. Polygon-oriented data-structures as edge-oriented, node-oriented, or winged-edge data-structures can be applied for the storage of [BRep](#page-134-11)s.

Constructive Solid Geometry: With Constructive Solid Geometry ([CSG](#page-134-10)), [3D](#page-134-2) bodies are described by trees of Boolean operators and primitives. These are also called [CSG](#page-134-10)-trees. Each primitive can then be defined by either [BRep](#page-134-11)s or half-space models. The display of [CSG](#page-134-10)-objects as trees is closely linked to the construction of the objects. The nodes of the tree represent regulatory Boolean

Figure 4.11: Model primitives. As for any body model primitives, the dimension parameters (here a cube with: xL, yL, and zL) as well as the position of the point of origin (x, y, and z) define a primitive unambiguously. Adapted from Encarnação et al. [\(1997b\)](#page-146-0).

set operations or transformations in space. The leaves refer to the primitives. A data-structure would make use of the tree-structure and the particular primitive representation used.

Cell Model: In a cell model, a [3D](#page-134-2) body is divided into a set of non-overlapping neighbor cells. Such cells have various forms, sizes, positions, and orientations. Therefore, the division of cells depends on a distinct number of cell types and simple operations for their assembly. The single cells can be arbitrary objects that are topologically equivalent to a sphere in [3D](#page-134-2) space. A [3D](#page-134-2) body is therefore composed of a set of half-disjunct cells. Data-structures highly depend on the cell type used. Often tree-structures come into use as for quad-trees and octrees.

A base body primitive is well-defined by fixing the dimension parameters of a [3D](#page-134-2) primitive (see Figure [4.11\)](#page-54-0). In a cell model, the base bodies can have various shapes: cylinders, cones, spheres, toroids, or cuboids. Dimension parameters are e.g. three-side lengths for cuboids or one length and a diameter for a cylinder. Each object has a point of origin to a local coordinate system. This local coordinate system cannot be altered, since it is \dot{a} priori fixed for every base body primitive.

In [4DiCeS](#page-134-1) special constraints have to be applied to a cell model for it to be suitable for representing a simulation space (see Table [4.5\)](#page-55-1). First of all it is of major importance that the complete geometrical space of simulation is described in full. This means that there must exist an uninterrupted representation of both a global reaction and diffusion area. This area may be segmented into closed

Table 4.5: Requirements to the simulation space. In addition to the definition rules of a used body model there exist four further specifications for a geometry representation in [4DiCeS](#page-134-1).

compartments that may be partitioned again. Such a nested structure of [3D](#page-134-2) region descriptions ensures that a model can be simulated. Another constraint is that compartments must not overlap. Furthermore all defined structures must provide slots to link in other parts of the application data model. A simulation space is complete if all given constraints are obeyed. During the creation and the initialization, the application data model relies on this additional information.

Actually, all body models could be applied to [4DiCeS](#page-134-1) if they obey the given requirements. But as a first reference implementation to such a model, a specialized form of the cell model was used that is specified in the subsequent section.

4.2.1.2 Geometrical Data Structure

The cell model described earlier (see Section [4.2.1.1\)](#page-53-0) was used for a reference implementation of a geometry model in [4DiCeS](#page-134-1). A cuboid was used to define the cell type. Thus the [3D](#page-134-2) simulation space is divided into a [3D](#page-134-2) grid of equallysized sub-volumes. In the literature such a structure is also referred to as a grid of voxels, which is an aggregation of the terms "volume" and "pixel" (Encarnação et al., [1997a\)](#page-146-1). Often voxels are also called Volume Elements ([VE](#page-136-3)s), in accordance to a grided volume of distinct elements. A generalized scheme

Figure 4.12: Simulation voxel space. A biological cell is subdivided into discrete cubes, called [VE](#page-136-3)s. These [VE](#page-136-3)s can harbor multiple particle species and may have a membrane as their surface area. To allow for various geometries, the [VE](#page-136-3)s are designed to be freely scalable in size.

of such a diced model can be seen in Figure [4.12](#page-56-0) that shows a schematized biological cell following this [3D](#page-134-2) geometry description.

The [VE](#page-136-3)s used here are able to form all sorts of complex compartments. This disposition of distinct but regular cubes has enormous benefits to further modeling and simulation. The size of the [VE](#page-136-3)s can be specified during initialization that helps to keep computational time low. Also such a geometry can be easily persisted in a simple [2D](#page-134-3) data-structure.

The [VE](#page-136-3) defines the smallest geometrical structure within the simulation space. It can also specify boundary conditions that can then be used as membranes for particle diffusion. Every side of the cuboid voxel can have a membrane area. Thus a [VE](#page-136-3) can be considered to be an open system if no membranes are set. Conversely the system is closed and reflecting if a [VE](#page-136-3) is fully covered by membrane. The total coverage of a [VE](#page-136-3) needs six membrane faces. Hence two neighboring [VE](#page-136-3)s with adjacent membranes are analogous to a cellular lipid bi-layered membrane. The membranes function either as barriers for definable particle species or as binding areas for transmembranous proteins (see Figure [4.13\)](#page-57-0). Opposite marginal [VE](#page-136-3)s on the grid, if not divided by membranes, allow the movement of particles from one side of the grid to the other (i.e. exiting the far right and reentering to the left of the grid again).

Every cube is able to hold particle concentrations for each particle species applied to the model. Unbound particles are allowed to diffuse freely in all

Figure 4.13: Boundary conditions. Every [VE](#page-136-3) may be covered by membranes on every one of its six faces. Two neighboring [VE](#page-136-3)s and membranes make a lipid bi-layer. Therefore it is possible to have both free and membrane-bound particles. The permeability of particles through the membranes may be individually adjusted.

directions. Membrane bound particles can only diffuse laterally within the membrane (see Figure [4.13\)](#page-57-0). A voxel is composed of six direct faces, 12 edges and eight corners. While each face has one direct neighbor, the edges have three and the corners eleven. The diffusion from one [VE](#page-136-3) to another is handled such that the probability of leaving a [VE](#page-136-3) is higher with the face neighbors than with edge neighbors. The corners could be neglected since statistically corner diffusion is constantly small and therefore often dispensable.

The allocation of particles is uniformly distributed within a [VE](#page-136-3) at initialization. This means that it is not possible to further specify particle positions within a voxel. However the number of particles per species and the membrane coverage can be set by the user.

To store and manage the [3D](#page-134-2) voxel geometry, a common [2D](#page-134-3) graphics format was applied. The Tagged Image File Format $(TIFF)^1$ $(TIFF)^1$ $(TIFF)^1$ $(TIFF)^1$ was chosen due to its ability to handle different color spaces, to utilize powerful as well as lossless data compression algorithms, and to allow for multi-page-[TIFF](#page-136-4)s. [TIFF](#page-136-4) is designed to work well in many viewing applications, such as the Internet, so it is fully streamable with a progressive display option. [TIFF](#page-136-4) is robust providing both full file integrity checking, as well as simple detection of common transmission errors [\(Miano,](#page-152-0) [1999\)](#page-152-0). [TIFF](#page-136-4) can store gamma and chromaticity data for improved color

¹[TIFF](#page-136-4) Specification: <http://partners.adobe.com/public/developer/en/tiff/TIFF6.pdf>

Figure 4.14: Placement of relevant geometrical information. [\(a\)](#page-58-0) The number of particles per voxel is represented as a pixel in an image. [\(b\)](#page-58-1) Represents the membrane-setting for all cubes of the grid. Here, membranes are binary-encoded as can be seen at Table [4.6.](#page-59-0) (a) + [\(b\)](#page-58-1) All the $x-y$ slices (the z stack) are spread in the picture's x direction. [\(a\)](#page-58-0) Thereby, every single particle species can be repre-sented in its [3D](#page-134-2) geometry in the [TIFF](#page-136-4)'s y direction. Each new step in time produces a new particle [TIFF](#page-136-4) page.

matching on heterogeneous platforms. Last but not least, [TIFF](#page-136-4) does not require licensing and can therefore be freely used [\(Murray and Ryper,](#page-153-2) [1994\)](#page-153-2).

The challenge of converting the [3D](#page-134-2) data into a [2D](#page-134-3) picture format was resolved by displaying a [3D](#page-134-2) stack of [VE](#page-136-3)s in two dimensions (see Figures [4.14](#page-58-2) and [4.15\)](#page-59-1). The grid of [VE](#page-136-3)s has its three spatial dimensions $(x, y, \text{ and } z)$. Every VE can be directly addressed via these unique space coordinates in [3D](#page-134-2). It is adequate to speak of $x-y$ -layers in the grid's z axis. These layers are organized sequentially in one row of a [TIFF](#page-136-4) image. The [VE](#page-136-3) information of the number of particles is stored as a 64bit integer pixel. Every modeled particle species is given by its own row of x–y-layers (see Figure [4.14\(a\)\)](#page-58-0). The result is a static [3D](#page-134-2) image of the [VE](#page-136-3)s' particles of the grid. Every [VE](#page-136-3)-grid-snapshot in time is stored as a separate [TIFF](#page-136-4)-page (see Figure [4.15\)](#page-59-1). The reusability of the output is enhanced, because a simulation can always be restarted with any of the previously produced [TIFF](#page-136-4) as input data again.

Similar as to the particle storage in [TIFF](#page-136-4) images, the structural information of membranes is also captured in a [TIFF](#page-136-4). Again, the $x-y$ -layers in the grid's z axis are saved. In contrast to the before mentioned, only the six faces are stored as a sequence of 'set' and 'unset' bits (see Figure [4.14\(b\)](#page-58-1) and Table [4.6\)](#page-59-0).

The [TIFF](#page-136-4)-file can also store additional textual information. This capability is used to provide auxiliary data on the geometry of the [3D](#page-134-2) grid representa-

Table 4.6: Membrane binary codes. Stored to a [TIFF](#page-136-4)-page, membranes are binary-encoded. Each of the six cube faces has a unique identification bit. Combined together into a pixel value, this encoding represent the coverage of a [VE](#page-136-3) by membrane. A fully covered voxel with membrane would therefore have the value 0x3F. A plain voxel without any membranes is represented by a value of zero $(0x00)$.

Figure 4.15: Transformation of [3D](#page-134-2) into [2D](#page-134-3) data. The picture shows the conversion of [VE](#page-136-3)s into two [2D](#page-134-3) [TIFF](#page-136-4)-pages. Here, the underlying information of all particles and the membrane geometry is stored as pixel values of an image file. The different colors yellow $\{x_0, y_0, z_0\}$, red ${x_1, y_0, z_0}$, brown ${x_0, y_1, z_0}$, blue ${x_1, y_1, z_0}$, cyan ${x_0, y_0, z_1}$, purple $\{x_1, y_0, z_1\}$, orange $\{x_0, y_1, z_1\}$, and green $\{x_1, y_1, z_1\}$ denote specific [VE](#page-136-3)s at their corresponding positions in X, Y , and Z respectively. Each voxel V includes a constant set of particle species S. After every iteration cycle t of time T the same structure is stored as a separate page to the [TIFF](#page-136-4) image.

tion. The underlying activity description, geometrical details, links to applied algorithms, Globally Unique IDentifiers ([GUID](#page-135-2)s), and version numbers can be stored as "tags" to the [TIFF](#page-136-4) image files. The grid representation is stored as three integer values that mark the length of each grid axis $(x, y, \text{ and } z)$. The

underlying cell model is tagged as the file name of the modeling file relative to the [TIFF](#page-136-4) directory position.

[TIFF](#page-136-4) files can be loaded with a great variety of existing image processing applications. Furthermore the implementation of an analysis tool for this kind of file format is not difficult. Several image processing libraries in various pro-gramming languages offer accordant [TIFF](#page-136-4) support^{[2](#page-60-1)}. The compression rate of [TIFF](#page-136-4) images allows for huge amounts of time series pictures to be created with relatively low storage requirements. As compared with flat text files of the same information, the [TIFF](#page-136-4) used here achieves a compression rate of approximately 1/200. This partially addresses the commonly encountered problem of storing large amounts of data.

4.2.2 Activity Description

As seen in the previous section the representation of [4DiCeS](#page-134-1) geometry was achieved by analyzing as well as utilizing existing [3D](#page-134-2) modeling methodologies. The computational modeling of activity descriptions requires a similar approach. Again model description files have to be utilized for an adequate persistence mechanism to the underlying system. To be useful as a formal characterization of biological systems understanding, mathematical models are put into a format that may be transferred successfully between different software tools [\(Bower and Bolouri,](#page-144-0) [2004\)](#page-144-0). Such a format then addresses a number of problems facing simulation and modeling (three of such major problems are stated in Table [4.7\)](#page-61-0).

One problem occurs when users have to handle multiple resources from various software tools during a project. Different tools have different capabilities and strengths. Working with different tools today mainly requires the re-encoding of models in each tool that is a very error-prone and time-consuming process. Another problem arises with the need for obtaining the model definitions in electronic form. This especially can be found with journal publications where authors use different modeling environments and model representation languages. Such definitions are often not straightforward to examine, compare, test, or reuse. A researcher typically must manually transcribe the model into his software tool's appropriate format. Last but not least, a problem exists if a simulation tool exceeds support. Models developed on such systems may become obsolete and therefore are often abandoned. The development of new

²Free graphics libraries: <http://www.thefreecountry.com/sourcecode/graphics.shtml>

Table 4.7: Problems facing simulation and modeling. This table states three problems facing simulation and modeling.

tools will only intensify these problems unless the issue of standard formats or clear and automatic translation between them is addressed [\(Hucka et al.,](#page-149-0) [2004\)](#page-149-0).

The development of standard modeling languages are being conducted in effort to address such problems. Standardized formats can then be used to communicate and exchange models between different software tools. By supporting such general formats as input and output formats, different software systems can all operate on the identical representation of a model. This provides a common starting point for simulations as well as analysis and eliminates the possibility of translation errors.

Model-data storage is currently accomplished either with proprietary data formats or by derivatives of markup languages such as [XML](#page-136-5) (see Table [4.8.](#page-62-0) Recently an attempt was made to use [UML](#page-136-2) as a biological model description language [\(Webb and White,](#page-158-0) [2005\)](#page-158-0). Most of the biological model description formats are freely available standards. While numerous language definition draft documents are still copyrighted, a non-commercial use is usually accepted and even often designated.

The popularity of [XML](#page-136-5) in the area of computational biology has grown in the past decade. [XML](#page-136-5) provides the capability of representing data in a comprehensive and standardized data structure. The structure of [XML](#page-136-5) documents defined using a [XML](#page-136-5) Document Type Definition ([DTD](#page-135-3)) however is limited to representing data in a hierarchical tree fashion. Such an approach imposes severe limitations on both the structure, as well as the ability to validate a document. The utilization of the World Wide Web ([WWW](#page-136-6)) Consortium [XML](#page-136-5) "schema" approach to document definition thus overcomes many of such limitations.

Table 4.8: Input description file standards. The description file formats displayed in this table are either cellular modeling languages or their specific helper-languages. They can be roughly described as proprietary and markup languages.

Many data formats have been proposed for representing models in the life-sciences including the BIOpolymer Markup Language ([BIOML](#page-134-14))^{[3](#page-62-1)} [\(Fenyo,](#page-146-2) [1999\)](#page-146-2), the Chemical Markup Language $(CML)^4$ $(CML)^4$ $(CML)^4$ $(CML)^4$, the [EML](#page-135-6) $((Sakurada et al., 2002)),$ $((Sakurada et al., 2002)),$ $((Sakurada et al., 2002)),$ $((Sakurada et al., 2002)),$ $((Sakurada et al., 2002)),$ the MicroArray Gene Expression Markup Language ([MAGE-ML](#page-135-9))^{[5](#page-62-3)} [\(Spellman](#page-157-0) [et al.,](#page-157-0) [2002\)](#page-157-0), the Proteomics Experiment Markup Language ([PEML](#page-136-8)) [6](#page-62-4) [\(Taylor](#page-157-1) [et al.,](#page-157-1) [2003\)](#page-157-1), the Protein Markup Language ([ProML](#page-136-9)) [7](#page-62-5) [\(Hanisch et al.,](#page-148-0) [2002\)](#page-148-0), and the Proteomics Standards Initiative's Molecular Interaction ([PSI-MI](#page-136-10)) [8](#page-62-6) lan-

³[BIOML](#page-134-14): <http://xml.coverpages.org/bioml.html>

⁴[CML](#page-134-15): <http://www.XML-CML.org/>

 5 [MAGE-ML](#page-135-9): <http://www.mged.org/MAGE/>

⁶[PEML](#page-136-8): <http://www.ccbm.jhu.edu/>

⁷[ProML](#page-136-9): <http://www.scai.fraunhofer.de/>

⁸[PSI-MI](#page-136-10): <http://psidev.sourceforge.net/>

guage [\(Hermjakob et al.,](#page-149-1) [2004\)](#page-149-1). There are also a few general-purpose modeling definition formats that can be found for biological model documentations. Two very popular examples are the Petri-Net Markup Language ([PNML](#page-136-11)) [9](#page-63-0) and the Resource Description Framework ([RDF](#page-136-12))^{[10](#page-63-1)}.

There are only two [XML](#page-136-5)-based formats that are designed to unambiguously specify, store, and exchange biological models in a form that is both computer and human readable – [CellML](#page-134-7) [\(Crampin et al.,](#page-145-0) [2004\)](#page-145-0), and [SBML](#page-136-1) [\(Hucka et al.,](#page-149-0) [2004\)](#page-149-0). The Dynamic Signaling Maps Language ([DSML](#page-135-5)) that could be used in the same way is not [XML](#page-136-5)-based but a proprietary flat-file language.

Both [SBML](#page-136-1) and [CellML](#page-134-7) support the Mathematical Markup Language ([MathML](#page-135-10)) [\(Ausbrooks et al.,](#page-142-0) [2003\)](#page-142-0). [MathML](#page-135-10) is an [XML](#page-136-5) application for describing mathematical notation and capturing both its structure and content. The goal of [MathML](#page-135-10) is to enable mathematics to be served, received, and processed, just as the HyperText Markup Language ([HTML](#page-135-11)) has enabled this functionality for text on the [WWW](#page-136-6).

[4DiCeS](#page-134-1) has currently reference implementations for two model description formats. The first is a proprietary description language designed during the development of [4DiCeS](#page-134-1) and can be considered the systems natively supported activity description format. This flat-file format called [4DiCeS](#page-134-1) Model Description ([FMD](#page-135-12)) is further discussed in Section [4.2.2.1.](#page-63-2) The second file format supported at present is [SBML](#page-136-1) (see Section [4.2.2.2\)](#page-65-0). The activity considerations will be closed by a closer look at the formal integration of these two description file formats into the [4DiCeS](#page-134-1) framework (see Section [4.2.2.3\)](#page-66-0).

4.2.2.1 [4DiCeS](#page-134-1) Model Description ([FMD](#page-135-12)) File Format

During early development of the [4DiCeS](#page-134-1) system, a custom-made modeling file format was used together with the kernel-parsing interface. This format allows for the design of models consisting of both data for the reaction as well as diffusion processes. The format is a flat-file that may integrate [TIFF](#page-136-4) geometry files as discussed in Section [4.2.1.2.](#page-55-0) In the following section, the four-tuple grammar $G = (T, N, P, S)$ of this flat-file modeling language is defined using the Backus–Naur Form ([BNF](#page-134-16)) (see Appendix [B\)](#page-126-0):

 9 [PNML](#page-136-11): <http://www.informatik.hu-berlin.de/top/pnml/> 10 [RDF](#page-136-12): <http://www.w3.org/RDF/>

$$
T = \{ 0, \dots 9, a, \dots z, A, \dots Z, +, >, |, @, *, \sim, -, \dots \} \tag{4.1}
$$

$$
N = \{ \langle outerproduct \rangle, \langle innerproduct \rangle, \langle educt \rangle, \langle particle \rangle, \\ \langle species \rangle, \langle volumeelement \rangle, \langle reaction \rangle, \langle rate \rangle, \\ \langle product \rangle, \langle letter \rangle, \langle digit \rangle, \langle amount \rangle, \langle sign \rangle, \\ \langle name \rangle, \langle membrane \rangle \} \}
$$
(4.2)

$$
S = \langle volume element \rangle, \tag{4.3}
$$

where T denotes the set of allowed terminal symbols. N specifies the set of all possible non-terminal symbols, and S defines the starting symbol as further defined in [4.10.](#page-64-0) The production system P is further specified by:

$$
\langle letter \rangle ::= 'a'..'z'|'A'..'Z'
$$
\n(4.4)

$$
\langle digit \rangle ::= '0'..'9'
$$
\n
$$
(4.5)
$$

$$
\langle percent \rangle ::= ['0']'.'\{\langle digit \rangle\}|'1'
$$
\n(4.6)

$$
\langle sign \rangle ::= ' *' |' \sim' |' -' |'.
$$
\n
$$
(4.7)
$$

$$
\langle name \rangle \ ::= \ ['''] \langle letter \rangle \{ \langle letter \rangle |'']
$$

$$
\langle digit \rangle | \langle sign \rangle \}
$$
 (4.8)

(4.9)

for base primitives, and:

$$
\langle volume element \rangle ::= \langle reaction \rangle \{ \langle 'reaction \rangle \} \tag{4.10}
$$

$$
\langle reaction\rangle ::= \langle rate\rangle \langle eluct\rangle \langle product\rangle \tag{4.11}
$$

$$
\langle rate \rangle ::= '(' \langle percent \rangle')'
$$
\n(4.12)

$$
\langle educt \rangle ::= \langle species \rangle \{ '+' \langle species \rangle \} \tag{4.13}
$$

$$
\langle species \rangle ::= \langle amount \rangle \langle particle \rangle \tag{4.14}
$$

$$
\langle particle \rangle ::= \langle name \rangle [\langle membrane \rangle] \tag{4.15}
$$

$$
\langle amount \rangle ::= \{ \langle digit \rangle \} \tag{4.16}
$$

$$
\langle membrane \rangle ::= '@'\langle name \rangle \tag{4.17}
$$

$$
\langle product \rangle ::= \langle innerproduct \rangle [\langle outerproduct \rangle]]
$$

$$
\langle outerproduct \rangle [\langle innerproduct \rangle]
$$
(4.18)

$$
\langle innerproduct \rangle ::= ' >' \langle educt \rangle \tag{4.19}
$$

$$
\langle outerproduct \rangle ::= \langle \langle \langle educt \rangle \rangle \tag{4.20}
$$

where a reaction (4.11) is defined by a set of educts (4.13) , products (4.18) , and a reaction rate [\(4.12\)](#page-64-0). A product can further be divided into outer [\(4.20\)](#page-64-0) and inner products [\(4.19\)](#page-64-0) that define a reaction output to be either in the inside or the outside of the current reaction compartment. Both educts and products are specified by particle species [\(4.14\)](#page-64-0) and their momentary concentration [\(4.16\)](#page-64-0). Particle species can be tagged as being membrane-bound by concatenating a membrane (4.17) , by the use of an ' $@$ ' symbol, to them.

This description specification is of course by far not complete but applicable for many usage scenarios of [4DiCeS](#page-134-1) in its present state. [FMD](#page-135-12) was especially useful for testing as well as verification purposes during implementation and can be considered to be the native description file format of activities in [4DiCeS](#page-134-1).

4.2.2.2 Systems Biology Markup Language ([SBML](#page-136-1))

The Caltech ERATO Kitano systems biology project has developed the Systems Biology Markup Language ([SBML](#page-136-1)) for the representation and modeling of information components in cellular systems. [SBML](#page-136-1) serves as an attempt to specify a common, model-based description language for systems biology simulation software [\(Hucka et al.,](#page-149-0) [2004\)](#page-149-0). At the moment the world-wide [SBML](#page-136-1) community counts over 80 software tools supporting this language in one or the other form ([SBML](#page-136-1) Software Guide^{[11](#page-65-1)}). The SBML representation language is organized around five categories of information: model, compartment, geometry, species, and reaction. The intent is rather to cover the range of data structures needed by the collection of all of the simulators examined so far [\(Finney and Hucka,](#page-147-0) [2003\)](#page-147-0). A major drawback is that [SBML](#page-136-1) does not currently have the means to represent Partial Differential Equation ([PDE](#page-136-13))-level models or diffusion terms [\(Hucka et al.,](#page-150-0) [2003\)](#page-150-0). Cellular geometry was also rather neglected. To surround such deficiencies, especially the lack of support for geometry definitions, great efforts have to be made on part of [4DiCeS](#page-134-1) (see the following section [4.2.2.3\)](#page-66-0). How other application are dealing with such is further shown in Section [5.4.](#page-106-0)

¹¹[SBML](#page-136-1) Software Guide: [http://sbml.org/SBML](http://sbml.org/SBML_Software_Guide/) Software Guide/

In the following the import to as well as the export from [4DiCeS](#page-134-1) of the previously presented two file formats [FMD](#page-135-12) and [SBML](#page-136-1) is going to be presented in more detail.

4.2.2.3 Import and Export

As for applying a model to the [4DiCeS](#page-134-1) suite, there are three specific user scenarios (see Figures [4.16\)](#page-67-0). Depending on what type of data is to be included to the application data model, there is one use-case for loading an activity description without any further geometry definitions, then a use-case with rudimentary geometry to allow for concurrent algorithms, and lastly a use-case describing the full interaction with both activity and geometry description in [3D](#page-134-2).

Plain Activity Model: If a user simply applies a given activity model file to [4DiCeS](#page-134-1), then there is no further need for any geometry information. The only thing that has to be specified in addition to the activity data is what kind of reaction algorithm should be applied to the simulation process. Within such a scenario, there is no algorithmic concurrency allowed, because of missing segmentations that would be defined by a geometry description. The use-case can be seen in Figure [4.16\(a\).](#page-67-1) Such a simulation can be referred to as a "single segment simulation".

Rudimentary Geometry: If compartments with concurrently running algorithms are required, but no further [3D](#page-134-2) geometry is needed, then a limited form of the geometry description must be applied in addition to the activity description. Only defined compartments are mapped to segments of the geometry to allow for setting individual algorithms for each compartment. The use-case is visually described by Figure [4.16\(b\).](#page-67-2) This type of simulation is called here a "[2D](#page-134-3) multiple segments simulation".

Full Input Description: When a trackable [3D](#page-134-2) geometry is required in addition to the activity and concurrency of algorithms, then the user has to input two complete description files of both geometry and activity. The use-case for this type of model can be found in Figure [4.16\(c\).](#page-67-3) This model type leads to a full "[3D](#page-134-2) multiple segments simulation".

All three use-cases have an activity description and the application data model in common. While use-case Figure $4.16(a)$ requires only a separate setting of the applied reaction algorithm to establish the simulation engine, the other two use-cases (Figures [4.16\(b\)](#page-67-2) and [4.16\(c\)\)](#page-67-3) require both a geometry description and a set of diffusion algorithms. With Figure [4.16\(b\)](#page-67-2) the application data

(a) Single Segment Model

Figure 4.16: Use-cases: Modeling. Depending on what type of data is applied to [4DiCeS](#page-134-1), there are three different use-case scenarios. Figure [\(a\)](#page-67-1) shows the input of an activity description without any further geometry definitions. Figure [\(b\)](#page-67-2) displays a use-case with limited geometry to allow for concurrent algorithms, and Figure [\(c\)](#page-67-3) describes the full interaction with both activity and geometry description in [3D](#page-134-2).

model will attempt to set the needed connections itself, since no further geometry is needed. With Figure [4.16\(c\)](#page-67-3) the user must explicitly set the geometry description so that a simulation engine can be established.

A [UML](#page-136-2) sequence diagram (see Figure [4.17\)](#page-68-0) displays the behavioral specification in more detail. The conditions are highlighted that produce an alternate behavior of the application data model during initialization. The three conditions are summarized in Table [4.9.](#page-69-0) It should be noted that if any of the three alternatives fails to connect to the remaining data sets, then the simulation engine

Figure 4.17: Sequence diagram: Model alternatives. This figure displays the conditions that lead to the three alternative model instances. On initialization, the kernel checks for the existence of a geometry description. If and only if such a description file exists the application data model instructs the geometry description to connect to the remaining data set. Otherwise, the kernel will further check for an existing set of diffusion algorithms. If and only if that set exists a simple geometry description is created, which again connects to the remaining data set. Only if both test are negative the kernel instructs the activity description to connect to the corresponding reaction algorithm. In either case, if and only if a final connection to the remaining data sets is successful, then a simulation engine is instantiated.

will not be produced and a simulation of the model will not be feasible. The

Condition	Resulting $Use-Case$	Description
$\Phi \neq \emptyset$	4.16(c)	3D Multiple Segments Model
$\Phi = \emptyset \land \Theta \neq \emptyset$	4.16(b)	2D Multiple Segments Model
$\Phi = \emptyset \land \Theta = \emptyset$	4.16(a)	Single Segment Model

Table 4.9: Model alternative conditions. This table summarizes the conditions that results in the three alternative use-case scenarios. Here Φ is a set of segments and Θ defines a set of diffusion algorithms. These sets can either be empty or filled with their corresponding elements. The case $\Phi \neq \emptyset \wedge \Theta = \emptyset$ results into an error event, the initialization process is interrupted, and the application data model remains in its uninitialized state.

application data model thus stays in its uninitialized state until the missing data is added and an additional initialization call is invoked.

As already mentioned in the preceding section the [SBML](#page-136-1) is missing an advanced functionally in supporting geometry. Therefore this model description format is handled somewhat different. The single segment model as well as the [2D](#page-134-3) multiple segments model, there initialization routines, and the actual simulation work very similar as described before. But in contrast to the [3D](#page-134-2)-[FMD](#page-135-12) handling the [3D](#page-134-2) multiple segments model is currently not facilitated. Also concurrent algorithms cannot be inserted by [SBML](#page-136-1) into [4DiCeS](#page-134-1), since this file-format is missing such a concept completely. Anyhow, at least with the later on described [GUI](#page-135-0) implementation for [4DiCeS](#page-134-1), the user is able to connect different algorithms to the system manually. This change of course cannot be exported to [SBML](#page-136-1) again, which occasionally results in a loss of persistency of the data model.

Noteworthily, the application data model offers the entire bandwidth of "getter" and "setter" functions. This is therefore the place where the loading, saving, and validation of models occur. Data sets, parameters, boundary information, and the simulation space can all be manipulated through this one object instance. The application data model actually offers the most direct and powerful set of methods to the kernel. With this the security of the [4DiCeS](#page-134-1) system becomes a concern, since the system is not always able to safely process erroneous model data through the parsing interface. Fall-back mechanisms

within the system will therefore run a controlled model rollback if syntactical errors are detected during parsing or during initialization.

The now following section will deal with the algorithmic back-end of the [4DiCeS](#page-134-1) framework. The modeling part of the system's description will be hence left and turned towards the depiction of simulation mechanisms.

4.2.3 Algorithm Handling

Both ports to the diffusion and reaction modules are actually designed as interfaces within interfaces. The reason for this architecture is the algorithms' handles that are needed to perform further tasks. There are algorithms that focus only on single compartments, whilst others need the entire simulation space to work on. Thus there is a low-level interface allowing for full access to the simulation space and its particles. In case of a compartment algorithm, there is a second high-level interface that handles the traversing of the nested segments. Therefore, it is possible to integrate various different algorithms to the [4DiCeS](#page-134-1) kernel without code modification. The high-level interface then allows for simulating a model with different algorithms in concurrency.

The following Section [4.2.3.1](#page-70-1) defines both the high and the low-level interfaces. Two sections will thereafter focus on the connection of diffusion (Section [4.2.3.2\)](#page-73-0) and reaction (Section [4.2.3.3\)](#page-74-0) algorithms to the [4DiCeS](#page-134-1) kernel in greater detail.

4.2.3.1 Algorithm Interface Design

The interface-within-interface design introduced earlier is displayed in Figure [4.18](#page-71-0) by using a [UML](#page-136-2) component diagram (see Appendix [C.4\)](#page-132-0). It can be seen here that there is a base interface providing a set of low-level methods to the application data model. This set of low-level commands actually maps the application data model with all its data sets very closely. However the directness of control is a trade-off against convenience functions that the high-level interface then provides. As can be seen, the high-level interface connects to the low-level interface using its direct functionality. Even though not all methods are passed on from the low-level to the high-level command set, the high-level functions are enriched by methods that allow for the traversal of given model data geometries.

The interfaces are connected through inheritance, where a general low-level interface description feeds all other application data model interfaces (see Figure

4 Definitions and Implementation

Figure 4.18: Component diagram: Interface within interface. The given figure shows an [UML](#page-136-2) component diagram on the interface architecture for the [4DiCeS](#page-134-1) application data model. While both the geometry as well as the activity description sets are linked directly to the low-level interface, the algorithm sets for either reactions or diffusion also have the choice to connect to a high-level command layer. Here, the high-level interface is linked with the low-level commands. It provides functions for the automatic traversing of the input descriptions. Therefore, the higher-level interfaces allow for the user to work on segments and thus enable the possibility for algorithms to run concurrently.

[4.19\)](#page-72-0). Only the algorithm interfaces present a high-level interface to their lowlevel counterparts. As will be later seen, the user interface description will also be based upon this interface inheritance connection and doing so ensures that no data redundancies occur.

A realization of either reaction or diffusion algorithm modules can be implemented in a straight forward fashion. First a module must inform the kernel about its presence and nature. Next the user is able to access global information regarding the algorithm, particularly the algorithm's name, a brief description, a list of references, as well as the names of both the algorithm inventor, and the author of the concrete implementation. A very important piece of information that is also directly used by the application is the type of algorithm (i.e. if it is stochastic or deterministic). Depending on this information the application data model can ensure that parameters (e.g. rate constants) are automatically transformed into the underlying type of algorithm. If this first information ini-

Figure 4.19: Static diagram: Interface inheritance. As can be seen in this static [UML](#page-136-0) diagram, the interfaces are linked to each other through a chain of inheritance. On the low-level side (dash-doted rectangle), all interfaces (including the interfaces for the input descriptions) are specializations to a global low-level interface. Only the reaction and the diffusion algorithm interfaces provide a high-level interface. These are in turn specializations of their low-level counterparts. The application data model is linked through a composition to the general low-level interface, but has aggregation linkage to concrete realizations of algorithms.

tialization has finished, an algorithm iteration cycle can be induced by calling the methods 'react()' or 'diffuse()' respectively. As for the current implementation of [4DiCeS](#page-134-0), each iteration cycle will activate both present reaction and diffusion algorithms for every segment. This may lead to the problem that alternatively switching between the two algorithm forms is not desirable. In some cases this will also result in faulty output. A solution to this scenario is a coupled algorithm allowing for reaction and diffusion of particles at the same time. This, of course, could easily be done by connecting both interfaces to the same algorithm that is capable of processing both simultaneously.

The next two sections discuss applicable algorithms to the interfaces. While the concrete implementations of such algorithms will then be introduced in Section [4.3](#page-79-0) and discussed in greater detail within Appendix [A.](#page-118-0)

4.2.3.2 Diffusion

The [4DiCeS](#page-134-0) diffusion interfaces are able to handle various different algorithms. To achieve such a behavior, the interface depends on the separation into high and low-level interfaces as described earlier.

Table 4.10: Diffusion methods in biology. This table lists methods that are used for simulating spatial movement. The type column determines if a simulation is of stochastic or deterministic nature. The scale column distinguishes microscopic, macroscopic, as well as mesoscopic schemes. The primitive column differentiates between data primitives that are used in the simulation. The time column depicts between discrete events, discrete time-steps, and the results of a numeric solution. The non-spatial Gillespie Methods and [ODE](#page-136-1)s were included for comparison. Adapted from [Takahashi et al.](#page-157-1) [\(2005\)](#page-157-1).

In the case of random-walk simulations, where every particle is treated separately, the global grid handle or the global particle handle is sufficient. The segments' interface would however not be very useful in such a scenario. Ratewalk diffusions, where only numbers of particles or concentrations are shifted from one segment to another, can take advantage of this. Other deterministic methods could make use of the interface of segments also. With every iteration calculated results are passed back to the kernel through the interfaces, which must then take care of the reallocation of particles to their compartments.

The input and the output of the number of particles per species or their respective concentrations then highly depend on the interface used. For a global handling the number of particles (or concentrations) are returned for every explicitly-specified segment and particle species. In case of the segments' interface, the compartments' data is pushed to the algorithms and will be traversed by the simulation space itself. With global particle access, all particles of every particle species are delivered and retrieved by the interface at once. The reintegration of particles into the simulation space and its nested segments is handled internally.

Diffusion is normally referred to as Brownian motion [\(Brown,](#page-144-1) [1828\)](#page-144-1) of either uncharged or charged particles in solution. The first quantitative description of such a process was Fick's Law of Diffusion [\(Fick,](#page-146-1) [1855\)](#page-146-1). Later, an explanation for Fick's law was discovered [\(Einstein,](#page-146-2) [1905;](#page-146-2) [von Smoluchowski,](#page-158-2) [1906\)](#page-158-2). Today, there are several methods to model diffusion with computational assistance (see Table [4.10\)](#page-73-0). Prominent examples are deterministic approaches as for [ODE](#page-136-1), [PDE](#page-136-2), as well as Molecular Dynamics ([MD](#page-135-1)). Next to them a considerable number of stochastic methods exist such as the Brownian Dynamics ([BD](#page-134-1)), Gillespie derivatives, some [CA](#page-134-2) variations, and others. A stochastic random-walk simulation was chosen as a reference implementation for the [4DiCeS](#page-134-0) application. A brief description of the algorithm can be found within Appendix [A.2.](#page-123-0)

4.2.3.3 Reactions

Similarly to the integration of diffusion algorithms reaction interfaces are able to handle various different algorithms in [4DiCeS](#page-134-0). The interface definition again depends on the splitting of command sets into high and low-level interfaces functions. But in contrast to diffusion, the reaction interface must provide several specific methods that allow for the connection of reaction algorithm modules. Given reaction equations from an activity description can always be used, independent of the underlying algorithm.

The reaction equations contain information about the reaction educts, the resulting products, and a variable that specify the kinetics of a reaction. The products occur as a result of a reaction within a segment, or they can be passed through membranes. Admittedly only elementary chemical reactions with educts up to and including the second order are allowed with [4DiCeS](#page-134-0), because collisions of three or more molecules at an instant are considered to be fairly improbable [\(Gillespie,](#page-148-0) [1976,](#page-148-0) [1977\)](#page-148-1). Therefore the system needs to be modeled by using only three types of elementary reactions: first order, second order, and the homodimer formation [\(Gillespie,](#page-148-2) [1996\)](#page-148-2). Hence only reaction equations that match the following three patterns will be allowed to pass through the reaction interface:

$$
A \longrightarrow \dots \tag{4.21}
$$

$$
A + B \longrightarrow \dots \tag{4.22}
$$

$$
2A \longrightarrow \dots \tag{4.23}
$$

The reaction equation [\(4.21\)](#page-75-0) is of first order, which includes isomeric reactions (e.g. $A \longrightarrow B$) where a single educt is converted into a single product, and degradations (e.g. $A \longrightarrow B + C$) where the single educt is split into multiple products. The second order equation [\(4.22\)](#page-75-0) is then assigned to reactions involving two different educts (e.g. $A + B \longrightarrow C$, $A + B \longrightarrow C + 3D$, ...). The homodimer formation reaction (4.23) is finally a special form of the second order equation with two educts of the same type (e.g. $2A \rightarrow B$, $2A \rightarrow B+2C$, ...). The reason why this last kind of reaction cannot be treated in the same way as a second order reaction is that there exist different numbers of possible educt combinations. For a system of N_A particles of A and N_B particles of B the number of possible distinct educts within B is $N_A \times N_B$ whereas the number of possible distinct encounters among A is $N_A(N_A-1)/2$. In [4DiCeS](#page-134-0) third and higher order reactions can be reasonably estimated by the combination of multiple reactions of the supported three types.

From a mechanistic perspective the actual reactants (reaction partners or educts) are piped to the interface by three matrices: one matrix for either of the educts M_E , the internal products $M_{P_{in}}$, or the external products $M_{P_{out}}$. The rows of each matrix specify the reaction equations and the columns define the reaction particle species. As an example the following five reactions can be considered:

(a) Equations		(b) M_E				(c) $M_{P_{in}}$					(d) $M_{P_{out}}$				
Reactions				A P X Y Z		A	\mathbf{P}	X	\mathbf{Y}	Ζ	A	\mathbf{P}	X		Z
(4.24)		$\overline{0}$	$\overline{0}$	$\mathbf{1}$	θ	θ	$\overline{1}$	$\overline{1}$	θ	$\overline{0}$	$\overline{0}$	$\overline{0}$	$\overline{0}$	θ	θ
(4.25)	0	θ			θ	θ	$\overline{2}$	θ	θ	θ	θ	θ	θ	θ	θ
(4.26)		θ		$\overline{0}$	$\overline{0}$	θ	$\overline{0}$	$\overline{2}$	θ	$\overline{2}$	θ	Ω	θ	$\left(\right)$	θ
(4.27)	θ	0	$\overline{2}$	$\overline{0}$	θ	1	1	$\overline{0}$	θ	θ	θ	θ	O	θ	θ
(4.28)	0	θ	θ	θ	1	θ	θ	θ	1	θ	$\overline{0}$	$\left(\right)$	θ		θ

Table 4.11: Internal representation of reaction equations. In [4DiCeS](#page-134-0), reaction equations are represented as matrices. There are three matrices for each element of a reaction equation: one for the educts M_E [\(b\);](#page-76-1) one for the inner products $M_{P_{in}}$ [\(c\);](#page-76-2) and one for the outer products $M_{P_{out}}$ [\(d\).](#page-76-3) The rows specify the different equations [\(a\)](#page-76-4) and the columns $(b)+(c)+(d)$ $(b)+(c)+(d)$ $(b)+(c)+(d)$ indicate the used particle species in the activity description. The cells of the matrices may have numbers from zero to two. Each number indicates the corresponding number of particles reacting in an equation.

$$
A + Y \longrightarrow X + P \tag{4.24}
$$

$$
X + Y \longrightarrow 2P \tag{4.25}
$$

$$
A + X \longrightarrow 2X + 2Z \tag{4.26}
$$

$$
2X \longrightarrow A + P \tag{4.27}
$$

$$
Z \longrightarrow fY \tag{4.28}
$$

These five reactions [\(4.24](#page-76-0) to [4.28\)](#page-76-0) were devised by [\(Belousov,](#page-143-2) [1958\)](#page-143-2) and are called the Belousov–Zhabotinsky Reaction ([BZR](#page-134-3)) today. They generate sustained temporal oscillations in the concentrations of intermediates X , Y , and Z. In an alphabetical ordering of reactants, the three matrices would look like those in Table [4.11.](#page-76-5) There are no external products in the reaction equations (4.24) to (4.28) . Therefore $M_{P_{out}}$ remains as a zero matrix.

The input and output of the number of particles per species (or their respective concentrations) in the current compartment work in much of the same way as with the reaction channel matrices. A matrix of amounts or concentrations is passed to the reaction algorithm were every row represents a particle species. The result is returned as two matrices – one matrix for the resulting internal $M_{R_{in}}$ and another matrix for external changes $M_{R_{out}}$.

Table 4.12: Reaction methods in biology. The columns represent almost the same information as those in Table [4.10.](#page-73-0) With column "Type" there is a new term "hybrid" introduced, which states algorithms that use both stochastic and deterministic methodologies. It generally can be seen that there are in addition to deterministic approaches various derivatives of the Gillespie and the τ -leaping method.

Similar to diffusion there exist various different methods for computational reaction processing. Table [4.12](#page-77-0) provides an overview of important candidates in this category. Deterministic approaches for [ODE](#page-136-1)s as well as [PDE](#page-136-2)s are again commonly utilized. Also stochastic methods such as the Gillespie derivatives and some τ -leaping variations exist. Attempts were made to combine deterministic and stochastic reaction methods to form "hybrid" approaches. Therefore one deterministic and several stochastic methods were chosen as reference implementations for the [4DiCeS](#page-134-0) application. These are Gillespie's "First Reaction Method", Gillespie's "Direct Method", as well as Gibson and Bruck's "Next Reaction Method" working on particles and the Runge–Kutta 4^th order method using concentrations. A brief description of these two algorithms can be found in Appendix [A.1.](#page-118-1) In addition to the two reference implementations, there is an implementation of the hybrid COntrollable Approximative STochas-tic (reaction-algorithm) ([COAST](#page-134-4)) [\(Wagner et al.,](#page-158-3) [2006;](#page-158-3) Möller, [2006\)](#page-153-0) by kind courtesy of Dr. Mark Möller, Bielefeld University, Germany.

4.2.4 User Interfaces

The previous sections described interfaces that can be used to insert application data model relevant modules. This section now provides more detail regarding the human-machine interdependency with the [4DiCeS](#page-134-0) system. The system requires user interaction since various functional and exchangeable modules for diffusion reaction and input description parsing have to be defined manually sometimes. There are two usage scenarios for user interfaces – the Command-Line Interface ([CLI](#page-134-5)) and a Graphical User Interface ([GUI](#page-135-2)). Both use-case scenarios are displayed in Figure [4.20.](#page-79-1)

The [CLI](#page-134-5) (see Figure [4.20\(a\)\)](#page-79-2) allows for initiating simulation processes from a shell-based environment. Alternatively, the application can be opened and further controlled by a [GUI](#page-135-2) (see Figure [4.20\(b\)\)](#page-79-3). While a [GUI](#page-135-2) has the advantage of being fully-manipulatable even during simulation, a [CLI](#page-134-5) allows for the onetime input of command-line options to run the program. Each simulation has to be newly initiated by [CLI](#page-134-5) parameters that can only be interrupted but not manipulated any further by the user. Such [CLI](#page-134-5) simulations can thus be used for multiple high-throughput simulations if embedded into batch-jobs. This implies that one or more application data models exist already and that the user is only interested in the simulation output. With a [GUI](#page-135-2) the user is in a position to manage both the data and the simulation simultaneously. Figures [4.21](#page-80-0) and [4.22](#page-81-0) highlight this difference through the use of two [UML](#page-136-0) sequence

Figure 4.20: Use-cases: User interfaces. The two [UML](#page-136-0) use-case diagrams display the two possible ways the [4DiCeS](#page-134-0) application can be communicate with. The shell-based [CLI](#page-134-5) [\(a\)](#page-79-2) is the simpler one of both scenarios. Here the user activates a simulation via parameters on the command line and receives the result at termination. Beyond that the [GUI](#page-135-2)-based approach [\(b\)](#page-79-3) can additionally manipulate the data the simulation is working on.

diagrams for a [CLI](#page-134-5) and a [GUI](#page-135-2) respectively.

Both sequence diagrams have the application data model instantiation and deletion in common. Their difference lies in the additional interaction points for the user. For the [CLI](#page-134-5) the only additional manipulation the user is able to do besides simulation initialization is canceling the simulation manually before it auto-terminates. The [GUI](#page-135-2) provides the user with much more flexibility as well as the complete control over the data set and simulation methods.

As can be seen with the now following Section [4.3,](#page-79-0) there exist reference implementations for both [CLI](#page-134-5) and [GUI](#page-135-2) interactions. A mixture of both scenarios should be possible as well but was not investigated any further.

4.3 Application Details

Previous sections of this chapter defined the formal structure of the [4DiCeS](#page-134-0) system. This section provides details on how these formal definitions were realized and what the user can expect from it. More significant than other implementation decisions, there was the goal of having good performance and extendability

Figure 4.21: Sequence diagram: [CLI](#page-134-5). As the simpler of the two usage scenarios, this figure presents an [UML](#page-136-0) sequence diagram for a [CLI](#page-134-5). Here it can be seen that the lifetime of the [CLI](#page-134-5) spans the complete model instantiation and simulation process. The user is only able to initialize or abort a simulation. If simulation output is produced, the user can collect it by the time the [CLI](#page-134-5) will terminate.

for the [4DiCeS](#page-134-0) application. Both topics were addressed by choosing a native programming language for the kernel implementation. $C/C++$ came to use as a programming language supporting object-oriented features (C_{++}) and allowing for very performing compiles (C). A drawback to this decision is that the kernel has to be recompiled for every operating system that the program is used on. This was resolved by using only standard elements of $C/C++$ and by continuously cross-compiling the source code on at least one Windows and one UNIX platform during development.

Dynamically extending a program in runtime with new subclasses on different platforms is one of the main design aspects of [4DiCeS](#page-134-0). On UNIX systems this makes use of the Executable and Linking Format ([ELF](#page-135-3)) and Shared Objects

Figure 4.22: Sequence diagram: [GUI](#page-135-2). This figure shows the [GUI](#page-135-2) usage as an [UML](#page-136-0) sequence diagram. In contrast to the [CLI](#page-134-5), the [GUI](#page-135-2) allows the user to perform alternative interactions with the [4DiCeS](#page-134-0) kernel. After an application data model has been set, the user has the choice between running a simulation on the model, manipulate it further, or load a completely different data set. Of course a simulation can also be interrupted by the user. If the user has completed his task, he can instruct the application to terminate. Object instantiation and deletion work for both scenarios equally well.

([SO](#page-136-3)s). It can also take advantage of Windows regular Dynamic Link Librarys ([DLL](#page-135-4)s). Having two incompatible strategies on two different architectures leads to cross-platform-incompatible code.

To better resolve this problem of cross-platform compatibility, Alex Holkner

Figure 4.23: Static diagram: [DynLoad](#page-135-5). The [DynLoad](#page-135-5) is designed as a dynamic link factory method that lets a class defer its instantiation to subclasses. This class checks the current operating system and loads the required file format accordingly. Hence, self-controlled [Dyn-](#page-135-5)[Load](#page-135-5) is able to open the necessary file and dynamically link it to the loading system framework.

[\(Holkner,](#page-149-1) [2002\)](#page-149-1) therefore introduced a simple API, called "[DynLoad](#page-135-5)", which allows developers to write dynamically-linkable code once and for all platforms. A new compilation of code is however required for each destination platform. Holkner's reference code consists of a dynamic link factory method (see Figure [4.23\)](#page-82-0). This method defines an interface for creating an object but allows subclasses to decide which class to instantiate. Here it determines the file format by checking the current operating system. It can hence automatically incorporate the necessary file and dynamically link it to [4DiCeS](#page-134-0).

The interfaces were implemented to allow for an easy connection to other programming languages. Special attention was paid to languages that do not need a recompilation if transferred to another platform. As a first prototype an interface connection to Java is provided. It compiles and runs by a Java Virtual Machine ([JVM](#page-135-6)) implementation, which is natively ported to most current operation systems. Python as another programming language can be used through the Java interface.

This open interface structure should allow for the integration of numerous implementations by other programmers in their favorite programming language. Updates to the modules can easily be integrated to the system for the same reason. This modularity also allows for much more since the system can be directly adjusted to suit the needs of a user. The user interface itself is therefore designed to allow various administrative and analytical methods on the [4DiCeS](#page-134-0) framework.

The following Section [4.3.1](#page-83-0) is concerned with the current realizations of user interfaces to the [4DiCeS](#page-134-0) application. Section (Section [4.3.2\)](#page-85-0) then states what a user can expect from the current state of the application and Section [4.3.3](#page-86-0) thereafter considers the overall computational complexity the [4DiCeS](#page-134-0) system might experience in regards to storage and performance. The last section of this chapter (Section [4.3.4\)](#page-87-0) gives details on where the [4DiCeS](#page-134-0) program and accompanying documentation can be obtained from.

4.3.1 User Interfaces

There are two user interface implementations for [4DiCeS](#page-134-0) at present. Figure [4.24](#page-84-0) demonstrates these realizations as well as one CAVE Automatic Virtual Environment ([CAVE](#page-134-6)) visualization (Figure [4.24\(b\)\)](#page-84-1) of [4DiCeS](#page-134-0) output data. Figure $4.24(a)$ shows a [CLI](#page-134-5) execution. As can be seen a full run of initialization and simulation iterations was performed. Figure $4.24(c)$ then shows a [GUI](#page-135-2) implemented to allow for convenient state-of-the-art user interaction with the software.

The [CLI](#page-134-5) can be used as a command-line shell approach to the [4DiCeS](#page-134-0) project. All necessary settings can be adjusted by giving it lines of textual commands either from keyboard input or from a script. The [4DiCeS](#page-134-0) [CLI](#page-134-5) can perform operations in a batch processing mode without user interaction. If there is thus the need to repeat certain simulations more often, a "script" could hold all the necessary information. The operation can thereafter be conducted without any further effort in analysis. In contrast [GUI](#page-135-2) users have to restart their exploration every time manually.

As for the [GUI](#page-135-2), the user is provided with a tool to view and verify the model and its geometrical structure. The [GUI](#page-135-2) was designed as a Tabbed Document Interface ([TDI](#page-136-4)) that allows for multiple documents to be contained within a single window (see Figure [4.24\(c\)\)](#page-84-3), using tabs to navigate between them. The child documents can be freely resized, positioned, and undocked from as well as docked within the parent window. The single child documents contain different graphical or textual interfaces for either manipulating the model, triggering the simulation environment, or the visualization of simulation output.

The [GUI](#page-135-2) described here was developed by making intensive use of the Qt library by Qt Software (formerly Trolltech)^{[12](#page-83-1)}. It should be easily possible to implement an

¹²Qt Software: <http://www.qtsoftware.com/>

(c) [GUI](#page-135-2) Realization

Figure 4.24: The user interface for [4DiCeS](#page-134-0). Controlling a program on the command line [\(a\)](#page-84-2) is not very convenient, but has the advantage of being able to use massive computing. This list is then automatically processed and runs without additional user input until completion. A [GUI](#page-135-2) [\(c\)](#page-84-3) is needed if there is a user who wants to see direct runtime results and needs to make changes to the system as it runs. Here the front panel of the current realization of the [4DiCeS](#page-134-0)-[GUI](#page-135-2) is shown. The [GUI](#page-135-2) allows [4DiCeS](#page-134-0) users to choose their desired application data model, and manipulate it during program run-time. Lastly a picture [\(b\)](#page-84-1) was taken of Julie Stromer (at the University of Calgary, Canada) handling a [4DiCeS](#page-134-0) visualization inside of the [4D](#page-134-7) virtual environment ([CAVE](#page-134-6)). 79 equivalent or even more powerful [GUI](#page-135-2) with another library and the currently supported set of programming languages. As a feasibility study, the [4DiCeS](#page-134-0) user interface mechanism was connected to a [4D](#page-134-7) virtual environment (see Figure [4.24\(b\)\)](#page-84-1) ([CAVE](#page-134-6)) [\(DeFanti et al.,](#page-145-2) [1993\)](#page-145-2). This implementation used Java and its [3D](#page-134-8) extension library Java3D.

4.3.2 Applicability

Although the [4DiCeS](#page-134-0) kernel is the central part of this work, the application itself is mainly accessed through user interfaces. Therefore this section specifies the human-machine interaction with [4DiCeS](#page-134-0), its possibilities, and, of course, its limitations in greater detail.

As already stated in the previous sections two distinct reference implementations of user interfaces exist at present. The first, the [CLI](#page-134-5), serves as an automation utility for massive throughput simulations. Modeling is not a concern of this user interface, since no further interaction can be accomplished after the application was started in this way. On the contrary the [GUI](#page-135-2) reference implementation mainly concentrates on visualizing the design of a model, geometry details, and the output of a simulation.

The model design functionality ranges from importing existing models to freely defining new model scenarios off the scratch. Here various child windows of the parent [TDI](#page-136-4) window support the user in editing reaction and diffusion parameters as well as simulation space settings. There currently exist textual editors for setting reactions, compartments, as well as initial parameters. All dialogs together form a tool to model reaction-diffusion systems and specify further dependencies as of connected algorithms and the geometry information.

Since both the [4DiCeS](#page-134-0) kernel and the application data model provide support for [3D](#page-134-8) simulations, the [GUI](#page-135-2)'s visualization output can also be displayed in [3D](#page-134-8). This task comes with some difficulties and limitations. As everything is simulated at once, it is difficult for the user to trace all the occurring events in a compact [3D](#page-134-8) image on a [2D](#page-134-9) display. Differentiating also between different particle species is limited to a small set of different particle species at a time due to the internal computer coloring scheme.

At present the underlying reference geometry description definition (cell model) is visualized as a [3D](#page-134-8) grid of cubes covered by a combination of the three fundamental colors obeying the red, green, and blue ([RGB](#page-136-5)) color scheme. Each color represents one particle species, and the shading of this color indicates the particle concentration in that specific voxel. Since there are 256 different shades for each color on a computer, this number limits proper differentiation of different particle concentrations. For the user to be able to get the most out of such a data display, [4DiCeS](#page-134-0) provides features to manipulate the proportionality scale of colors to represent logarithmic changes in concentration. The total change of concentration is displayed by a [2D](#page-134-9) plot next to the [3D](#page-134-8) visualization.

As can be seen from attempts made in the following Chapter [5](#page-88-0) in testing the system itself (see Section [5.2\)](#page-91-0) and the comparison study with related tools (see Section [5.3\)](#page-98-0) [4DiCeS](#page-134-0) is a well suited tool for the modeling and simulating of biochemical reaction-diffusion systems. Furthermore it is feasible for the analysis of highly dynamic models as of the oscillating [BZR](#page-134-3) (see Section [5.2.2.2\)](#page-96-0) and the two pool calcium oscillation model (see Section [5.2.2.3\)](#page-96-1). The now following section will consider the overall computational complexity the [4DiCeS](#page-134-0) system might experience with such experiments in regards to storage and performance.

4.3.3 Complexity Considerations

This section is concerned with the prediction of resources that the [4DiCeS](#page-134-0) application requires in total. The influence of the algorithms (see Appendix [A\)](#page-118-0) is actually very minimal, since they all grow asymptotically slower than n^2 and the traversal of [VE](#page-136-6)s in [3D](#page-134-8) during one simulation iteration already yields an $O(n^3)$ upper bound on the overall worst-case running time.

But the discrete Lévy process should then be considered separately. As already stated earlier this algorithm does not depend on the grid of [VE](#page-136-6)s in the simulation space, but only traverses over all particles during one iteration. Anyhow, by the time the result output has to be stored to the kernel then again all [VE](#page-136-6)s have to be traversed. Algorithms that have to perform some sort of initialization routines sum up this overhead during a simulation over more than one [VE](#page-136-6) over time.

Therefore it can be said that the overall worst-case performance of [4DiCeS](#page-134-0) converts in $O(t \times r \times p \times x \times y \times z)$, were t represents time, r denotes the number of reactions, p is the total number of particle species, and the product of x , y , as well as z describes the total number of [VE](#page-136-6)s in [3D](#page-134-8). The storage worstcase complexity is growing in a very similar asymptotically expansion, because the same six-tuple has to be considered when observing the overall memory consumption.

4.3.4 Availability

The [4DiCeS](#page-134-0) project, along with all its sources and relevant documentation is hosted by SourceForge^{[13](#page-87-1)}. The project's domain is [http://www.4DiCeS.de/.](http://www.4DiCeS.de/) The domain address is automatically forwarded to the project's web-content at SourceForge [\(http://four-dices.sourceforge.net/\)](http://four-dices.sourceforge.net/).

For compiling the [4DiCeS](#page-134-0) sources, additional programs, libraries, as well as extensions are required as can be seen in Table [4.13.](#page-87-2) A detailed description for various C/C++ Integrated Development Environments ([IDE](#page-135-7)s) is included in the [4DiCeS](#page-134-0) documentation. The current version of [4DiCeS](#page-134-0) has been tested and developed under Linux and Windows. As long as the tools and packages described above are in place, [4DiCeS](#page-134-0) should install and run smoothly.

Table 4.13: Required third-party tools. Perl and CxxTest are only for compiling and running [4DiCeS](#page-134-0) unit tests. Qt is required for compiling and running the [4DiCeS](#page-134-0)-[GUI](#page-135-2). The use of Java depends on the need for this specific programming language on plug-ins within [4DiCeS](#page-134-0).

The source-code as well as the compiled program(s) are licensed under the GNU Lesser General Public License $(LGPL)^{14}$ $(LGPL)^{14}$ $(LGPL)^{14}$ $(LGPL)^{14}$ $(LGPL)^{14}$ and are therefore free for noncommercial use. Commercial usage can only be granted after consultation with the current source-code administrator.

¹³SourceForge: <http://sourceforge.net/>

¹⁴[LGPL](#page-135-8): <http://www.gnu.org/licenses/lgpl.html>

Applications, Results, and Analysis

In this work the [4DiCeS](#page-134-0) framework was developed for modeling and simulation of cellular dynamic phenomena in [3D](#page-134-8). As described in the previous chapter, the implementation of this project includes a modular design of the [4DiCeS](#page-134-0) kernel and four modular interfaces. Various reference implementations of algorithms add up to these interfaces (see Table [5.1\)](#page-88-1).

Table 5.1: Interface reference implementations. The [4DiCeS](#page-134-0) framework comprises a variety of different reference implementations for the given four interfaces. Additionally to these modules an implementation of the COntrollable Approximative STochastic (reaction-algorithm) $(COAST)$ $(COAST)$ $(COAST)$ is also available by kind curtesy of Dr. Mark Möller.

With these first modules it was possible to test and verify the framework in its behavior and reliability. Of special interest was the impact of diffusion to well known reaction algorithms. The propagation of wave-like structures of oscillating model systems was expected but not detected. Although it was possible to generate massively random 3D grids out of uniformly distributed particles there were no concrete pattern formations perceived. A discussion of this finding and all other results will be handled in Chapter [6.](#page-108-0)

The following sections now want to give insight to further outcomes in regards to application development (Section [5.1\)](#page-89-0), applied models for simulation (Section [5.2\)](#page-91-0), as well as ten reference programs in comparison to [4DiCeS](#page-134-0) (Section [5.3\)](#page-98-0). The last Section [5.4](#page-106-0) summarizes the findings from the comparison of [4DiCeS](#page-134-0) with the other ten tools.

5.1 Application Development

As previously stated the main characteristic of the simulation tool is its central simulation kernel and the interfaces to the reaction as well as diffusion modules, the model input parsing environment, and the visualization (see Table [5.2\)](#page-90-0). Front-end user modules support both the [CLI](#page-134-5) processing and a [GUI](#page-135-2)-based version.

This section examines the resulting plug-in structure of the module interfaces. The possible connectivity to other programming languages then follows directly out of this architectural definition.

5.1.1 Plug-Ins

Every interface presented here (see Tables [5.1](#page-88-1) and [5.2\)](#page-90-0) has a special module extending the interface to expand to other programming languages (see Figure [5.1\)](#page-91-1). These modules make use of the underlying [4DiCeS](#page-134-0) kernel [API](#page-134-10). The [API](#page-134-10)'s functionality again is passed on to the respective language module plug-ins.

A plug-in generally consists of a module instantiation function and an object destruction facility to enable a safe shout-down behavior of the kernel. These two maintenance methods provide a simple interface to all plug-in interfaces. The instantiated modules are then specialized forms of either interface class. These classes bring along the connection to the underlying model description, kernel functionality, and offer distinct entry points to execute their particular operation.

In the case of the algorithmic modules the entry points are functions named diffuse and react. These functions either work with entire segmentations (compartments) or on single voxels. The kernel as well as the model binding are rather limited to methods for retrieving information. The modification of data

Table 5.2: [4DiCeS](#page-134-0) interface definitions. This table briefly describes the kernel interfaces. The reaction and diffusion processes are all handled within the simulation space. The simulation interfaces (reaction and diffusion) are computationally more demanding than the input parsing modules. There are normally only a few loading and saving events during a simulation. Hence there is no need for highly optimized computational efforts on this topic. The visualization interface can be of both types if [CLI](#page-134-5) and [GUI](#page-135-2) applications are compared.

is restricted to the extension that an algorithm must not alter the structure of a model. It should only modify its values.

In contrast the parsing and visualization interfaces have full control over a model. The main difference between these two is the connectivity to the kernel. While a parser plug-in is controlled by the kernel via its entry points, parse and save, the visualization has full control over the kernel. The kernel actually acts as a plug-in to the user interface rather than the user interface to the kernel. This structure allows for the easy integration of simulation results into other existing tools and frameworks for further analysis.

Figure 5.1: Static diagram: A general plug-in interface. The [4DiCeS](#page-134-0) [API](#page-134-10) provides users with the choice of which algorithm to use, and programmers with the ability to modify given modules or to implement new ones. The [API](#page-134-10) provides additional programming language plug-ins for each interface.

5.1.2 Programming Languages

 $C/C++$, Sun Microsystems Java, and Python were selected as the project's programming languages. All are commonly used in computational biology and have their respective benefits. $C/C++$ was used for the system's kernel and the interface implementation. The other programming languages were chosen as extensions to guarantee an easy [API](#page-134-10) to the [4DiCeS](#page-134-0) framework. Thus the widespread use of either language will allow many researchers to use [4DiCeS](#page-134-0) as a workbench for their implementation of algorithms and tools.

All libraries used were carefully chosen to allow for an easy migration of the [4DiCeS](#page-134-0) framework between operation systems. The current target platforms are the Microsoft Windows operating systems (95 and upwards) as well as current POSIX derivatives including Linux. It should be easy to port the program to other operation systems by a simple recompilation of the given C/C++ sources.

5.2 Modeling and Simulation

As described in Section [4.2.1,](#page-51-0) the geometrical representation of the [4DiCeS](#page-134-0) data model consists of a grid of [VE](#page-136-6)s. As each [VE](#page-136-6) can be covered by membrane, [3D](#page-134-8) compartmental substructures can be included to the grid. Hence it is possible to define particle barriers in the [3D](#page-134-8) simulation space. This also means that different kinds of boundary conditions can be modeled in the system. Before the two models from Section [5.3](#page-98-0) were applied a few test simulations examined

the overall functionality of the system.

The first subsection illustrates a simple diffusion application in [4DiCeS](#page-134-0) to test the system's [2D](#page-134-9) and [3D](#page-134-8) capabilities. Then the second subsection demonstrates diffusion-reaction systems from simple chemical reactions to finally a non-linear system such as the [BZR](#page-134-3). All grids are divided into 9, 261 [VE](#page-136-6)s $(21 \times 21 \times 21)$ grid) with an edge length of 100 nm if not stated otherwise. The analyzed time-points are enclosed with the figures showing the simulation results.

5.2.1 Simple Diffusion Simulation

The testing of the diffusion functionality of [4DiCeS](#page-134-0) was deemed early on as being of great value. The geometry of models and the diffusion within the modeled geometry had to be verified in case studies involving [2D](#page-134-9) as well as [3D](#page-134-8) space. It was therefore necessary to design test models of only the diffusion components at first. Two of these models will be presented in the following subsections. The membrane-bound diffusion model in lateral ([2D](#page-134-9)) and the [3D](#page-134-8) displacement model are illustrated first. Thereafter a more complex membranous system will be introduced, that tests if boundary conditions are securely handled by [4DiCeS](#page-134-0).

5.2.1.1 Lateral and [3D](#page-134-8) Diffusion

The simplest diffusion test described here consists of a cubic grid of [VE](#page-136-6)s. A continuous $x-y$ -layer of membrane is set near to the bottom of the grid defining two separate compartments within the grid. Four of the six simulation space boundaries are not limited by membranes. Their boundary conditions at these four borders allow for the direct migration of particles from one side of the simulation space to the opposite side. The only exceptions to this are the top and the floor layer. Membranes have to prevent particles from diffusing out of the top [VE](#page-136-6)s and into the bottom layer, which defines the other compartment. Thus the geometrical model is a closed system with a membranous barrier to the bottom of the grid. Two diffusible particle species are placed in this model. One is assumed to translocate laterally over the membrane, whereas the other is permitted to move freely in [3D](#page-134-8). The membrane is set to hinder the freely diffusing particles from switching from one compartment to the other.

As can be seen from Figures [5.2](#page-93-0) [\(a\)](#page-93-1) to [\(c\),](#page-93-2) this simple application of diffusion works as expected. Both the lateral and the [3D](#page-134-8) diffusion mechanisms show anticipated translocation. The dividing membrane stops the further diffusion

Figure 5.2: Diffusion application. The Figures [\(a\)](#page-93-1) to [\(c\)](#page-93-2) illustrate the simple diffusion of particle species on a lateral ([2D](#page-134-9)) membrane and in free [3D](#page-134-8) space. With Figures [\(d\)](#page-93-3) to [\(f\)](#page-93-4) the grid is divided into two compartments again, but with an opened tube in the membrane's middle. For both applications, the green particles are bound to the membrane of the model and can only diffuse along this barrier. Conversely, the red particle species is allowed to diffuse freely in [3D](#page-134-8) space with the only limitation of not being able to cross a membrane. In the second model the red particles can pass the tube's top opening (see arrow in Figure [\(e\)\)](#page-93-5)

of the [3D](#page-134-8) diffusing particle species to the other side of the boundary. Also the laterally diffusing particles stay to their membrane bound condition as expected.

5.2.1.2 Complex Membrane System

Here the diffusion is tested on a more complex membranous system. The overall model is still a closed system. Particles are still able to leave four sides of the grid at one side and reenter it on the opposite side. The top and bottom layer are blocked for diffusion. Instead of the previous plain membrane layer near the floor of the grid, a membranous tube is placed on top of the structure. The tube is uneven and opens to the top of the grid. The simulation of diffusion (see Figures [5.2](#page-93-0) [\(d\)](#page-93-3) to [\(f\)\)](#page-93-4) begins with freely diffusible particles at the bottom right corner of the grid. Next to the [3D](#page-134-8) diffusion, the [2D](#page-134-9) diffusion is started. The particles diffuse along the tube to its top opening. At the end of the simulation, the [3D](#page-134-8) particles were able to enter the opposite compartment through the tube's top opening.

5.2.2 Applying Diffusion-Reaction Systems

Besides diffusion the reaction capability of the [4DiCeS](#page-134-0) framework was also tested. Since reaction systems ignoring diffusion would only evaluate the reaction algorithm by its own and not the [4DiCeS](#page-134-0) framework completely, two diffusion-reaction systems were applied for testing. At first a simple chemical reaction was verified in a membrane-free grid as well as under boundary conditions. Then a [BZR](#page-134-3) was used to heavily test oscillating chemical systems.

5.2.2.1 Common Saturation Reactions

The test of a simple chemical reaction system of the form

$$
A \stackrel{B}{\rightleftharpoons} C \tag{5.1}
$$

was handled using two different geometrical models (see Figure [5.3\)](#page-95-0). The first model is completely free of any membranous boundaries, whilst the other contains a grid structure equal to that described in Section [5.2.1.1.](#page-92-0) The reaction consists of an educt A , a product C , and a catalyzing enzyme B . In the first simulation experiment (see Figures [5.3](#page-95-0) [\(a\)](#page-95-1) to [\(c\)\)](#page-95-2), all particle species (A, B, \mathcal{L}) and C) are allowed to diffuse freely in [3D](#page-134-8) space. In the second case (see Figures [5.3](#page-95-0) [\(d\)](#page-95-3) to [\(f\)\)](#page-95-4), the enzyme species B is limited to the dividing membrane.

Figure 5.3: A simple chemical reaction. The reaction system of the form $A +$ $B \rightleftharpoons B + C$ is tested with the [4DiCeS](#page-134-0) framework. Figures [\(a\)–](#page-95-1)[\(c\)](#page-95-2) illustrate the reaction without any membranous barriers and (d) – (f) show the same reaction system with a membrane layer. The colors green and red stand for the reactants in this system. In (d) – (f) the green particles (enzymes) are membrane bound and must diffuse laterally along the membrane. The Figures [\(c\)](#page-95-2) and [\(f\)](#page-95-4) show the reaction product C as a blue particle species. Here, in (c) both educts A as well as B and in [\(f\)](#page-95-4) only A were hidden away to have a better display of the product.

As can be seen, the reactions start with particles of both educts present in the hosting [VE](#page-136-6)s. The simulation was halted shortly after the first products appeared. As expected the educts A continuously react in the presence of enzyme B to the product C . The membrane-divided model needed more time for first reactions to occur. This was expected since the educts A had to move the entire distance to the anchored enzymes B.

5.2.2.2 Sustained Oscillating Reaction

The [BZR](#page-134-3) (see Equations [4.24](#page-76-0) to [4.28](#page-76-0) at Section [4.2.3.3\)](#page-74-0) is probably the most widespread oscillating reaction system both theoretical and experimental. On the theoretical side, the Field–Körös–Noyes ([FKN](#page-135-9)) model system [\(Field and](#page-147-1) [Noyes,](#page-147-1) [1974a](#page-147-1)[,b\)](#page-147-2) quantitatively mimics the actual [BZR](#page-134-3)s [\(Field et al.,](#page-146-3) [1972\)](#page-146-3). The Brusselator [\(Glansdorff and Prigogine,](#page-148-6) [1971;](#page-148-6) [Nicolis and Prigogine,](#page-153-1) [1971\)](#page-153-1) as well as the Oregonator [\(Noyes,](#page-154-0) [1976a](#page-154-0)[,b\)](#page-154-1) are then variations of the [FKN](#page-135-9) model system [\(Ipsen et al.,](#page-150-2) [1997\)](#page-150-2).

Boris Belousov stated that in a mix of potassium bromate, cerium sulfate, propanedioic acid and citric acid in dilute sulfuric acid, the ratio of concentration of the cerium(IV) (Ce^{4+}) (Ce^{4+}) (Ce^{4+}) and cerium(III) (Ce^{3+}) ions in combination with bromate (V) ([Br](#page-134-13)) oscillated.

Basically the reactions can be divided into two parts. The concentration [[Br](#page-134-13)-] determines which part is dominant at any time. While there is a high [[Br](#page-134-13)-] the first part is dominant. During this stage [Br](#page-134-13) is consumed and the cerium is mainly in the Ce^{3+} Ce^{3+} state. As $[Br]$ $[Br]$ $[Br]$ decreases further it passes through a critical value and then drops quickly to a low level. At this stage the second part takes over. The Ce^{3+} Ce^{3+} changes to Ce^{4+} . Ce^{4+} reacts to produce [Br](#page-134-13) again while it reverts to the Ce^{3+} Ce^{3+} state. Now $[Br]$ $[Br]$ $[Br]$ increases. By the time $[Br]$ is sufficiently high the first part gets dominant again [\(Murray,](#page-153-2) [2002a](#page-153-2)[,b\)](#page-153-3). The whole sequence is continually repeated and hence produces colored oscillations.

The test environment in the [4DiCeS](#page-134-0) framework considers no internal membrane barriers. Only the grid side walls are covered with membrane to change the model to inflexible boundary conditions. The particle species are initially distributed equally over the grid of [VE](#page-136-6)s. As expected it can be seen (see Figure [5.4\)](#page-97-0) that an oscillation of the reactants' concentrations takes place with a cycle duration of nearly 30 s in the hosting [VE](#page-136-6)s. The simulation shows a short latency of about 8 s at the very beginning. The oscillation process was halted after one full cycle $(t = 55 s)$. There are no distinguishable wavefronts, but that is due to the rather big [VE](#page-136-6)s compared to the rather small grid.

5.2.2.3 Two Pool Oscillatory Calcium Model

The two pool model (see Section [2.2.2\)](#page-20-0) describing oscillating Ca^{2+} Ca^{2+} changes within a cellular model was also applied to [4DiCeS](#page-134-0). Figure [5.5](#page-98-1) shows the output of four simulation iterations with slightly varying parameters. This difference in

Figure 5.4: An oscillating reaction. Here, the [BZR](#page-134-3) system was tested with the [4DiCeS](#page-134-0) framework. The images (a) –[\(j\)](#page-97-2) illustrate a stirred 60 ml beaker with BZ reactants. Adapted from [Winfree](#page-158-4) [\(1987,](#page-158-4) [2001\)](#page-158-5). The images (k) – (t) show the same [BZR](#page-134-3) from a [4DiCeS](#page-134-0) simulation. The times-step for the simulation was set to 1 s and the edge length of the [VE](#page-136-6)s was 5 mm.

output can only be seen, due to the fact that the stochastic model was applied. The deterministic version of the model does not show this behavior [Kraus et al.](#page-151-1) [\(1992\)](#page-151-1). All parameters used for the simulation are stated within the caption of Figure [5.5.](#page-98-1) The application of a diffusion term and the simulation of the pool model in [3D](#page-134-8) did not show any wave-patterns. Reasons for this behavior could be the coarse granularity of the simulation space $(21\times21\times21)$ or a bad selection of either the time-steps or the edge length of the [VE](#page-136-6). Although, varying these geometry parameters did not produce any better results. The overall particle number dynamics output was very similar to [Kraus and Wolf](#page-151-2) [\(1992\)](#page-151-2).

Figure 5.5: Stochastic Simulation of Two Pool Model. The two pool model was stochastically simulated with different strong stimulations of the receptors expressed through β . The diagrams show the number of Ca^{2+} Ca^{2+} ions in the cytosol (N_x) as a function of time t. The parameters for [\(a\)](#page-98-2) to [\(c\)](#page-98-3) are (see Table [2.4\)](#page-23-0): $n = m = 2, p = 4, v_0 = 1000s^{-1}$, $v_1 = 7300 \text{s}^{-1}, k = 10 \text{s}^{-1}, k_f = 1 \text{s}^{-1}, K_2 = 1000, V_{M2} = 32500 \text{s}^{-1},$ $k_R = 2000, K_A = 900, \text{ and } V_{M3} = 2.5 \times 10^5 \text{s}^{-1}$. The parameters of [\(d\)](#page-98-4) are chosen that way that they are by a factor of 1000 greater than [\(a\)](#page-98-2) to [\(c\).](#page-98-3) The specific β s can be found in the corresponding caption of every sub-figure. The times-step for the simulation was set to 0.01 s and the edge length of the [VE](#page-136-6)s was 0.095 μ m. The model description and all its reaction parameters were extracted from [Kraus and Wolf](#page-151-2) [\(1992\)](#page-151-2).

5.3 Comparing Tools

This section continues the comparison of the ten modeling and simulation tools from Section [3.2.4.](#page-33-0) In addition this study brings all presented tools in context with the [4DiCeS](#page-134-0) application, which was introduced in Chapter [4.](#page-36-0) The first Section [5.3.1](#page-99-0) will reconsider the criteria important for comparing all the applications. Then the following Section [5.3.2](#page-99-1) will compare the applications to each other in more detail; while, Section [5.3.3](#page-105-0) will deal with [4DiCeS](#page-134-0).

5.3.1 Comparison Criteria

Next to the 12 comparison criteria defined in Section [3.2.4](#page-33-0) it was crucial to see how the applications handle equal test scenarios. Therefore the final tests for every tool were two oscillating test models.

As a straight forward model, the previously described [BZR](#page-134-3) (see Section [5.2.2.2\)](#page-96-0) was applied either in the form of differential equations with concentrations or by reaction equations with particle numbers. This first test case requires no compartmentalization of the model in any way. The second test is based on the two pool model (as described in Section [2.2.2\)](#page-20-0). Here the simulation space is subdivided into two distinct compartments: the cytoplasm, and the [ER](#page-135-10).

Both models should show oscillation as a result similar to the one found in literature. Depending on the ability of the application to provide spatial simulations, it will further on be attended to apply both test cases to [3D](#page-134-8). Again wave-structures (as found in literature) should be observable.

5.3.2 Comparing Related Works

For this comparison it was important to consider the very different strategies of the applications, rather than bringing almost equal tools into contrast. The choice of the given applications is therefore a mixture from deterministic, stochastic, and hybrid approaches.

Some tools are already rather dated, some allow for [3D](#page-134-8) simulations, and some have proprietary or standardized file format extensions (see Tables [5.3,](#page-101-0) [5.4,](#page-102-0) and [5.5\)](#page-103-0). The ten simulators were tested against the test models as defined in the last section. Here the implementations of the test cases were not necessarily straight forward to all compared programs.

A feasible implementation of the model was however possible with almost all ten applications. The results of these simulations are found to be nearly convergent for most of the tools. Variations observed may be due to different methods for solving the equations or application data model incompatibilities. In the following the results are presented in further detail.

[BIOCHAM](#page-134-15): Installing [BIOCHAM](#page-134-15) was accomplished with a simple button click. The only prerequisite is for the computer to have a [JVM](#page-135-6). The application tested was version 2.5, which can be used under the terms of the GNU General Public License ([GPL](#page-135-11)). [BIOCHAM](#page-134-15) allows for simulations with either a boolean

[CA](#page-134-2) algorithm, the Runge–Kutta 4^{th} order solver, or the Rosenbrock's 3^{rd} order solver. With this pure deterministic simulator, the modeling and the results can only be displayed in one dimension. Models can either be defined in [BIOCHAM](#page-134-15)'s own description language (BM) or by [SBML](#page-136-8) files. The simulation results are displayed in a common [2D](#page-134-9) plot and can be viewed as a table. Explicit export functionality for the results is however missing. Both testing models could be applied to the application by setting up the reaction rules in BM. Importing a model through valid [SBML](#page-136-8) files did not work well. Both the boolean and the numeric solvers gave the same predicted results, and it can be said to be accurate and very fast $(< 1 sec)$.

[BioNetS](#page-134-16): The download of [BioNetS](#page-134-16) is split into several packages. For the application to run, all these packages have to be installed. In particular for version 2.0 of [BioNetS](#page-134-16), a [JVM](#page-135-6) and GNUStep, a cross-platform, object-oriented framework for desktop application development must be present. Since the application produces C++ code from given models, there is always a compilation iteration preceding any simulation attempt. Models can then be simulated with either a Langevin equation solver or by the use of the Next Reaction Method [\(Gibson](#page-148-4) [and Bruck,](#page-148-4) [2000\)](#page-148-4). The description of models is only possible with [BioNetS](#page-134-16)'s proprietary [XML](#page-136-9) file format BNET, but a modeling [GUI](#page-135-2) takes away any need to cope with that directly. The setup of the two testing models was difficult for another reason. Both solvers can handle only very limited particle amounts that is why the models' particle numbers had to be condensed a 1000-fold. Doing so resulted in the predicted results for both solver classes. Simulation times ranged from below one second for the Langevin solver and several minutes for the next reaction method. The application can be used under the terms of the Berkeley Software Distribution ([BSD](#page-134-17)) license.

[Copasi](#page-134-18): This simulation tool was tested in the beta version 4.0.19 and was easy to install. The tool's dependency on the rather expensive [GUI](#page-135-2) library Qt is only of importance if the user wants to recompile the application from source code. This is possible since the source code is provided under the [Copasi](#page-134-18) Non-Commercial License ([CN-CL](#page-134-19)) and is therefore free for academic usage. Import and export file formats are in CPS, GPS, and [SBML](#page-136-8). As for simulation methods, the programm ranges from a pure [ODE](#page-136-1) solver, to stochastic algorithms, to hybrid forms. The modeling of reactions leaves nothing to be desired. The tool is able to perform steady-state, time course, as well as parameter estimation simulations. In the case of the two applied test models, [Copasi](#page-134-18) provided expected results in the time-range of seconds to several minutes depending on the algorithm used. The results are displayed in plots and tables but cannot be exported as such directly.

Table 5.3: Tested simulation tools. The table lists all simulation tools tested in this chapter. Here the specific features for simulation are applied methods and algorithms (Column 2), the version number of the tested software (Column 3), the license agreement depending on the version number (Column 4), and further software dependencies to the application are included. The methods of Column 2 are further described in Table [5.4](#page-102-0) with the exception of the xCellerator software, which utilizes the equation solvers provided by Mathematica. A superscripted uncapitalized 'h' determines an hybridized form of the given method. If there are both variants (hybridized and plain), then the uncapitalized 'h' is parenthesized. A subscripted number gives further information about the order of consistence an [ODE](#page-136-1) solver is working on. A superscripted 'b' with version numbers in column 3 denotes a beta-version. The dependencies in column 5 contain special footnote-marks that indicate that the additional package is either provided as a separate download (†), or is already included (∗) in the installation routine of the software package. All other dependencies are prerequisites to either installation or running of the program and must be obtained and installed separately. Further information and links are provided by Section [3.2.](#page-26-0)

Table 5.4: Contained algorithms. This table provides mnemonics for the methods and algorithms applied to the different software packages described in Table [5.3.](#page-101-0) Further information on such algorithms can be obtained from Section [4.2.3.3.](#page-74-0)

E-Cell: The E-Cell programm comes in version 3.1.105 under the terms of the [GPL](#page-135-11) and can be easily installed. The prerequisite packages of the Python scripting language and the GTK-[GUI](#page-135-2) library are automatically installed if not present. The modeling is aided graphically. Model description languages are the E-Cell's own EM/EML formats and [SBML](#page-136-8). For simulation the user can choose between Differential-Algebraic Equation ([DAE](#page-135-15)) solvers and an implementation of the next reaction method. The applied test models had to be remodeled with the help of E-Cell's model editor but ran as expected afterwards. Simulation time ranged from a few seconds to several minutes depending on the used simulation method. The output can be visualized both by plots as well as tables. The export of the data is thus not explicitly possible. It should be noted that the application could not be restarted after the computer was turned off and on again after the first installation.

[Gepasi](#page-135-12): As the oldest of all the tested applications, [Gepasi](#page-135-12) is also the smallest in size and the fastest when it comes to comparison of [ODE](#page-136-1) solvers. As of the

Table 5.5: Supported file extensions. All tested applications have their own proprietary file formats (denoted by an '∗'). Most of the packages support the import and export of [SBML](#page-136-8) files. In the case of [Copasi](#page-134-18), the precursor file format from [Gepasi](#page-135-12) is still supported. Further on, the VitualCell allows for model exports to the [CellML](#page-134-21), Mathlab, and the value chain markup language (VCML). This last file format is primarily used as a business collaboration standard and not for modeling in particular.

last released version 3.30, which is provided as freeware, the installation was easy. The only simulation methodology provided by [Gepasi](#page-135-12) is the LSODA solver for stiff and non-stiff systems. A straight forward [GUI](#page-135-2) and the import as well as export formats for its own model description file GPS and the support for earlier versions of [SBML](#page-136-8) round out its appearance. The tested models both gave good results, which could be displayed in a [2D](#page-134-9) plot. The plot is produced by the required GNUPlot package, which comes with [Gepasi](#page-135-12) already. As stated before the simulations were very fast and ranged for both test cases within a second of total execution time.

[MCell](#page-135-13): As a purely stochastic application with its own reaction and diffusion algorithm implementation, [MCell](#page-135-13) stands apart from all other simulation programs considered here. It is a command-based tool in version 2.50. Running under the terms of the National Resource for Biomedical SuperComputing Soft-

ware License ([NRBSC-SL](#page-136-10)), [MCell](#page-135-13) is free for academic usage and can be run in combination with a visualization utility named DReAMM. The dependencies to other software packages are limited to Cygwin for [MCell](#page-135-13) alone and OpenDX as well as for UNIX X-Server in the case of DReAMM. The download page provides a single Cygwin-[DLL](#page-135-4) that makes installation very simple. Since [MCell](#page-135-13) only supports its own proprietary file format [MDL](#page-135-16) and gives no further [GUI](#page-135-2) tool at hand, the modeling of the test cases was very complex. Obtaining each of the simulation results took almost a day of computational time and results were far from what was expected. Both test cases showed no oscillations even after several modifications of the model.

SmartCell: Next to [MCell](#page-135-13) SmartCell was not able to provide results on the simulation test cases as expected. The reaction equations, which can only be set up by SmartCell's graphical model editor, seem to always lose connection between reactants and their multipliers. Also SmartCell only supports [SBML](#page-136-8) as a helper description that requires a SmartCell's [XML](#page-136-9) reaction placement description in addition. Without this file SmartCell will neither let the user load nor run a single [SBML](#page-136-8) model. SmartCell has many current versions ranging between version 2.5 and version 3.0. The current beta version of SmartCell 3.0 was tested and did not produce any output for the applied models. Consequently the promising [3D](#page-134-8) features, which make use of the required Java3D, could not be tested to their full extent. SmartCell is available under the terms of the Academic Software License ([ASL](#page-134-20)) and utilizes the stochastic next reaction method for both reaction and diffusion process with version 2.5.

[StochSim](#page-136-11): Next to [Gepasi](#page-135-12) [StochSim](#page-136-11) is the oldest stochastic simulation tool of these considered here. It comes with its own algorithm for simulating every two possible reaction partners in discrete time steps. The installation of [StochSim](#page-136-11) is similar to [BioNetS](#page-134-16) with regards to unpacking and putting together of different files. If this is done and the required scripting language Perl and the [GUI](#page-135-2) library Tcl/Tk are installed beforehand, the application is ready to run. In version 1.6 [StochSim](#page-136-11) runs under the terms of the [LGPL](#page-135-8) and can therefore be used freely. As for a model description format, it supports [SBML](#page-136-8). The test case models applied to the system showed the expected behavior although under high computational costs. The overall simulation time for both models reached nearly an hour for every simulation turn. This is even more surprising, since [StochSim](#page-136-11) gives no support for [3D](#page-134-8) modeling. The results can be either directly viewed on the [GUI](#page-135-2) or further analyzed by third-party table-calculation utilities.

VirtualCell: As the most powerful application of all discussed with respect to modeling editors and [3D](#page-134-8) geometry, the VirtualCell is also very different to the others. A server-based approach was chosen that handles the model storage and provides the necessary computational power. Then the Java-based client helps with the design and the setup of the model. The tested beta version 4.3 of the client can be used freely, but a user has to log on to the server every time he wants to use the application. This presupposes an internet connection for program use. VirtualCell provides various [ODE](#page-136-1) solvers, a very sophisticated model editing environment, and [3D](#page-134-8) functionality. The test case could be easily applied via the [SBML](#page-136-8) import and gave good results in [1D](#page-134-22). Although the models both oscillated in [3D](#page-134-8), the particles did not diffuse in space as predicted. This of course can only be a fault in the model design. Another feature of the VirtualCell is its ability to export models and output results directly to predefined formats including [CellML](#page-134-21), Mathlab, and the value chain markup language (VCML). This last file format can be used for the output data.

xCellerator: The xCellerator is very special because it is an 'add-on' to the Wolfram Mathematica package. Here the great benefit is the fact that all solvers and further features of input/output or visualization provided by Mathematica can directly be utilized. A major drawback is the fact that Mathematica is only available commercially. The description of models is done by Mathematica files that obey the additional palette rules of xCellerator. All resulting output for both simulations was as expected, and the simulation time sets standards. Both test case models needed less than a second to be solved. The tested program was of version 0.27 and can be used under the terms of the [LGPL](#page-135-8) (provided that a valid license for Mathematica is obtained).

5.3.3 In Comparison to [4DiCeS](#page-134-0)

As with all the other ten simulation applications, [4DiCeS](#page-134-0) was critically checked for all the given test criteria (see Section [3.2.4\)](#page-33-0). It has to be said that [4DiCeS](#page-134-0) is not able to compete with other applications in single disciplines. Other simulation applications were implemented under the precondition to solve one or a few very specific model scenarios.

In contrast to this, [4DiCeS](#page-134-0) was implemented as an open tool to solve multiple scenarios correctly but not fully optimized. Therefore [4DiCeS](#page-134-0) is a trade-off between features as well as functionality on the one hand, and on the other hand optimized computational demands. The feature list is not fully comparable. [4DiCeS](#page-134-0) provides some functionality that none of the other applications can compete with. Such features are the simulation with algorithmic concurrency

and the open interface design. Contrarily [4DiCeS](#page-134-0) is missing certain functionality other tools have such as a graphical model editor and a data output interface.

Compared by the test criteria and the test models specified in Section [5.3.1,](#page-99-0) [4DiCeS](#page-134-0) has initial implementations for both a [GUI](#page-135-2) and [CLI](#page-134-5). It was shown that the user interface is capable of applying other user interaction scenarios combined with other programming languages through the implementation of a [CAVE](#page-134-6) front-end.

Then the question of the ease of use is a hard task for the programmer of exactly this application. Of course the intention was to have a clearly arranged user interface with as much assistance to the user as possible. In fact first user reactions to the system will show if further any improvement is necessary. For now it is reasonable to say that [4DiCeS](#page-134-0) clearly ranges among the tools that were indicated to be for experienced users. Admittedly [MCell](#page-135-13) is the only tool that was defined to be hard to handle since it lacks a [GUI](#page-135-2).

[4DiCeS](#page-134-0) was able to simulate both test cases correctly, even in [3D](#page-134-8). However first user reactions are needed in order to make further statements. This holds true for two important reasons. Firstly the test models were chosen by the comparator. This is the complete opposite of a "double-blind study" known from pharmaceutical research. Secondly, many mistakes happen unobserved by application architects. Here only bug testing by people, other than the actual programmer, can unearth possible inconsistencies and problems. Nonetheless [4DiCeS](#page-134-0) ranged in times for both test cases of a [1D](#page-134-22) simulation within two seconds to five minutes and within a complete day for the [3D](#page-134-8) counterparts.

5.4 Related Work

Based on the comparison (see Section [5.3\)](#page-98-0), the following findings can be summarized for the ten compared programs. The usability of an application is somewhat critical from the user's perspective. For example [Gepasi](#page-135-12) and [BIOCHAM](#page-134-15) are easy to use, whereas a program such as xCellerator requires further programming experience with Mathematica.

Additionally the lack of standards and interfaces between tools becomes apparent. For example the support for external file formats, such as [SBML](#page-136-8), will become more important in the future as the amount of available data increases - "tower of Babel" problem. An example is SmartCell, which uses [SBML](#page-136-8) in combination with its own [XML](#page-136-9) file format in such a way that an import of plain [SBML](#page-136-8) files to the application is no longer possible. Also it was often not feasible to save a model by one application as an [SBML](#page-136-8) file and reload it with another. This was the case for [BIOCHAM](#page-134-15) and [Copasi](#page-134-18) in both directions.

Then a sufficient documentation and transparency of implementation details are fundamental principles of scientific practice. Here e.g. SmartCell is lacking further information regarding the algorithms it applies. Simulation tools also differ in their ability to utilize external triggers.

Another challenging area is the computational parameter estimation for models. Presently, \dot{a} priori information of the parameters may be unavailable and the parameter values are adjusted either with some semi-automatic methods or by manually varying them. The two tools [Gepasi](#page-135-12) and [Copasi](#page-134-18), provide methods for computational parameter estimation already.
Conclusions

A well structured and extensible platform for systems biology was developed by combining a set of functional components. The already implemented modules for the [4DiCeS](#page-134-0) application allow for simulation of preliminary applications of biochemical models.

Figure 6.1: Lego drawing of a generalized nucleus. An illustration of the modeling of the nucleus by Lego blocks. The blocks can be compared to the [4DiCeS](#page-134-0) [VE](#page-136-0)s that form a complex structure of cellular compartments. In contrast to Lego-blocks however, the [VE](#page-136-0) can be adjusted in size. Thus, the granularity of the system can be adapted to the scientific requirements at hand.

Systems biology poses new challenges for visualization as the data types often contain [4D](#page-134-1) information. The representation of such data is currently unresolved. This work provides an approach to easily handle multi-dimensional data. The presented format could become a standard for [3D](#page-134-2) modeling environments. In respect thereof this approach works similar to the construction of an object from a box of Lego^{[1](#page-108-0)} blocks, in which Lego blocks are exchanged against [VE](#page-136-0)s (see Figure [6.1\)](#page-108-1). Also the general system's design can be seen such that each part of the [4DiCeS](#page-134-0) application is exchangeable but will not work alone.

¹The Lego Group: <http://www.lego.com/>

6.1 Design Decisions

With [4DiCeS](#page-134-0) a modular approach was chosen to allow for the integration of diverse algorithms and file formats. With [4DiCeS](#page-134-0) models can be composed on different scales and complexity. It allows for the integration of numerous different algorithms to perform on either single particles, particle concentrations, or distinct [VE](#page-136-0)s. More importantly it is also possible to mix different algorithmic approaches within one simulation space.

In this work the [4DiCeS](#page-134-0) framework was developed for modeling and simulation of signal transduction networks in [3D](#page-134-2). The implementation of this project includes a modular design of the [4DiCeS](#page-134-0) kernel and various interfaces as described below:

• [4DiCeS](#page-134-0)-Kernel:

The kernel of the [4DiCeS](#page-134-0) framework is the core of the project. It is responsible for securely interfacing with other modules. Furthermore it provides a platform for whole cell modeling as a result of its underlying data structure. Here [3D](#page-134-2) geometry, reactions, and diffusion are supported.

• User Interface:

The user interface facilitates the adaption of tailored user-interaction solutions to the [4DiCeS](#page-134-0) system. Both [GUI](#page-135-0) as well as [CLI](#page-134-3) approaches are supported.

• Model Parsing Interface:

The input parsing interfaces provide the capability to extend the [4DiCeS](#page-134-0) framework to several modeling input description languages. Thus there is the potential to include models from various data repositories.

• Reaction Interface:

The reaction interface allows for the integration of reaction algorithms. These plug-ins provide the actual modular simulation functionality to the [4DiCeS](#page-134-0) system beside the diffusion interface.

• Diffusion Interface:

The diffusion interface assists the combination of both known and new diffusion techniques. It is essential for simulations in [3D](#page-134-2) to have the ability of translocating particles.

Each interface has its own API, including other programming language extensions facilitating the implementation and integration of native code as well as other language plug-ins. The complete [API](#page-134-4)s offer a purpose-specific control over the [4DiCeS](#page-134-0) platform.

[4DiCeS](#page-134-0) aims to serve as a universal simulation framework, which can integrate any set of different simulation algorithms, including differential-equation-based models, diffusion-reaction, stochastic algorithms [\(Gillespie,](#page-148-0) [1976,](#page-148-0) [1977,](#page-148-1) [2001;](#page-148-2) [Gibson and Bruck,](#page-148-3) [2000\)](#page-148-3) as well as many from [CA](#page-134-5) [\(Wurthner et al.,](#page-158-0) [2000\)](#page-158-0) to GMA/S-Systems (Hernández-Bermejo et al., [2000\)](#page-149-0).

The tool is designed to conduct efficient cellular signaling simulations on a cell model consisting of segments with their specialization on different simulation tasks.

Successful attempts were also made to introduce the visualization environment to a [CAVE](#page-134-6) (see Section [4.2.4\)](#page-78-0) for virtual reality. The user interface from the [4DiCeS](#page-134-0) package was used for displaying the simulation results onto the four stereo-enabled screens of a [CAVE](#page-134-6).

Finally the usefulness and utility of this framework was shown for two sample diffusion and reaction applications in comparison to other related tools.

6.2

In comparison to other existing systems, the [4DiCeS](#page-134-0) workbench has its benefits in its modularity and extendability. Other simulation tools are either limited by mathematical constraints or do not support [3D](#page-134-2) structures (see Section [5.3.2\)](#page-99-0). The [4DiCeS](#page-134-0) framework has no limits in this respect. It furthermore provides the potential to include any type of diffusion and/or reaction algorithm. The integration of existing simulation tools is also possible via the [API](#page-134-4) definitions as well.

Only a few of the existing cell simulators conceptualize the division of a cell into its representative compartments, organelles. Most of the simulation tools allowing for such an approach do not offer a "real" [3D](#page-134-2) analogous representation of exactly such segments. They do however make use of encircling parts of the model to define confined compartments. The only cell simulation tools that currently support [3D](#page-134-2) geometry (see Figure [6.2\)](#page-111-0), besides the one developed here, are [MCell](#page-135-1) (see Section [3.2.2.3\)](#page-31-0), SmartCell (see Section [3.2.2.4\)](#page-31-1), and the VirtualCell (see Section [3.2.1.4\)](#page-29-0).

Figure 6.2: [3D](#page-134-2) visualizations of cell simulators. [\(a\)](#page-111-1) shows the communication between neurons in the brain occurs at synapses with [MCell](#page-135-1) [\(Cog](#page-145-0)[gan et al.,](#page-145-0) [2005\)](#page-145-0), [\(b\)](#page-111-2) displays a very simple model of a cell with its nucleus created by SmartCell, and [\(c\)](#page-111-3) presents the [3D](#page-134-2) surface visualization viewer of the VirtualCell project showing a neuronal cell filled with calcium indicator collected by confocal microscopy at University of Connecticut Health Center Farmington, USA, by Ion Moraru.

Here VirtualCell uses a purely deterministic approach for [2D](#page-134-7) as well as [3D](#page-134-2) modeling and simulation. Within [MCell](#page-135-1) every compartment is modeled using [3D](#page-134-2) vector schemes [\(Casanova et al.,](#page-145-1) [2004\)](#page-145-1). This modeling approach comes with [3D](#page-134-2) representations that approximate reality but at the costs of being computationally high demanding. Visualizing every particle and its trajectory over time increases complexity and resource requirements dramatically. Lastly SmartCell utilizes a very similar approach as [4DiCeS](#page-134-0) does. The reaction space is a grid of [VE](#page-136-0)s that can define separate compartments.

With this approach the computational time does not necessarily depend on the number of particles and their trajectories but on the number of defined [VE](#page-136-0)s. If either the [VE](#page-136-0)s are scaled to the size of particles or the internal computation of [VE](#page-136-0)s is handled by calculating trajectories, then the simulation might run as slow as [MCell](#page-135-1).

The grid layout offers another benefit. It is possible to run different reaction and diffusion algorithms on the model. [MCell](#page-135-1) is limited to the calculation of trajectories of single particles and VirtualCell can only handle differential equations. With SmartCell and [4DiCeS](#page-134-0), arbitrary algorithms can be chosen and executed on the same model all the time. This holds true for SmartCell in theory only since it is lacking a well-defined programming interface for plugable algorithms. In contrast [4DiCeS](#page-134-0) offers these advantageous interfaces and even goes a step further in providing a framework where not only algorithms can be plugged to the system, but the modeling, visualization, and analysis engines may be also exchanged. Although [4DiCeS](#page-134-0) is delivered with a preliminary set of algorithms and tools, the programming languages' [API](#page-134-4)s feature a good extendability to future needs. [Copasi](#page-134-8) (see Section [3.2.2.2\)](#page-31-2) as the only competitor offers a similar extendability, but then lacks the [3D](#page-134-2) support.

	Accuracy	Concurrency	Designer	Exchange	Extendability	Features	GUI	Methodology	Performance	$\operatorname{Scripting}$	S paciality	$\label{eq:bsability} \textbf{Usability}$	Score
4DiCeS			\bullet	\bullet		\bullet	٠		\bullet	D		\bullet	9.5
Copasi		∩	\bullet		∩					\bullet	\bullet	\bullet	8.0
E-Cell		O	◑		\bullet			∩	\bullet	Ð	\bullet		8.0
BioNetS		∩	\bullet	\bullet	\bullet	\bullet	٠			\bigcirc	∩		7.0
VirtualCell		∩			∩			∩	\bullet	◯		\bullet	7.0
Gepasi		O	\bullet	\bullet	∩			◯		\bigcirc	∩		6.0
SmartCell	\bullet	∩	◑	\bullet	∩	◯	٠		\bullet	◯			6.0
x Cellerator		∩	\bullet	\bullet	\bullet	\bullet	\bullet		\bullet	\bullet	∩	\bullet	6.0
StochSim		O	\bullet	\bullet	∩	◯	٠	\bullet	◯	n	\bullet	\bullet	5.5
BIOCHAM	\bullet	∩	∩	\bullet	∩	O		∩	\bullet		()	\bullet	4.0
MCell			∩	∩	∩	\bullet	C	\bullet	∩			∩	4.0

Table 6.1: [4DiCeS](#page-134-0) comparison. Equal to Table [3.4](#page-35-0) filled $\left(\bullet \right)$ circles are generally superior to their half-full $\left(\mathbb{O} \right)$ counterparts. Empty $\left(\bigcirc \right)$ circles indicate the need for improvement or a total absence. In methodology hybrid approaches are considered most sufficient. Designers are either defined as textual $(\mathbf{0})$, graphical $(\mathbf{0})$, both $(\mathbf{0})$, or none (\circ) at all. Additionally, the score was calculated and printed to this table also. Each score is the sum of circle weights of the respective application. The weights are defined with empty $= 0$, half $= 0.5$, and $full = 1$ circles. All comparison criteria were assessed equally. The applications are then sorted by their individual score in a top-down order. If two or more application have an equal score then they are sorted in alphabetical order.

Table [6.1](#page-112-0) provides an overview of how [4DiCeS](#page-134-0) is rated against the 12 comparison criteria from Section [3.4.](#page-35-0) If the same weighting scheme with zero (0), a half (0.5), and one (1) is applied then [4DiCeS](#page-134-0) exceeds all other related application with a comparison value of 9.5. The two best tools ([Copasi](#page-134-8) and E-Cell) only reached a value of eight.

In summary it can be said that [4DiCeS](#page-134-0) might not outperform single other simulation applications in their specialized disciplines but rather provides a tool that incudes the ability of allowing all the different models run on one single simulation framework only. The modular design and the well-defined interfaces make it easy for software developers to extend [4DiCeS](#page-134-0) to their needs. The use of cross-platform libraries makes the application executable on almost all recent operation systems. Also the [4DiCeS](#page-134-0) system is not limited to client applications alone but may be used as a server-based system as well.

6.3 Challenges and Accomplishment

In recent years whole cell simulation has been declared as a new scientific goal [\(Tomita,](#page-157-0) [2001\)](#page-157-0). The E-Cell project (see Section [3.2.1.2\)](#page-29-1) with its deterministic approach is only one such example. Although still not within reach, many approaches exist that use deterministic as well as stochastic modeling methods for simulating at least biochemical processes within a cell.

Stochastic simulators consume more computational power than deterministic simulators and are thus generally used for smaller models. [Endy et al.](#page-146-0) [\(2000\)](#page-146-0) simulated a bacteriophage by stochastic models. The stochastic simulation algorithms by [Gillespie](#page-148-0) [\(1976,](#page-148-0) [1977\)](#page-148-1) were further optimized [\(Gibson,](#page-148-4) [2000\)](#page-148-4) such that "whole cell simulation" goal comes within reach. To allow for an appreciation of how long it might take for current programs and computers to simulate such a challenge, the following sections consider the complexity of whole cell models and their calculation time.

6.3.1 Model Complexity

For the simulation of whole cell models, the complexity of such an approach should be first estimated. For example the base limit for the particle numbers involved can be obtained from the genome of an arbitrary organism. As for

Escherichia coli (an intestinal bacteria) the total number is currently counted near 4, 500 genes [\(Karp et al.,](#page-150-0) [2000,](#page-150-0) [2002\)](#page-150-1). The cytoplasm of E. coli includes according to an estimation [\(Goodsell,](#page-148-5) [1997\)](#page-148-5) around 22, 500 proteins, 15, 000 ribosomes, over 170, 000 ribonucleic acid ([RNA](#page-136-2)) molecules, 15, 000, 000 small organic molecules as nucleotides, amino acids, sugar, and others, as well as 25, 000, 000 ions. Then, about 70% of the bacteria's cell volume consist of water.

During one full cell cycle of E. coli (\sim 50 sec), the number of biochemical reactions executed was estimated [\(Endy and Brent,](#page-146-1) [2001\)](#page-146-1) to be somewhere between 10^{14} and 10^{16} . The complexity of other cell types is about a factor thousand smaller for the smallest known cell types (mycoplasms) and about a factor of thousand larger for typical plant and animal cells [\(Schwehm,](#page-156-0) [2001\)](#page-156-0). Also the relation of compartments to the cellular volume has to be taken into account if one considers higher organized cells (see Table [6.2\)](#page-114-0).

Table 6.2: Membranous compartments. The table shows the relative volumes occupied by the major membranous compartments (organelles) in a liver cell (Hepatocyte). An organelle is a specialized subunit within a cell that has a specific function, and is separately enclosed within its own lipid membrane. The first column holds the names of the intracellular compartments. The second and third columns display the percentage of due to the total cell volume and the approximate number of compartments per cell. Adapted from [Alberts et al.](#page-142-0) [\(2003\)](#page-142-0).

6.3.2 Performance Evaluation

Analyzing the given numbers, the question arises how far scientific research is away from realizing whole cell simulation. A very optimistic estimation implies an appropriate whole cell model with 10^{14} biochemical reactions for a single cell cycle of E. coli [\(Endy and Brent,](#page-146-1) [2001\)](#page-146-1). With a mean of 250, 000 processed reactions per second, [\(Schwehm,](#page-156-0) [2001\)](#page-156-0) a stochastic whole cell simulation of only one E. coli cell cycle would take 4 ± 10^8 sec or nearly 13 years for a single computer processing unit.

Thus taking into account a factor of thousand for typical eukaryotic cells the simulation of whole cells is even more complex. Here either the modularization of the problem domain or much faster algorithms as well as heuristics are required. Combined heterogenous simulation methods, as done with [4DiCeS](#page-134-0), is going to complement such approaches. Unfortunately future enhancements in computer power or the utilization of highly distributed computing will again feature a linear speedup. However considering that it once will be possible to simulate such a process within a few hours, the discrepancy to the actual 50 sec of a cell cycle is still far to great.

6.4 Outlook

Further extension as well as further development of this software could be directed towards incorporating additional specialized interfaces for (partly) automated pathway construction, a geometrical input tool for real cellular geometries from microscopy, model and simulation analysis, verification tools, and an own repository of models in addition to the collected kinetic data on molecules, reaction, diffusion, and pathways.

These interfaces should be implemented in a dynamically-linked manner as used throughout the existing components of the [4DiCeS](#page-134-0) framework. Module updates or supplements can be easily incorporated. The following paragraphs briefly illustrate some possible future extensions.

Distributed Computing: The simulation kernel could distribute its work onto heterogenous computer-networks. On a very low level, it is already possible to distribute the simulation process by using a screensaver-based approach of distributed computing – called Models@Home^{[2](#page-115-0)} [\(Krieger and Vriend,](#page-151-0) [2002\)](#page-151-0).

²Models@Home: <http://www.cmbi.kun.nl/models/>

Models@Home is building a network of idle computers while allowing each of them to work on small pieces of a scientifically demanding project. The world's largest distributed computing project, which is the Search for ExtraTerrestrial Intelligence ([SETI](#page-136-3)) [3](#page-116-0) , works on a very similar but much larger scale.

Model Construction Interface: A model construction interface could facilitate the construction of new or the extension of existing models for [4DiCeS](#page-134-0) users. Plug-ins could connect the simulation environment with common pathway and molecular databases (see Section [3.1\)](#page-24-0) to automatically retrieve information needed for the current model being developed. The construction kernel allows for the manual design of such models as well.

Geometry Input Interfaces: To allow for close-to-reality [3D](#page-134-2) cell structures, files from different microscopy techniques such as [CLSM](#page-134-10) imaging are planned to be importable. Additionally it would be possible to take time-dependant pictures of living cells in [4D](#page-134-1). This can then be used for verification in modeling and simulation. As a preliminary approach, it could be designed to have a few different input modules for various [CLSM](#page-134-10) picture stack formats from leading microscopy companies such as $Leica⁴$ $Leica⁴$ $Leica⁴$, Nikon^{[5](#page-116-2)}, Olympus^{[6](#page-116-3)}, TILL Photonics^{[7](#page-116-4)}, and $Zeiss⁸$ $Zeiss⁸$ $Zeiss⁸$.

Analysis Interface: Analysis is normally done through third party products such as spreadsheet analysis programs and graph drawing utilities. A future implementation for an analysis interface would not provide a full replacement package for these powerful programs with a possibility to analyze the current output of a given simulation in real time. Further massive analysis of the data has to be done externally but could then use such an interface for automated data transportation.

Programming Language Bindings: Further programming language bindings could enlarge the usability and acceptance of [4DiCeS](#page-134-0). Microsoft's .NET framework for example could easily open this tool to all .NET languages as for C_{\sharp} , VisualBasic, and Managed C++. The framework is also ported by the third-party Mono project^{[9](#page-116-6)} and can therefore be used on various other operation systems as well.

³Seti@Home: <http://setiathome.berkeley.edu/>

⁴Leica Microsystems: <http://www.leica-microsystems.com/>

⁵Nikon: <http://www.nikon.com/>

⁶Olympus: <http://www.olympus.com/>

⁷TILL Photonics: <http://www.till-photonics.com/> (now part of Agilent Technologies)

⁸Carl Zeiss: <http://www.zeiss.de/>

⁹Mono Project: <http://www.mono-project.com/>

[4DiCeS](#page-134-0) Repository: The database storage could provide future [4DiCeS](#page-134-0) users with data from past simulations and models. This is meant to keep the model construction time small and the reusability high. The simulation process delivers different kinds of data, and the simulator can always reuse a previous output as a new input.

Algorithms in Detail

This chapter gives a brief overview of the algorithms that were applied as first module implementations to [4DiCeS](#page-134-0). The next two sections explain stochastic approaches regarding reaction and diffusion processes and an [ODE](#page-136-4)-solver used for reaction equations.

A.1 Reaction Algorithms

Continuous biochemical rate equations do not accurately predict cellular reactions since they rely on bulk reactions that require the interactions of millions of particles. They are typically modeled as a set of coupled ordinary differential equations. In contrast dynamic [MC](#page-135-3) algorithms allow for a discrete and stochastic simulation of a system with few reactants, because each reaction is explicitly simulated.

It has to be noted that the reaction constants for dynamic [MC](#page-135-3) algorithms are not the traditional macroscopic or deterministic rate constants. Rather they are mesoscopic rate constants, which are only related to the macroscopic rate constants. Macroscopic rate constants depend on concentrations of particles, while mesoscopic reaction constants are based on numbers of particles.

For [4DiCeS](#page-134-0) a numeric [ODE](#page-136-4)-solver based on the Runge–Kutta 4^{th} order algorithm and three dynamic [MC](#page-135-3) methods were implemented.

A.1.1 Runge–Kutta Method

This explicit method numerically integrates [ODE](#page-136-4)s by using trial steps at the midpoint of an interval to cancel out lower-order error terms. Let an initial value problem be specified as follows

$$
\frac{dy_i}{dt} = f_i(t, y_1, ..., y_n)
$$
\n(A.1)

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with $t \in [a, b]$ and the initial values $y_i(a) = y_{a_i}, 1 \le i \le n$. Then, the fourthorder formula for this problem is given by

$$
y_{n+1} = y_n + h \times K \tag{A.2}
$$

where

$$
K = \frac{k_1 + 2k_2 + 2k_3 + k_4}{6} \tag{A.3}
$$

with

$$
k_1 = f(t_n, y_n), \tag{A.4}
$$

$$
k_2 = f(t_n + \frac{1}{2}h, y_n + \frac{1}{2}k_1), \tag{A.5}
$$

$$
k_3 = f(t_n + \frac{1}{2}h, y_n + \frac{1}{2}k_2), \text{ and} \t(A.6)
$$

$$
k_4 = f(t_n + h, y_n + k_3) \tag{A.7}
$$

[\(Runge,](#page-155-0) [1895;](#page-155-0) [Kutta,](#page-151-1) [1901\)](#page-151-1). The next value y_{n+1} is determined by the present y_n , an estimated slope K, and the product of the size of the interval h. K is a weighted average of slopes. Here k_1 is the slope at the start of h. k_2 is the slope at the midpoint of the interval, using slope k_1 to define the value of y at the point $tn + \frac{h}{2}$ $\frac{h}{2}$ using Euler's method. Again k_3 is the slope at the midpoint however using the slope k_2 to determine the y-value. k_4 is the slope at the end of h, with its y-value provided using k_3 . A greater weight is given to the two slopes at the midpoint.

The total accumulated error is on the order of $h⁴$, while the error per step has order h^5 . The Runge–Kutta Method is reasonably robust as well as simple and can be generally used for numerical solution of differential equations if combined with a flexible step-size method. For a given set of N differential equations and M intermediate steps this algorithm takes time proportional to $N \times M$ to solve the equations for t_{n+1} [\(Press et al.,](#page-155-1) [1992\)](#page-155-1).

A.1.2 Dynamic [MC](#page-135-3) Algorithms

Dynamic [MC](#page-135-3) algorithms can be used for modeling the dynamic behaviors of particles by comparing the rates of individual steps by random numbers. Such

algorithms generate statistically correct trajectories of stochastic equations. The 'Direct Method' and the 'First-Reaction Method', both developed and published by [Gillespie](#page-148-0) [\(1976,](#page-148-0) [1977\)](#page-148-1), are algorithms to simulate chemical or biochemical reaction systems accurately as well as efficiently. These two algorithms are computationally demanding. Therefore many adaptations and modifications exist. These include the 'Next-Reaction Method' [\(Gibson and](#page-148-3) [Bruck,](#page-148-3) [2000\)](#page-148-3), τ -leaping, as well as hybrid techniques where abundant reactants are modeled with deterministic behavior.

Both Gillespie algorithms as well as the Next-Reaction Method by Gibson and Bruck were implemented in [4DiCeS](#page-134-0). In the following a well-stirred system of N chemical species $S_1, ..., S_N$ is undergoing M chemical reactions $R_1, ..., R_M$. The current state of the system is specified by the vector $\mu = (\mu_1, ..., \mu_N)$, where μ_i is the current number of S_i molecules in the system. Each reaction channel R_j is characterized by its propensity function $a_j(\mu)$ and its state-change vector $v_j = (v_{1j},...,v_{Nj})$. Here $a_j(\mu)d\tau$ gives the probability that the system will experience an R_j reaction in the next infinitesimal time $d\tau$, and v_{ij} is the change in the number of S_i molecules caused by one R_j reaction.

A.1.2.1 Direct Method

Gillespie's Direct Method calculates explicitly which reaction occurs next and when it appears. This is achieved by specifying the probability density $P(\mu, \tau)$ that the next reaction is μ that occurs at time τ . The algorithm is direct in the sense that it generates μ and τ instantaneously. It can be shown that

$$
P(\mu, \tau) d\tau = a_{\mu} e^{\left(-\tau \sum_j a_j\right)} d\tau \tag{A.8}
$$

Integrating $P(\mu, \tau)$, for all τ , $0 \leq \tau \leq \infty$ results in

$$
Pr(\mu) = \int P(\mu, \tau)dt = \frac{a_{\mu}}{\sum_{j} a_{j}}.
$$
 (A.9)

Then, summing $P(\mu, \tau)$ over all μ gives

$$
P(\mu, \tau) d\tau = \left(\sum_j a_j\right) e^{\left(-\tau \sum_j a_j\right)} d\tau.
$$
 (A.10)

These two distributions lead to Gillespie's direct algorithm:

- 1. Initialization: Set initial numbers of particles and $t_0 = 0$.
- 2. Calculate the propensity functions a_i for all i.
- 3. Choose μ according to the distribution in Eq[.A.9.](#page-120-0)
- 4. Choose τ by an exponential with parameter $\sum_j a_j$ as in Eq[.A.10.](#page-120-1)
- 5. Change the number of particles according to μ . Set $t_{n+1} = t_n + \tau$.
- 6. Go to Step 2 unless the simulation time exceeded or a reactant's number is zero.

This algorithm uses two random numbers per iteration, takes time proportional to M to update the a_i s, and takes time proportional to M to calculate $\sum_j a_j$ as well as to generate a random number according to the distribution in Eq[.A.9](#page-120-0) [\(Gillespie,](#page-148-0) [1976\)](#page-148-0).

A.1.2.2 First-Reaction Method

Gillespie's First-Reaction Method generates for each reaction a putative time τ_{μ} at which the reaction μ occurs. Thereafter, the reaction μ^* with the smallest τ_{μ}^* is chosen and executed. Formally, the algorithm is as follows:

- 1. Initialization: Set initial numbers of particles and $t_0 = 0$.
- 2. Calculate the propensity functions a_i for all i.
- 3. For each i, generate a putative time τ by an exponential distribution with parameter a_i .
- 4. Let μ be the reaction whose putative time τ_{μ} is least.
- 5. Let τ be τ_{μ} .
- 6. Change the number of particles according to μ . Set $t_{n+1} = t_n + \tau$.
- 7. Go to Step 2 unless the simulation time exceeded or a reactant's number is zero.

For a given set of M reactions this algorithm uses I random numbers per iteration, takes time proportional to I to update the a_i s, and takes time proportional to I to identify the smallest τ_{μ} [\(Gillespie,](#page-148-1) [1977\)](#page-148-1).

A.1.2.3 Next-Reaction Method

As a modification of Gillespie's two methods the Next-Reaction Method is based on:

- Store τ_i , not just a_i .
- Recalculate a_i only if it changes.
- Re-use τ_i s where appropriate.
- Switch from relative time between reactions to absolute time.
- Store a_i s and τ_i s so that updating will be very efficient.

Here a *dependency graph* is introduced to update the minimum number of a_i s and an *indexed priority queue* stores all a_i s and τ_i . These modifications lead to the algorithm:

- 1. Initialization:
	- **a** Set initial numbers of particles and set $t_0 = 0$. Generate a dependency graph G.
	- **b** Calculate the propensity functions a_i for all i.
	- **c** For each i, generate a putative time τ according to an exponential distribution with $parameter a_i$.
	- **d** Store the τ_i values in an indexed priority queue P.
- 2. Let μ be the reaction whose putative time τ_{μ} stored in P is least.
- 3. Let τ be τ_{μ} .
- 4. Change the number of particles according to μ . Set $t_{n+1} = t_n + \tau$.
- 5. For each edge (μ, α) in the dependency graph G,
	- **a** Update a_{α} .
	- **b** If $\alpha \neq \mu$, set $\tau_{\alpha} = (a_{\alpha}, old/a_{\alpha}, new)(\tau_{\alpha} t) + t$.
	- **c** If $\alpha = \mu$, generate a random number ρ according to an exponential distribution with parameter a_{μ} . Set $\tau_{\alpha} = \rho + t$.
	- **d** Replace the old τ_{α} value in P with the new value.
- 6. Go to Step 2 unless the simulation time exceeded or a reactant's number is zero.

The total number of operations per iteration is at most $c_{2,3,4,5a,6} + c_{5b}(k-1)$ + $c_{5c} + c_{5d}(k)(2 \log M)$, where each c is a machine specific constant, k is the times of executing c_5a , and M is the number of reactions. This results in $O(\log M)$ [\(Gibson and Bruck,](#page-148-3) [2000\)](#page-148-3).

A.2 Diffusion Algorithm

The mathematics of Brownian motion, as diffusion processes are also called, is often studied using simple models. The simplest is a random-walk where equal sized steps can be taken in any direction. For purposeful movement the start-to-end displacement increases at a rate proportional to time t leading to a Gaussian distribution of particles. A Gaussian random-walk is self-similar and such processes were originally discussed by the mathematician Paul Lévy [\(1948\)](#page-152-0). Random-walks with self-similar dynamics and power-law scaling are therefore called discrete Lévy processes.

The approach used here for simulating the diffusion of particles either bound to membranes or freely distributing is a straight forward random-walk implementation.

A.2.1 Discrete Lévy Process

Let (X_n) , $n \in \mathbb{N}_0$ be a stochastic process. Applies $|X_{n_1+h}-X_{n_1}| \sim |X_{n_2+h}-X_{n_2}|$ for all n_1 , n_2 , and $h \in \mathbb{N}_0$ then X is a process of stationary increase with the given representation

$$
X_{n+1} = X_n + \sum_{i=1}^{t} Y_i.
$$
 (A.11)

With a mean quadratic shift of Ω using the diffusion coefficient $D = \frac{K_B T}{f}$ with the Bolzmann constant K_B , a coefficient of friction f and temperature T, it is possible to change the position $X_n = (x_{n1}, x_{n2}, x_{n3})$ of a particle over two random angles α and β in [3D](#page-134-2). For a lateral diffusion in [2D](#page-134-7) with $X_n = (x_{n1}, x_{n2})$ only one angle α is required. For an ideally sphere shaped particle $f = 6\pi\eta R_S$ (Einstein-Stokes Equation) [\(Einstein,](#page-146-2) [1956\)](#page-146-2), where η is the viscosity of water and R_S is the Stokes-radius, a [2D](#page-134-7) diffusion is described by

$$
X_{n+1} = X_n + \Omega\left(\begin{array}{c} \cos(\alpha) \\ \sin(\alpha) \end{array}\right) \tag{A.12}
$$

where $\Omega = \sqrt{4Dn}$ and $-\pi \leq \alpha \leq \pi$. The [3D](#page-134-2) diffusion is then described by

$$
X_{n+1} = X_n + \Omega \begin{pmatrix} \cos(\alpha)\cos(\beta) \\ \sin(\alpha)\cos(\beta) \\ \sin(\beta) \end{pmatrix}
$$
 (A.13)

where $\Omega = \sqrt{6Dn}$, $-\frac{\pi}{2} \leq \alpha \leq \frac{\pi}{2}$ $\frac{\pi}{2}$, and $-\pi \leq \beta \leq \pi$. The total number of operations n per iteration i is proportional to the numbers of particles p existing in the system.

Backus–Naur Form

The Backus–Naur Form ([BNF](#page-134-11)) is a meta-syntax used to express context-free grammars in a formal way. It was named after two pioneers in computer science, John Warner Backus and Peter Naur [\(Knuth,](#page-151-2) [1964,](#page-151-2) [2004\)](#page-151-3). [BNF](#page-134-11) is widely used as a notation for the grammars of communication protocols, computer programming languages, and instruction sets. A [BNF](#page-134-11) specification is a set of derivation rules written as

$$
\langle S \rangle ::= \langle E \rangle \tag{B.1}
$$

where $\langle S \rangle$ is a non-terminal symbol, and the $\langle E \rangle$ is an expression consisting of sequences of other symbols and/or a set of symbols. Sets are separated by a vertical bar $'$, indicating a choice. Symbols that never appear on a left side are called terminals [\(Backus,](#page-142-1) [1959;](#page-142-1) [Nauer,](#page-153-0) [1960\)](#page-153-0). There are many extensions and variants of the [BNF](#page-134-11), where the Extended [BNF](#page-134-11) ([EBNF](#page-135-4)) is a very common derivative. Often such specifications include some of the additional syntax rules stated in Table [B.1.](#page-126-0)

Rule	Description	
Alternative	Choices in a production are concatenated by a '.	
Grouping	Simple parenthesis enclose groups of symbols.	
Option	Options are enclosed in squared brackets.	
Repetition	Repeats are characterized by a following '*'.	
$0+$ Repeat	Zero or more time repeats are enclosed in curly brackets.	
$1+$ Repeat	One or more time repeats are followed by a $+$.	
Typeface	Terminals appear in bold and non-terminals in plain text.	

Table B.1: [BNF](#page-134-11) extensions. Today there exist many [BNF](#page-134-11) grammar variants. This table describes the most commonly found extensions used by such extended derivatives [\(Wirth,](#page-158-1) [1977\)](#page-158-1).

Unified Modeling Language

The Unified Modeling Language ([UML](#page-136-5)) is a general-purpose modeling and specification language for object-oriented software engineering. It was designed to specify, construct, visualize and document software-intensive systems and includes a standardized graphical notation for the creation of common concepts like use-cases, components, classes, generalization, aggregation, and behaviors. [UML](#page-136-5) is officially defined at the Object Management $Group¹$ $Group¹$ $Group¹$ by the UML metamodel.

It is noteworthy to differentiate between the [UML](#page-136-5) model and the set of [UML](#page-136-5) diagrams of a system. A diagram is only a partial graphical visualization of the model. Next to such diagrams the model may also contain textual documentation. In [UML](#page-136-5) 2.0 there exist 13 types of diagrams (see Table [C.1\)](#page-129-0), which can be grouped into the three categories of behavior, interaction, and structure. [UML](#page-136-5) does not restrict its element types to a certain diagram type. Actually every [UML](#page-136-5) element may appear on almost all types of diagrams [\(Scott,](#page-156-1) [2004\)](#page-156-1).

The following sections give details on the used [UML](#page-136-5) diagram types in this work.

C.1

An important part of the [UML](#page-136-5) is the possibility of drawing use-case diagrams. Normally use-cases are utilized during the analysis phase of a project to identify the functional parts of a system. However use-cases can be used to demonstrate the actual systems functionality. Since a graphical [UML](#page-136-5) use-case has only a very low information content it is often accompanied by its textual representation.

Generally [UML](#page-136-5) use-case diagrams separate a system into the respective usecases and actors. Actors represent the roles of users (i.e. human beings, other hardware, or software) of the system. The actors are external to the system described by the use-cases. They trigger the system and may receive output from it. To emphasize the externality of actors use-cases are framed into a box.

¹Object Management Group: <http://www.omg.org/>

Table C.1: [UML](#page-136-5) diagram types. This table describes the 13 types of diagrams currently defined by the [UML](#page-136-5) in brief. For a better structuring the diagrams are often categorized into the three types of behavioral, statical, and interaction deportment [\(Larman,](#page-151-4) [2001;](#page-151-4) [Scott,](#page-156-1) [2004\)](#page-156-1). The four diagram types used in this work are marked with an asterisk $($ *').

In [UML](#page-136-5) use-cases are represented by ovals and the actors are drawn as simple stick figures. The actors are connected with the use-cases by plain lines. Then the use-cases may have dashed relationship arrows of two kind. The first is a uses assignment, which can be compared to a function call or subroutine in programming languages. The other link-up is the extends assignment, which manipulates the original use-case to fit other purposes. Both relationships are placed on top of their arrows surrounded by doubled angle brackets (i.e. \ll uses and \ll extends \gg) [\(Larman,](#page-151-4) [2001;](#page-151-4) [Scott,](#page-156-1) [2004\)](#page-156-1).

C.2 Class Diagrams

[UML](#page-136-5) class diagrams have the purpose to depict the classes within a model. In an object-oriented application classes have attributes, methods, and relationships to other classes. The fundamental element of a class diagram is an icon that depicts two of these components into compartments of one class-rectangle. The topmost compartment contains the class name, the middle comprises all attributes and the bottom includes a list of methods. Often the bottom two compartments are omitted or do only present attributes and methods that are meaningful for the current diagram. If a class is abstract its name will be indicated in an italic typeface.

The relationship to other classes is represented by different kinds of lines. [UML](#page-136-5) distinguishes aggregations, associations, compositions, dependencies, inheritance, and interfaces. The inheritance relationship in [UML](#page-136-5) is depicted by a peculiar triangular arrowhead with a plain line. The arrowhead points to its base class connecting it to its derived class(es). The composition relationship depicts a strong form of an aggregation and is also called a by-value-relation. It is represented by a black diamond at the class that composes another class. The other class is pointed to by an ordinary arrowhead, if the relationship is only navigable in one direction. The weak form of an aggregation is denoted by an open diamond. An other form of a containment shows no diamond at all. This plain arrow is called an association. Both the aggregation and the association are often referred to by-reference-relations. Sometimes the relationship between two classes is very weak. In this case they are only implemented as method arguments. Such relationships are represented by dashed arrows and are called dependencies. An [UML](#page-136-5) way to display interfaces (abstract classes with no attributes at all) is the so-called "lollypop" notation. The interface has a line with an empty circle at its end. This circle is the point to which other classes draw their dependency arrows to [\(Larman,](#page-151-4) [2001;](#page-151-4) [Scott,](#page-156-1) [2004\)](#page-156-1).

C.3

A [UML](#page-136-5) sequence diagram emphasizes the sequence of messages between objects. Rectangles represent objects. The names of the objects are underlined to distinguish them from classes. Additionally the object name is separated from the class name by a colon. If objects do not have a particular object name then the colon precedes the class name without an object name in front of it.

Below of each object is a dashed line known as the "lifeline". Such lines define time axes of the diagram. By convention time proceeds in a downward direction. A variation could though be to tip a sequence diagram on its side so that time proceeds to the right. The lifelines depict how long the objects that they are connected to are in existence. If lifelines extend from the very top of the diagram to the very bottom then this implies that the objects portrayed in the diagram are in existence before the start of time of the diagram and remain in existence beyond the end of the diagram.

The arrows between lifelines represent messages being sent between objects. Sequence numbers are permitted but not necessary. White arrow rectangles indicate that the arrow terminates on a called activation. They show the duration of the execution of a method in response to a message. The methods implicitly return to their caller at the end of the activation. Such returns can be displayed by an unlabeled arrow that extends from the bottom of the activation back to the lifeline of the calling object. In the case of asynchronous messages the end of an activation does not imply a return.

Narrow rectangles on top of lifelines that enclose groups of messages define an iteration. The looping condition for such an iteration is then shown at the bottom of the rectangle.

The creation of objects is denoted by a message arrow that terminates on an object's rectangle. Deletion is likewise denoted by a message arrow that terminates on a capital 'X' at the end of the object's lifeline.

Incomplete (half) arrowheads denote asynchronous messages. An asynchronous message is a message that the receiving object executed the method in a separated thread. That threat could be in existence prior to the sending of the asynchronous message and just waiting for something to do. This gives sequence diagrams the power to display concurrent multi-threaded interactions. Wherever concurrency is present, race conditions are possible. Race conditions occur when a single thread or object receives messages from two competing sources. If not handled properly, the participating objects can get quite confused.

C.4 Component Diagrams

Components are qualified in [UML](#page-136-5) as reusable program code that implements well-defined duties. Therefore components are very similar to [UML](#page-136-5) packages. They combine thematically related elements to a set. As a general rule, this topical relationship is also reflected in similar characteristics of elements, as for their attributes, interfaces, and methods. Whereas packages represent a view of the contents, components highlight the software-technical aspects of elements' commonality. Particularly the structure of the program is of major importance.

Components establish a unity with one or more interfaces to the outside. Component diagrams visualize this graphically. They describe dependencies between software components and their interfaces. Here a component is displayed as a box. Additionally, the box can be marked with the phrase 'component' surrounded by doubled angle brackets (i.e. \ll component \gg) or by a component symbol. Then interfaces are connected through the "lollypop" notation. Component diagrams are also used to describe static dependencies between programs, as for example compiler-dependencies. Such dependencies between programs are illustrated by dashed arrows. The dependent components points at the independent component.

Components can contain additional elements, as for objects, other components, or nodes. For comprehension as regards content, components are often substantiated by other diagrams, as for class diagrams and use-case diagrams. For technical aspects of the implementation, deployment diagrams are utilized.

List of Abbreviations _____________

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