

Nr. 202

The Nucleus of Homogeneous Games
with Steps

by

Joachim Rosenmüller and Peter Sudhölter

April 1991



University of Bielefeld
4800 Bielefeld, Germany

SE 050
US B51
202

The Nucleus of Homogeneous Games with Steps

J. Rosenmüller and P. Sudhölter

Abstract

Homogeneous games were introduced by VON NEUMANN-MORGENSTERN [15] in the constant-sum case. PELEG [7,8] studied the kernel and the nucleus within this framework. However, for the general non-constant-sum case OSTMANN [6] invented the unique minimal representation, ROSENMÜLLER [10] gave a second characterization and SUDHÖLTER [14] discovered the "incidence vector". Based on these results PELEG-ROSENMÜLLER [9] treated several solution concepts for "games without steps". The present paper treats the case of games "with steps". It is shown that with a suitable version of a "truncated game" the nucleus of a game is essentially the one obtained by truncating behind the "largest step". As the truncated version has "no steps", the case "with steps" is reduced to the one "without steps" - which is treated in [9].

UB BIELEFELD

030/1559746+1



Section 0 Homogeneous games

The material of this paper is organized as follows. SECTION 0 serves as an introduction to the theory of homogeneous games and provides the necessary concepts and notations. SECTION 1 deals with certain families of representations of homogeneous games "with steps". These families put an increasing amount of weight at the players within the lexicographically first minimal winning coalition (Theorem 1.11). As a consequence, it turns out that in games with steps, the system of minimal winning coalitions cannot be (weakly) balanced (Corollary 2.6). This is of course important in context with the structure of the nucleolus; thus SECTION 2 discusses some simple properties of the nucleolus. However, Corollary 2.6 is not sufficient to explain the structure of the nucleolus of a homogeneous game "with steps". Therefore, SECTION 3 explains the "reduction theory" of the nucleolus. First, a truncation procedure is necessary. This, after some preliminary work, is described by Definition 3.6. Lemma 3.7 explains the nature of the truncated game. Finally, by Theorem 3.8 and Corollary 3.9 we collect the material available so far and prove that the nucleolus of a game with steps reduces to the one of the truncated version.

This section serves as an introduction to the theory of homogeneous (simple) games. All of the material presented may be found in the literature, see e.g. OSTMANN [6], ROSENMÜLLER [10], SUDHÖLTER [14].

Let $N = \{1, 2, 3, \dots\}$ denote the "universe of players". For the "grand coalition" we choose some "interval" $\Omega = [a, b] = \{i \in N \mid a \leq i \leq b\}$. $\underline{P} = \underline{P}(\Omega) = \underline{P}([a, b]) = \{S \mid S \subseteq \Omega\}$ is the system of *coalitions*. If

$$v : \underline{P} \rightarrow \mathbb{R}, \quad v(\emptyset) = 0$$

is a mapping on \underline{P} , then $(\Omega, \underline{P}, v)$ is a *game*; somewhat sloppily we refer to v as to "a game". v is *simple* if $v : \underline{P} \rightarrow \{0, 1\}$ holds true.

Unions of coalitions and players are written $S \cup i$ instead of $S \cup \{i\}$; $S + T$ and $S + i$ denotes disjoint unions. Similarly, $i < T$ denotes $i < j$ ($j \in T$) ($S, T \in \underline{P}$, $i, j \in \Omega$).

Given a simple game v , $\underline{W} = \underline{W}(v) = \{S \in \underline{P} \mid v(S) = 1\}$ is the system of *winning coalitions* while

$$\underline{W}^m = \underline{W}^m(v) = \{S \in \underline{W} \mid v(T) = 0 \text{ for } T \not\subseteq S\}$$

is the system of *minimal winning coalitions* ("min-win coalitions").

A vector $M = (M_i)_{i \in \Omega} \in \mathbb{R}_+^{\Omega}$ is tantamount to a function on \underline{P} via $M(S) = \sum_{i \in S} M_i$ ($S \in \underline{P}$) (thus, it is a non-simple "game") and hence called a "measure" (M is additive).

Games and in particular measures, may be restricted on subsets $T \subseteq \Omega$, the notation is $v \upharpoonright_T$ or $M \upharpoonright_T$; e.g.

$$v \upharpoonright_T(S) = v(T \cap S) \quad (S \in \underline{P})$$

$$v \upharpoonright_T(S) = v(S) \quad (S \in \underline{P}(T));$$

the version living on $\underline{P}(\Omega)$ and the one living on $\underline{P}(T)$ are not distinguished. We tolerate $v \upharpoonright_{\emptyset}$.

If M is a measure and $\lambda > 0$, then (M, λ) is a *representation* of v if

$$v(S) = \begin{cases} 1 & M(S) \geq \lambda \\ 0 & M(S) < \lambda \end{cases}$$

holds true, in this case we write $v = v_{\lambda}^M$. Of course, integer representations are of particular interest.

A measure M is said to be *homogeneous* w.r.t. $\lambda \in \mathbb{R}_+$ (written " M hom λ ") if, for any $T \in \underline{P}$ with $M(T) > \lambda$, there is $S \subseteq T$ with $M(S) = \lambda$.

A game v is *homogeneous* if there exists a representation (M, λ) with M hom λ and $v(\Omega) = 1$. (The definition is due to VON NEUMANN-MORGENSTERN [15].)

We assume all representable games to be *directed*, i.e., there exists a representation (M, λ) such that $i < j$ implies $M_i \geq M_j$ ($i, j \in \Omega$).

Thus, the "strong" or "large" players (the ones with big weight M_i) are first in enumeration (or "index"). In particular, if $I(S) = \max \{i | i \in S\}$ for $S \in \underline{P}$, then $I(S)$ is the "weakest", "smallest", or last player in S .

While players are ordered according to "size", coalitions are ordered lexicographically. In particular, the lex-max min-win coalition is the lexicographically first minimal winning coalition; in a homogeneous game with a homogeneous representation (M, λ) this coalition is sometimes denoted by $S^{(0)}$ or $S^{(\lambda)}$ (an interval with measure $M(S^{(\lambda)}) = \lambda$).

Players i and j are of the same type (written $i \sim j$ or $i \sim^v j$), if for all $S \subseteq \Omega - (i+j)$, $v(S+i) = v(S+j)$ holds true. A representation (M, λ) is symmetric if $i \sim j$ implies $M_i = M_j$ ($i, j \in \Omega$).

Player $i \in \Omega$ is a dummy if $v(S \cup i) = v(S)$ for all $S \in \underline{P}$. All dummy players are of the same type. Note that the game is assumed to be directed; thus, the definition of types induces a decomposition of Ω into intervals

$$\Omega = T_1 + \dots + T_r$$

of players of one type: $i \in T_\rho$ is also expressed as " i is of type ρ "; thus

$$R = \{1, \dots, r\}$$

denotes the set of types. If dummies are present, then r is "the dummy type" and, in a natural way, type ρ is "stronger" than type $\rho + 1$ ($\rho \in R - r$).

We shall refer to "dummy" as to a character that may or not be attached to a player. There are two further characters, "sum" and "step", which we are going to explain now.

To this end, fix a non-dummy player $i \in \Omega$

Among all min-win coalitions containing i , let $L^{(i)}$ be a one with minimal length, i.e.

$$|L^{(i)}| = \min \{|S| \mid S \ni i, S \in \underline{W}_m\}$$

Then

$$C^{(i)} := [|L^{(i)}| + 1, b]$$

is the domain of i and $M^{(i)} := M|_{C^{(i)}}$ is its satellite measure.

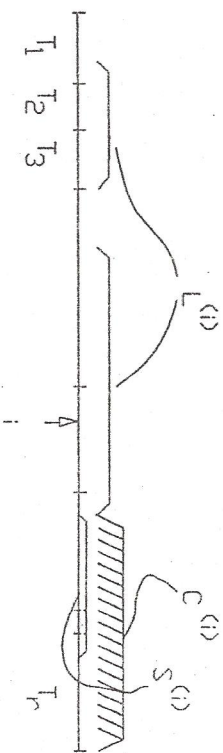


Fig. 1

Now, if $M^{(i)}(C^{(i)}) \geq M_i$, then i is a sum ("his character is sum"), since he may be replaced in a min-win coalition by a coalition of smaller players, his weight being the sum of the weights of the smaller players.

In this case, we call $v^{(i)} := v_{M_i}^{M^{(i)}}$ the satellite game of i (a homogeneous game). Also, $S^{(i)} = S_{M_i}^{M^{(i)}}$ is the coalition of its satellites; this is the lex-max min-win coalition of $v^{(i)}$.

Otherwise, if $M^{(i)}(C^{(i)}) < M_i$, then i is a step. In this case his ("pseudo") satellites are the members of his domain, i.e., we put $S^{(i)} := C^{(i)}$.

"Sum" and "step" are possible characters of a player - like dummy. From this, there results a further decomposition of Ω into the sets of characters

$$\Omega = \Sigma + \Pi + \Delta,$$

where $\Sigma = \Sigma(v) = \{i \in \Omega \mid i \text{ is a sum}\}$, $\Pi = \Pi(v) = \{\text{steps}\}$ and $\Delta = \Delta(v) = \{\text{dummies}\}$.

Σ as well as Δ may be empty while Π is not.

Remark 0.1: The following is well known (OSTMANN [6], ROSENTHALER [10], SUDHÖLTER [14]).

1. The smallest nondummy player is always a step. Its domain may be the empty set. If v is a constant-sum-game, then the smallest nondummy is the only step.

2. A homogeneous game has a unique minimal representation $(\bar{M}, \bar{\lambda})$ (e.g., in the sense that $(\bar{M}, \bar{\lambda})$ is integer and $\bar{M}(\Omega)$ is minimal), this representation is symmetric and attaches weight 0 to dummies.

3. A pair (M, λ) is a homogeneous representation of v , iff there exist real numbers $\Delta_i \geq 0$ with $\Delta_i > 0$ ($i \in I$), $\Delta_i = 0$ ($i \in \Sigma$) such that

$$\begin{aligned} M_i &= \Delta_i & (i \in \Delta) \\ M_i &= \Delta_i + M(S^{(i)}) & (i \in \Sigma \cup I) \end{aligned}$$

holds true. (Δ_i ($i \in \Omega$) is the "jump at i "). The unique minimal representation is obtained by putting $\Delta_i = 0$ ($i \in \Delta$), $\Delta_i = 1$ ($i \in I$).

4. Let $j \in \Sigma$ and let $i \in C^{(j)}$ be a nondummy w.r.t. $v^{(j)}$ (suitably, we write $i \notin \Delta^{(j)} := \Delta^{(v^{(j)})}$). Then i has domain, satellites etc. w.r.t. $v^{(j)}$; let $C^{(i,j)}$, $M^{(i,j)}$ denote his domain and satellite measure w.r.t. $v^{(j)}$. Then

$$\begin{aligned} C^{(i)} &= \bigcup \{C^{(i,j)} \mid j \in \Sigma, i \in C^{(j)}, i \notin \Delta^{(j)}\} \\ M^{(i)} &= \max \{M^{(i,j)} \mid j \in \Sigma, i \in C^{(j)}, i \notin \Delta^{(j)}\} \end{aligned}$$

holds true with an obvious interpretation of "max".

5. If $i \notin S^{(\lambda)}$, then $i \in \Sigma$ iff $i \in \Sigma^{(j)}$ for some $j \in S^{(\lambda)}$.

Also, $i \in \Delta$ iff $i \in \Delta^{(j)}$ for all $j \in S^{(\lambda)}$.

6. Let $j \in S^{(\lambda)}$ and let $\lambda^j = \lambda(S^j)$. Then $C^{(j)} = [1^j + 1, b]$. $S^{(\lambda)} \cap I$ is the coalition of "inevitable players" (i.e., those that are present in every min-win coalition). If all players in $S^{(\lambda)}$ are steps (inevitable), then v is the unanimous game of the members of $S^{(\lambda)}$ (with minimal representation $(\bar{M}, \bar{\lambda}) = (1, \dots, 1, 0, \dots, 0; \bar{\lambda})$).

Apart from the inevitable players, no further steps occur in $S^{(\lambda)}$.

7. In every homogeneous representation (M, λ) of v , sums of the same type have the same weight. Steps of the same type may have different weight, but then they appear or do not appear simultaneously ("as a block") in every min-win coalition.

Section 1 Monotone representation

Representations of homogeneous games are essentially defined by prescribing the "jumps" at the various steps. This section serves to study the consequences if the jumps are considered to be (positive) affine functions of a real parameter.

Lemma 1.1: ("small steps belong to its domain")

Let v be a homogeneous game and let $i < \tau$ be two players of different type. If $\tau \notin \Sigma$, then $\tau \in C(i)$.

Proof:

By induction on the number of types. If there is just one type, nothing has to be proved.

Otherwise, let $S(i)$ be the lex-max min-win coalition. τ cannot be an inevitable player (since i precedes him and is of different type). Also, τ cannot be one of the sums in $S(i)$. Hence, $\tau \notin S(i)$, and if $i \in S(i)$ the proof is done.

If $i \notin S(i)$, then consider, for every sum $j \in S(i)$ the satellite game $v(j)$. In at least one of these satellite games, i and τ are of different type. Since $\tau \notin \Sigma$, we have $\tau \notin \Sigma(j)$. By induction, $\tau \in C(i,j)$. Since

$$C(i) = \bigcup C(i,j),$$

our claim follows.

Definition 1.2:

Let $i, \tau \in \Omega$ and $i < \tau \in \Pi$. We shall say that τ is the next step following i if i and τ are of different type and there is no step in $[i+1, \tau-1]$.

Lemma 1.3:

Let (M, λ) be a homogeneous representation of v . Also, let $i \in \Omega$ and let $\tau \in \Omega$ be the next step following i . Denote the jump at i by Δ_i (thus, $\Delta_i = 0$ if $i \in \Sigma$).

Then there is $0 < c_i \in \mathbb{R}$ and $C = (c_i)_{i \in \Omega}$, $0 \leq c_i \in \mathbb{R}_0$ ($i \in [\tau+1, b]$) such that

$$(1) \quad M_i = \Delta_i + c_i M_1 + \sum_{j=1}^b c_j M_j = \Delta_i + c_i M_1 + CM.$$

Proof:

By the recursive property of M we have

$$M_i = \Delta_i + M(S(i)),$$

and by Lemma 1.1 we know that $\tau \in C(i)$ holds true.

Now, if τ is the first (largest) player in $S(i)$, then we are already done.

Otherwise, the first player in $S(i)$ has to be a sum and τ is in his domain (Lemma 1.1). Replacing his weight by a sum of weights of his satellites we observe players entering that are either sums preceding τ , τ himself and possibly smaller players - but never smaller players without τ (again by 1.1).

Proceeding this way, we replace the weight of all sums preceding τ by a sum of weights of smaller players. Thus, eventually the largest weight occurring in our sum is the one of τ - and necessarily with a positive coefficient. q.e.d.

Note that $CM = \sum_{j=1}^b c_j M_j$ is not necessarily the weight of a coalition (we cannot necessarily write " $M(T)$ " for this expression). For, although some M_j ($j < \tau$) is a sum of smaller players weights, these smaller players must not necessarily be available in $S(i)$ - hence the above mentioned process does not necessarily generate a coalition.

Nevertheless, the term CM as it involves weights of players smaller than τ looks rather similar to a sum of weights of smaller coalitions. Conveniently, we shall therefore write $\tau < C$ to indicate that $C = (c_i)_{i \in \Omega}$ is an integer vector with coordinates indexed by players smaller than τ . This will be important beginning with Lemma 1.8.

Similarly, we write $C \in C(i)$ to indicate that $C = (c_i)_{i \in C(i)} \in \mathbb{R}_0^{C(i)}$.

Definition 1.4:

Let v be a homogeneous game. A family of representations

$$(M(a), \lambda(a))_{a \in \mathbb{R}_{++}}$$

is said to be *affine*, if there are constants $A_i > 0$ ($i \in \Pi \cup \Delta$) and $B_i > 0$ ($i \in \Pi$) $B_i > 0$ ($i \in \Delta$) such that $M_i = M_i(a)$ satisfies

$$(2) \quad M_i = A_i a + B_i \quad (i \in \Delta)$$

$$(3) \quad M_i = A_i a + B_i + M^{(i)}(S^{(i)}) \quad (i \in \Pi)$$

$$(4) \quad M_i = M^{(i)}(S^{(i)}) \quad (i \in \Sigma)$$

that is, the jump at every step is a ("positive") affine function on \mathbb{R}_{++} .

Remark 1.5:

1. For a $e \in \mathbb{R}_{++}$, $(M(e), \lambda(e))$ is a homogeneous representation.

2. By induction it is easily seen that there are vectors $E_i, F_i \in \mathbb{R}_+^{\Omega}$ such that $F_i > 0$ and

$$(5) \quad M_i = M_i(a) = E_i a + F_i \quad (i \in \Omega),$$

hence

$$(6) \quad M(a)(S) = a E(S) + F(S) \quad (S \in \underline{E}).$$

Thus, an affine family of representations is equivalent to an affine mapping from \mathbb{R}_{++} into the homogeneous representations of v (regarded as a subset of \mathbb{R}^{Ω}).

For short, we shall write $(M(\cdot), \lambda(\cdot))$ to indicate an affine family of representations (an a.f.r.).

3. Satellite measures, lex-max-coalitions etc. do not depend on a . Therefore, it makes sense to state that, for some $i \in \Delta$,

$$(M^{(i)}(\cdot), g_i(\cdot))$$

is an a.f.r. of the satellite game $v^{(i)}$ etc.

Definition 1.6:

Let v be homogeneous. An a.f.r. $(M(\cdot), \lambda(\cdot))$ is said to be *monotone* if, for every $i \in \Pi$ of type ρ , the constants A_i and B_i as required by (3) satisfy

$$(7) \quad \frac{A_i}{B_i} \geq \frac{A_i'}{B_i'} \quad (i < i' \in \Pi \cap T_\rho)$$

$$(8) \quad \frac{A_i + B_i}{B_i + F_i} \geq \frac{A_i'}{B_i'} \quad (i, i' \in C^{(i)}).$$

A monotone $(M(\cdot), \lambda(\cdot))$ is *strictly monotone* "at $i \in \Pi$ " if (7) and (8) are strict inequalities and i is not of the type of the smallest step.

A monotone $(M(\cdot), \lambda(\cdot))$ is *strictly monotone* if it is strictly monotone at some $i \in \Pi$ which is *not* of the type of the smallest step.

Note 1.7:

For nonnegative reals a, b, c, d with $b, d > 0$ and $m, n \in \mathbb{N}$ it is clear that

$$(9) \quad \max\left(\frac{a}{b}, \frac{c}{d}\right) \geq \frac{ma + nc}{mb + nd} \geq \min\left(\frac{a}{b}, \frac{c}{d}\right)$$

Now, since (8) is required for all $j, l \in C^{(i)}$, it follows at once that for all coalitions $S, T \in C^{(i)}$

$$(10) \quad \frac{A_i + E(S)}{B_i + F(S)} \geq \frac{E(T)}{F(T)}.$$

In addition, if one inequality in (9) is strict, then so are all of them. Thus, if (8) is strict for all j, l in some $C^{(i)}$, then so is (10) for all $S, T \in C^{(i)}$.

The same holds true if instead of conditions $S, T \in C^{(i)}$ we use vectors $C = (c_l)_{l \in C^{(i)}}$ and $D = (d_l)_{l \in C^{(i)}}$ and $D = (d_l)_{l \in C^{(i)}}$ with coordinates $c_l, d_l \in \mathbb{N}_0$ - i.e., $C, D \in C^{(i)}$, a notation we have introduced with 1.3. The analogue to (10) is

$$(11) \quad \frac{A_i + CE}{B_i + CF} \geq \frac{DE}{DF}.$$

Monotone representations enjoy a certain monotonicity property: essentially the quotients $\frac{E_i}{F_i}$ are increasing with weights (i.e. from the right to the left).

Now, since the weight of a sum averages smaller weights we cannot expect that quotients increase while i moves through sums (from right to left) - in view of Note 1.7.

However, (8) ensures that there is a *significant* jump in the quotient *at every step* and that, thereafter, the quotient stays at least above all levels that have been attained by previous steps.

That is, from step to step the quotient increases significantly (from right to left) and in between it does not fall too rapidly.

Thus, we imagine the following picture

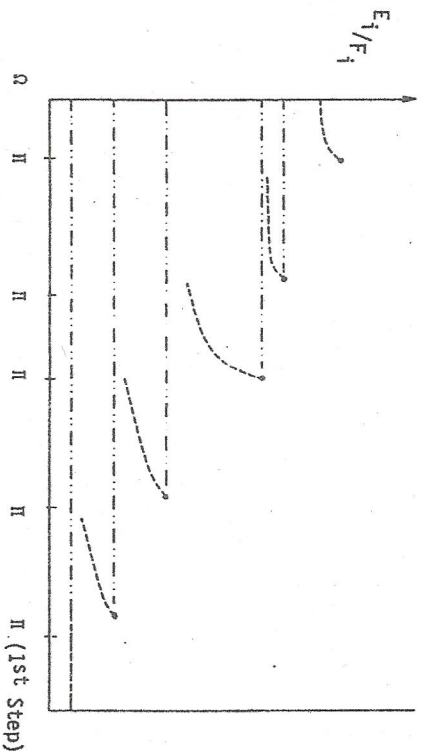


Fig. 2

The following exhibition serves to clear up this intuitive idea.

Lemma 1.8:

Let v be homogeneous and let $(M(\cdot), \lambda(\cdot))$ be a monotone a.f.r. Let $i \in \Pi$ and let $|j|$ be of smaller type. Then

$$(12) \quad \frac{A_i + E_j}{B_i + F_j} > \frac{E_j}{F_j}$$

In addition, if $(M(\cdot), \lambda(\cdot))$ is strictly monotone at $i \in \Pi$, then (12) is a strict inequality for all $|j|$ following τ which are of smaller type.

Proof:

If i is the smallest step, then there is nothing to prove.

If j and l follow behind the next step to the right of i , then (12) follows from (11) and Lemma 1.1, since the next step is in i 's domain and thus $|j|$ are in i 's domain $C^{(i)}$.

We shall consider the case that $|j| < \tau$ where τ is the next step to the right of i , and thus, j and l are sums (the remaining case will be clear, also the statement concerning strict inequalities is analogous).

Now writing $M = M(a) = aE + F$ we obtain by Lemma 1.3 with suitable $c, d \in M$, $C, D \in C^{(i)}$

$$(13) \quad \begin{aligned} M_j &= cM_\tau + CM \\ M_l &= dM_\tau + DM. \end{aligned}$$

(As $\tau < C, D$ it follows that $C, D \in C^{(i)}$ by 1.1).

A similar relation holds true for E_j, F_j , e.g.

$$\begin{aligned} E_j &= cE_\tau + CE \\ F_j &= cF_\tau + CF \end{aligned}$$

and the same for l, d, D .

Hence, in order to prove (9), we come up with

$$\frac{A_i + E_j}{B_i + F_j} = \frac{A_i + cE_\tau + CE}{B_i + cF_\tau + CF} > \frac{cE_\tau + DE}{cF_\tau + DF} = \frac{E_j}{F_j}$$

since $\tau \in C^{(i)}$ and CE, DE etc. are multiples of weights smaller than τ .

q.e.d.

Theorem 1.9:

Let v be homogeneous and let $(M(\cdot), \lambda(\cdot))$ be a monotone a.f.r.. Also, let $\tau \in \Pi$ and let $i \in \Omega$ be of smaller type.

Then, for all $i \leq \tau$

$$(14) \quad \frac{E_i}{F_i} > \frac{E_1}{F_1}$$

holds true. The inequality is strict if monotonicity is strict at τ .

Proof:

Assume first of all that i is a step of the same type as τ . Then i and τ have the same satellites and domains, thus, in view of (7), it follows that

$$\frac{E_i}{F_i} = \frac{A_i + E(C^{(i)})}{B_i + F(C^{(i)})} > \frac{A_\tau + E(C^\tau)}{B_\tau + F(C^\tau)} = \frac{E_\tau}{F_\tau} > \frac{E_1}{F_1},$$

the last inequality using Lemma 1.8.

Next, assume that τ is the next step following $i \in \Sigma$ (— the remaining cases will be omitted).

Choose c_1, C according to Lemma 1.3 such that $\tau < C$ and

$$M_i = c_1 M_1 + CM.$$

Then

$$\begin{aligned} E_i &= cE_\tau + CE \\ F_i &= cF_\tau + CF \end{aligned}$$

and

$$\frac{E_i}{F_i} = \frac{cE_\tau + CE}{cF_\tau + CF} = \frac{cA_\tau + cE(S^{(c)}) + CE}{cB_\tau + cF(S^{(c)}) + CF} > \frac{E_1}{F_1},$$

where again the last inequality is established by Lemma 1.3.

q.e.d.

Clearly, Theorem 1.9 resembles the statement suggested by Fig. 1.

Corollary 1.10:

If $(M(\cdot), \lambda(\cdot))$ is a monotone a.f.r. of a homogeneous game v , then

$$(15) \quad \frac{E_i}{F_i} > \frac{E(C^{(i)})}{F(C^{(i)})} \quad (i \in \Sigma \cup \Pi).$$

Proof:

If $i \in \Pi$, then $C^{(i)} = S^{(i)}$ and $E_i = A_i + E(S^{(i)})$, $F_i = B_i + F(S^{(i)})$, hence

$$(16) \quad \frac{E_i}{F_i} = \frac{A_i + E(S^{(i)})}{B_i + F(S^{(i)})} > \frac{E(S^{(i)})}{F(S^{(i)})}$$

follows from Definition 1.6 (formula (8)).

Now, let $i \in \Sigma$. Suppose first of all, that there are no steps to the right of i apart from the smallest one, say τ_0 . Then, M_i is a multiple of M_{τ_0} , E_i is a multiple of E_{τ_0} etc. Hence, an equation holds true in (15).

Next, assume that the largest step to the right of i is not of smallest type. Call this step $\tau \in \Pi$

Now, let

$$C^{(i)} = S^{(i)} + R^{(i)}$$

be the decomposition of its domain into the satellites and the remainder. Then consider all players of τ 's type: by Lemma 1.1, they are in $C^{(i)}$ and since "steps rule their followers" (OSTMANN [6]) they must appear in $S^{(i)}$. For simplicity we treat the case that τ is alone in his type only.

Thus, with a suitable $d_1 \in \mathbb{N}$ and suitable integer vector D , $\tau < D$, we have

$$M(S^{(i)}) = dM_\tau + DM.$$

From this, we conclude first of all:

$$(17) \quad \frac{E(S^{(i)})}{F(S^{(i)})} = \frac{dE_\tau + DE}{dF_\tau + DF} = \frac{dA_\tau + dE(S^{(d)}) + DE}{dB_\tau + dF(S^{(d)}) + DF} > \frac{E(R^{(i)})}{F(R^{(i)})}$$

by Lemma 1.8. Finally, we come up with

$$(18) \quad \frac{E_1}{F_1} = \frac{E(S^{(i)})}{F(S^{(i)})} + \frac{E(R^{(i)})}{F(R^{(i)})} = \frac{E(G^{(i)})}{F(G^{(i)})}, \quad \text{q.e.d.}$$

Essentially, Lemma 1.10 can be visualized by inspecting Fig.2. In fact, for any interval of players of the form $[i,j]$ the quotient

$$\frac{E([i,j])}{F([i,j])}$$

ranges between the one supplied by the largest step within this interval and the one of the *second largest step* to the right of this interval (more precisely: the smallest member of the second largest type) (cf. Fig.3).

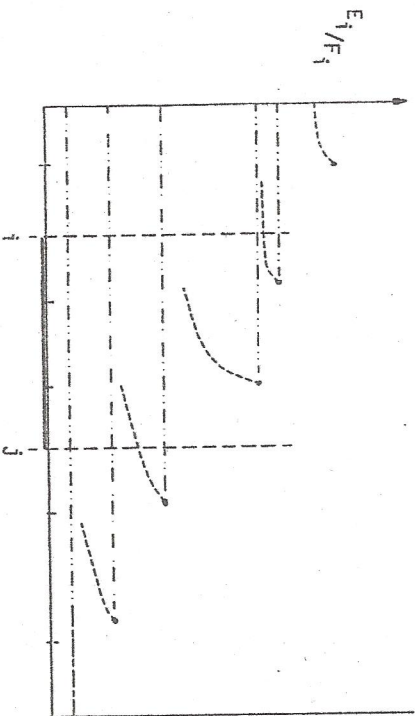


Fig.3

In particular, the quotient $E([i,j]) / F([i,j])$ is seen to decrease (strictly) at every step.

In particular, if we have a game without dummies, and at least two steps of different type, we may construct strictly increasing affine families of representations. In such representations it turns out that, with increasing parameter, an increasing amount of weight is accumulated on the lex-max min-win coalition.

Theorem 1.11:

Let v be homogeneous and $(M(\cdot), \lambda(\cdot))$ an a.f.r.. Define

$$(19) \quad Q : \mathbb{R}_{++} \rightarrow \mathbb{R},$$

$$Q(a) = \frac{\lambda(a)}{M(a)(\Omega)}.$$

1. If $(M(\cdot), \lambda(\cdot))$ is monotone, then Q is a monotone increasing function in a .
2. If $(M(\cdot), \lambda(\cdot))$ is strictly monotone, then so is Q .

Proof:

Clearly (omitting the argument a)

$$(20) \quad Q(a) = \frac{M(S(\lambda))}{M(\Omega)} = \frac{aE(S(\lambda)) + F(S(\lambda))}{aE(\Omega) + F(\Omega)}.$$

In order to show that this is (strictly) monotone, it suffices to show that

$$(21) \quad \frac{E(S(\lambda))}{F(S(\lambda))} > \frac{E(\Omega)}{F(\Omega)}$$

holds true.

Both sets are intervals and $S(\lambda)$ does not include all steps of the second largest type. Thus, in a strict a.f.r., (21) is a strict inequality, q.e.d.

Section 2

The nucleolus: preliminary results

From now on we shall always assume that any homogeneous game under consideration has no dummies.

Therefore, the smallest player is always a step (SEC.0, Remark 0.1.1). Following the tradition of SUDHÖLTER [14] and PELEG-ROSENTHAL [9] we speak of a homogeneous game "without steps" if the smallest player is the only step. Note that for games "without steps" the representation is unique up to a multiple and that constant-sum games are games "without steps".

Definition 2.1:

(1) $\mathcal{G}^* = \{x \in \mathbb{R}^N \mid x(S) = 1\}$
 is the set of *pre-imputations*. Also

(2) $\mathcal{G} = \{x \in \mathcal{G}^*(\Omega) \mid x \geq 0\}$
 is the set of *pseudo imputations*.

The *nucleolus* of a game was introduced by SCHMEIDLER [12], see also MASCHLER-PELEG-SHAPILEY [5]; usually, it is defined with respect to a set of payoff vectors. Tentatively, the *pre-nucleolus* $\mathcal{N}^*(v)$ is meant to be the one defined with respect to \mathcal{G}^* and the *pseudo nucleolus* $\mathcal{N}(v)$ is meant to be defined with respect to \mathcal{G} .

In [12] it is shown that the pseudo-nucleolus consists of a unique pseudo imputation $\nu(v)$ (also called "the pseudo-nucleolus of v ").

Note that in our context of homogeneous games, we assume neither superadditivity of v nor do we exclude singletons to be winning coalitions. However, even if single players form winning coalitions ("are winning"), we do not encounter additional problems, for \mathcal{N}^* and \mathcal{N} are equal, as is stated by the following lemma.

For any game v on Ω and $x \in \mathbb{R}^N$ let us use the notation $e(S,x) = v(S) - x(S)$ to denote the excess of x (at S). Also let

(3) $\mu = \mu(x,v) = \max \{e(S,x) \mid S \in \underline{P}\}$

and

(4) $\mathcal{M} = \mathcal{M}(x,v) = \{S \in \underline{P} \mid e(S,x) = \mu(x,v)\}.$

Now, we have

Lemma 2.2:

Let v be a homogeneous game. Then $\mathcal{N}^*(v) = \mathcal{N}(v)$.

Proof:

Since $\mathcal{G} \subset \mathcal{G}^*$, it suffices to, given any $x^* \in \mathcal{G}^*$ with negative coordinates, construct $x \in \mathcal{G}$ such that

(5) $\mu(x,v) < \mu(x^*,v)$

holds true.

To this end, fix $x^* \in \mathcal{G}^*$ and define

(6) $P := \{i \in \Omega \mid x_i^* > 0\}, N := \{i \in \Omega \mid x_i^* < 0\},$

we assume $N \neq \emptyset$. Pick $x \in \mathcal{N}$ such that the following conditions are satisfied

(7) $0 \leq x_i < x_i^* \quad (i \in P)$

(8) $0 = x_i \quad (i \in \Omega - P).$

Now, consider $S \in \mathcal{M}(x,v)$. If $P \subset S$, then $0 \geq e(S,x) = \mu(x,v)$ is verified at once. As $e(\{i\}, x^*) > 0$ ($i \in N$), we are done, since (5) holds obviously true.

If, on the other hand $P \subset S$ prevails, then let $S' := S \cup N$. Clearly

(9) $v(S') \geq v(S)$

since any homogeneous game is monotone. Moreover

$$\begin{aligned} x^*(S^-) &= x^*(S \cup P) + x^*(N) \\ &= x^*(S \cup P) + x(P) - x^*(P), \end{aligned}$$

(since $x^*(N) + x^*(P) = 1 = x(P)$), and thus

$$\begin{aligned} (10) \quad x^*(S^-) &= x(x^-) - x^*(x^- - S) \\ &= (x - x^*)(P - S) + x(P \cap S) \\ &= (x - x^*)(P - S) + x(S) \\ &< x(S) \end{aligned}$$

(observe (7) and (8)). Combining (9) and (10) we obtain

$$\begin{aligned} e(S, x^*) &= v(S) - x^*(S^-) \\ &\geq v(S) - x^*(S^-) \\ &> v(S) - x(S) = e(S, x), \end{aligned}$$

which proves (5),

q.e.d.

According to KOHLBERG [4], a collection

$$\mathcal{A} = \{\underline{B}_0, \dots, \underline{B}_p\}, \quad \underline{B}_q \subseteq \underline{P} \quad (q = 0, \dots, p)$$

of systems of coalitions is called a *coalition array* if \underline{B}_0 contains only singletons and

$$\underline{B}_1 + \dots + \underline{B}_p = \underline{P}$$

holds true.

Given a homogeneous game v on Ω and a pseudo imputation $x \in \mathcal{X}$, a coalition array

$\mathcal{A}(x, v)$, i.e.,

$$\underline{B}_0 = \underline{B}_0(x), \underline{B}_1 = \underline{B}_1(x, v), \dots, \underline{B}_p = \underline{B}_p(x, v)$$

is specified as follows:

1. $\underline{B}_0(x) = \{\{i\} \mid x_i = 0\}$
- (11) 2. $e(S, x) = \text{const} \quad (S \in \underline{B}_j(x, v)) \quad (j = 1, \dots, p)$
3. $e(S, x) < e(T, x) \quad (S \in \underline{B}_j(x, v), T \in \underline{B}_{j-1}(x, v)) \quad (j = 2, \dots, p)$

A coalition array has *property I* if, for all $q \in \{1, \dots, p\}$ and $y \in \mathbb{R}^{\Omega}$ satisfying

$$(12) \quad y(S) \geq 0 \quad (S \in \bigcup_{j=0}^q \underline{B}_j)$$

$$(13) \quad y(\Omega) = 0$$

it follows that

$$y(S) = 0 \quad (S \in \bigcup_{j=1}^q \underline{B}_j)$$

holds true.

A coalition array has *property II* if, for all $q \in \{1, \dots, p\}$ there is a system of coefficients $c_S > 0 \quad (S \in \bigcup_{j=1}^q \underline{B}_j)$ and $c_S \geq 0 \quad (S \in \underline{B}_0)$ such that

$$(14) \quad \sum_{S \in \bigcup_{j=0}^q \underline{B}_j} c_S 1_S = 1_{\Omega}.$$

This means in particular that $\bigcup_{j=0}^q \underline{B}_j$ is weakly balanced.

The above exposition follows KOHLBERG [4]. For our purpose we quote some of his results as follows.

Theorem 2.3: ([4])

If v is a homogeneous game, then

$$\begin{aligned} \mathcal{N}(v) &= \{x \in \mathcal{X} \mid \mathcal{A}(x, v) \text{ has property I}\} \\ &= \{x \in \mathcal{X} \mid \mathcal{A}(x, v) \text{ has property II}\}. \end{aligned}$$

Again, it should be noted that the assumption $v(\{i\}) = 0 \quad (i \in \Omega)$ (i.e., there are no winning players) can be dropped without destroying the proofs.

Theorem 2.4:

Let v be a homogeneous game and let

$$k = \max \{i \in \Omega \mid \{i\} \in \underline{W}^m\}$$

be the smallest winning player ($\max \emptyset = a-1$).

1. If $\kappa = b$ then

$$v(v) = (\bar{v}_1, \dots, \bar{v}_n) \quad (n := b - a + 1)$$

2. If $\kappa < b$, let \bar{v} denote the homogeneous game on $\bar{\Omega} = [k+1, b]$ which is obtained by dropping the winning players. Also, let $\bar{v} = v(\bar{v})$ and

$$(15) \quad \begin{aligned} \bar{a} &:= 1 - \mu(\bar{v}, \bar{v}) \\ \bar{x} &:= (\bar{a}_1, \dots, \bar{a}_k, \bar{v}_{\kappa+1}, \dots, \bar{v}_b) / (k+1-a) \bar{a} + 1. \end{aligned}$$

$\underbrace{\hspace{10em}}_{k+1-a \text{ times}}$

Then $v(v) = \bar{x}$.

In other words, the pseudo nucleolus of v is obtained by computing the one on \bar{v} , then assigning \bar{a} to the winning players and finally rescaling.

Proof.

The case $\kappa = b$ and $\kappa = a-1$ is trivial; so we have to concentrate on the second case for $a \leq \kappa < b$.

Consider the coalition array $\mathcal{S}(\bar{x}, v)$, we would like to show that it enjoys property I.

To this end, fix $q \in [1, p]$ and let $y \in \mathbb{R}^\Omega$ satisfy (12) and (13).

First of all note that $y_j \geq 0$ for $j \in [a, \kappa]$. For, in view of $e(\bar{x}, \{j\}) = 1 - \bar{x}_j = \frac{\mu(\bar{v}, \bar{v})}{(k+1-a)\bar{a}+1}$, it turns out that $\{j\} \in \underline{B}(\bar{x}, v)$. Thus, y is nonnegative on $[a, \kappa]$.

Next, if $y([a, \kappa]) = 0$, then clearly $y(S) = 0$ for all $S \in \sum_{j=1}^{\kappa} \underline{B}_j(\bar{x}, v)$; this follows by the fact that \bar{v} is the nucleolus of \bar{v} .

Finally, if $0 < y([a, \kappa]) =: \bar{\beta}$, then define

$$(16) \quad \tilde{y} := (y_{\kappa+1} + \frac{\bar{\beta}}{\sigma_{\kappa}}, \dots, y_b + \frac{\bar{\beta}}{\sigma_{\kappa}}).$$

It is not hard to see that \tilde{y} indeed satisfies (12) and (13) with respect to the game \bar{v} and, say, $q = 1$. This is a contradiction to Theorem 2.3, since \bar{v} is the nucleolus of \bar{v} - hence this case cannot occur and we have finished our proof. q.e.d.

The last theorem shows that we may disregard the case that winning players are present.

Hence, from now we shall assume that all homogeneous games under consideration have no winning players (i.e. $M_i < \lambda$ ($i \in \Omega$) for any representation (M, λ) of some v).

Consequently, the prefixes "pre" and "pseudo" may be omitted, thus \mathcal{S} is the set of imputations and $v = v(v)$ the nucleolus.

Remark 2.5: (KOHLEBERG [4], Theorem 1.4)

Let κ denote for the moment the last player who gets a payoff with the nucleolus, i.e.,

$$(17) \quad \kappa = \max \{i \in \Omega \mid v_i > 0\}.$$

Then $\{S \cap [a, \kappa] \mid S \in \mathcal{K}(x, v)\}$ is strongly balanced.

If v is a homogeneous game "with steps of different type" (other than the smallest nondummy that is), then it can be inferred easily, that \underline{v}^m cannot be strongly balanced (see also Remark 5.4. in (PELEG-ROSENTHALER [9]).

Now, in view of the exhibition presented in SECTION 1, we can easily show that \underline{v}^m cannot be weakly balanced. In fact, we show a bit more:

Corollary 2.6:

Let v be a homogeneous game with steps of different type (no dummies, no winning players). Then 1_Ω is no linear combination of $(1_S)_{S \in \underline{v}^m}$.

Proof:

Let $(M(\cdot), \lambda(\cdot))$ be a strictly monotone a.f.r. and suppose that, for some system of coefficients $(c_S)_{S \in \underline{W}^m}$, we have

$$\sum_{S \in \underline{W}^m} c_S 1_S = 1_\Omega.$$

Then

$$(18) \quad M(a)(\Omega) = \sum_{S \in \underline{W}^m} c_S M(a)(S) = \lambda(a) \sum_{S \in \underline{W}^m} c_S.$$

Now, $\sum c_S$ is a constant, thus (18) contradicts Theorem 1.11. which states that the quotient $\frac{\lambda(a)}{M(a)(\Omega)}$ is a strictly increasing function, q.e.d.

Section 3 The nucleolus for games with steps

As we have mentioned, we will from now on only deal with homogeneous games without dummies and winning players.

The behavior of the nucleolus for games "without steps" has been described in PELEG-ROSENMÜLLER [9]. Here, we want to tackle the same problem when steps are present.

There is an inductive procedure involved in our method which (unlike the method of satellite games as explained in SECTION 0) uses a truncation procedure cutting off smaller players. To explain this version of "truncated games" we have to shortly recall the theory of the incidence vector of a homogeneous game (without steps), as developed by SUDHÖLTER [14].

To this end we fix $\Omega = [1, n]$ (i) throughout this section and focus our attention (initially only) on a homogeneous game v without steps. Let (M, λ) be its unique minimal representation so that $v = v_\lambda^M$.

Next, consider $S \in \underline{W}^m$ and let $\ell = \ell(S)$ again denote the last player in S . Suppose $j \in S$ is such that

$$(1) \quad [j, \ell] \subset S, S - j + [\ell+1, n] \in \underline{W}.$$

Then j is expendable; we may replace him in S by an interval of smaller players, thus generating a coalition

$$(2) \quad \rho_j(S) := S - j + [1(S)+1, \ell]$$

where ℓ is uniquely defined by $M(\ell(S)+1, \ell) = M_j$. This procedure is based on the BASIC LEMMA (ROSENMÜLLER [10]), see SUDHÖLTER [14].

On the other hand, let $T \in \underline{W}^m$ and suppose that $r \notin T$ satisfies

$$(3) \quad [r+1, \ell(T)] \subset T.$$



Then r is the last dropout (see ROSENTHALER [10]) and there is a unique $v \in [r+1, k(T)]$ such that

$$(4) \quad q(T) := T + r - [v, k(T)]$$

is min-win. That is, v inserts the last dropout and cuts off an appropriate tail of T as to generate a min-win coalition. And thus, ρ_j renders j to be the last dropout if he is expendable in S .

Clearly, if r is the last dropout in T , then (he is expendable in $q(T)$ and)

$$(5) \quad \rho_r(q(T)) = T.$$

Similarly, if j is expendable in S then (he is the last dropout in $\rho_j(S)$ and)

$$(6) \quad q(\rho_j(S)) = S$$

holds true.

According to SUDHÖLTER [14] we have

Lemma 3.1: (cf. [14], Theorem 2.3, Definition 2.4)

Let v be a homogeneous game (without dummies and winning players). Assume that v has no steps and is not the unanimous game of the grand coalition.

Then there is a unique sequence S_1, \dots, S_n of min-win coalitions defined by the following procedure.

1. $S_1 = S(N)$
2. For every $k \in [1, n-1]$, the system $\underline{S}_k := \{S_i \mid i \in [1, k], k \text{ is expendable in } S_i\}$ is nonempty.
3. Among all $S_i \in \underline{S}_k$ with minimal length $k(S_i)$, let S_i be the one with smallest (first) index.
4. $S_{k+1} = \rho_k(S_i)$.

Definition 3.2:

Let S_1, \dots, S_n be given by Lemma 3.1. Then

$$\ell = \ell(v) = (\ell_1, \dots, \ell_n) := (k(S_1), \dots, k(S_n))$$

is the incidence vector of v .

The incidence vector characterizes v uniquely (Theorem 2.10 of [14]). (The term "incidence vector" can be defined abstractly.)

Given the incidence vector, the game v can be obtained by "reversing" the procedure of Lemma 3.1. In other words, the sequence S_1, \dots, S_n can be constructed in a unique way and, since we are dealing with a game without steps, the unique minimal representation is obtained at once.

Let us shortly describe this "reversal procedure".

Given $\ell = \ell(v) = (\ell_1, \dots, \ell_n)$, the staircase corresponding to $\ell(v)$ (and hence to v) is the vector

$$\pi = \pi(v) \in \mathbb{N}^n$$

given by

$$(7) \quad \pi_k = \min \{\ell_j \mid j \leq k \leq \ell_j\} \quad (k = 1, \dots, n)$$

(with the convention that $\min \emptyset = 0$). If π is regarded as a function of k , then it is monotone and can be identified as a "quadratic step function" since the heights of jumps and the length of plateaus are equal (see [14], SECTION 3). E.G., if ℓ equals

$$(8) \quad \ell = (3, 7, 6, 5, 7, 7, 8)$$

then

$$\pi = (3, 3, 3, 5, 5, 6, 7, 8).$$

Thus π appears as a staircase with square steps that vary in height and width simultaneously (and with appropriate view, ℓ decreases on the plateaus of π but dominates π , cf. Fig. 4).

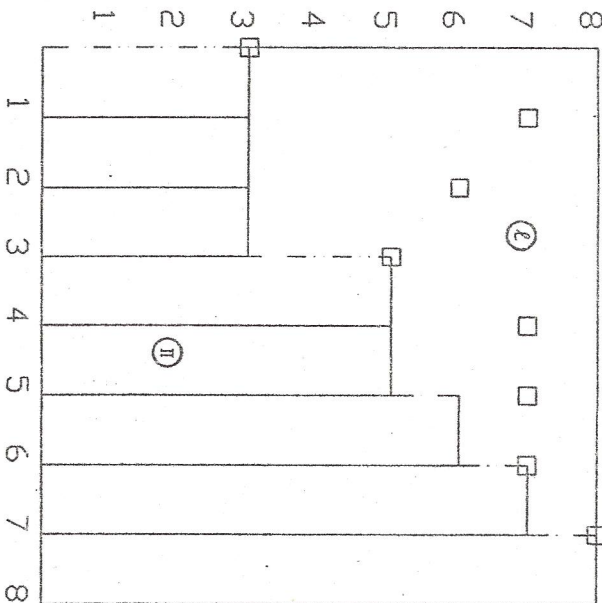


Fig. 4

(On the other hand, π_k denotes of course the minimal length of a coalition in \underline{S}_k - if we view Lemma 3.1).

Now, define the selector to be the vector $\omega = \omega^{(v)} = \omega^{(v)}$ which is given by

$$(9) \quad \omega_k = \min \{j \mid l_j = \pi_k\}$$

(again $\min \emptyset = 0$). Then ω selects the appropriate index i_0 in the formulation of Lemma 3.1. More precisely, given l , the sequence S_1, \dots, S_n as specified by Lemma 3.1 is given by

$$(10) \quad \begin{aligned} 1. \quad S_1 &= [1, l] \\ 2. \quad S_{k+1} &= S_{\omega_k} - k + [l_{\omega_k} + 1, l_{k+1}] = \rho_k(S_{\omega_k}). \end{aligned}$$

E.g., in the example suggested by (8), we obtain

$$\omega = (1, 1, 1, 4, 4, 3, 2, 8)$$

telling us that, e.g., in the fifth step of the construction suggested by Lemma 3.1 we have to render player 4 to become the last dropout in S_4, \dots

Remark 3.3: (SUDHÖLTER [14])

Let (M, λ) be the minimal representation of a homogeneous game with steps; assume $\Omega = [1, n]$ and write $M = (M_1, \dots, M_n)$. Let $\hat{M} := (M_1, \dots, M_n, 1)$. Then $v_{\lambda}^{\hat{M}}$ is an $(n+1)$ -person) homogeneous game without steps and (\hat{M}, λ) is its minimal representation.

Intuitively, if we add a player of weight 1, then his weight can just be used to close the "jumps" that appear at a step (cf. SEC.0).

Definition 3.4:

Let $\Omega = [1, n+1]$ and let l be an incidence vector (of a hom game for $n+1$ players without steps). Let $k \in [2, n]$. The truncation of l at k is the vector $l^{\circ} = l^{(k)} \in \mathbb{R}^{k+1}$ given by

$$(11) \quad l_i^{\circ} = \begin{cases} l_i & l_i \leq k \\ k & k < l_i \leq n, \pi_{i-1} < k \\ k+1 & \text{otherwise} \end{cases}$$

($\pi \in \mathbb{N}^{n+1}$ is the staircase corresponding to l)

Remark 3.5:

$l^{(k)}$ is an incidence vector.

Proof:

This follows immediately by the observation that $l^{(k)}$ enjoys a corresponding staircase, namely

$$(12) \quad \pi^{(k)} = (k \wedge \pi \mid [1, k], k+1)$$

where "min" ($= \wedge$) has to be taken coordinate-wise.

Definition 3.6:

Let v be a homogeneous game (with steps) on $\Omega = [1, n]$. Let $\kappa \in [2, n]$. The truncation of v at κ , $\hat{v}^{(\kappa)}$ is defined as follows.

1. Let (M, λ) be the minimal representation. Let $(\tilde{M}, \tilde{\lambda})$ be obtained by Remark 3.3 and let $\hat{\ell}$ be the incidence vector of $\hat{v} = v \tilde{\lambda}$.
2. Let $\hat{\ell} \in \mathbb{R}^{\kappa+1}$ be the truncation of $\hat{\ell}$ at κ as defined by 3.4. $\hat{\ell}$ generates a homogeneous game \hat{v} on $[1, \kappa+1]$ with minimal representation $(\hat{M}, \hat{\lambda})$, $\hat{M} \in \mathbb{R}^{\kappa+1}$.
3. $\hat{v}^{(\kappa)}$ is the game which is (minimally) represented by

$$\tilde{M} := \begin{pmatrix} \hat{M}_1 & \dots & \hat{M}_\kappa \\ \vdots & & \vdots \end{pmatrix} \Big|_{[1, \kappa]} \\ \tilde{\lambda} := \hat{\lambda}$$

Note that homogeneous games without steps indeed attach weight 1 to the smallest two players (w.r.t. the minimal representation). Of course the one-to-one correspondence between homogeneous games and incidence vectors is heavily used (cf. Theorem 2.10 of SUDHÖLTER [14]).

Our first aim is to obtain some insight into the structure of the truncated versions. The following lemma is an attempt to describe the min-win coalitions of some $\hat{v}^{(\kappa)}$.

To this end, let us slightly augment our notation:

If $S \in \underline{W}^m(v)$ and $r \notin S$, $r < l(S)$, (r is *arg dropout*), then:

$$(13) \quad \varphi(S) := (S \cap [1, r-1]) + [r, l] \in \underline{W}^m(v).$$

Thus, $\varphi(S)$ is the lexicographical first min-win coalition among all coalitions T with $T \cap [1, r-1] = S \cap [1, r-1]$.

Lemma 3.7:

Let v be a homogeneous game (on $\Omega = [1, n]$) and let $\kappa \in [2, n]$. Then $\hat{v}^{(\kappa)}$ has the following properties.

1. If $i \in [1, \kappa]$ is a step of $\hat{v}^{(\kappa)}$, then it is a step of v or $i \sim \kappa$. On the other hand, if $i \in [1, \kappa]$ is a step of v , then he is a step of $\hat{v}^{(\kappa)}$.
2. If $S \subseteq [1, \kappa]$ and $S \in \underline{W}^m(v)$, then $S \in \underline{W}^m(\hat{v}^{(\kappa)})$.
3. If $S \subseteq [1, \kappa]$ and $S \in \underline{W}^m(\hat{v}^{(\kappa)}) - \underline{W}^m(v)$, then $\kappa \in S$ and $S \notin \underline{W}(v)$.
4. If $S \subseteq [1, \kappa]$ and $S \in \underline{W}(\hat{v}^{(\kappa)})$, then $S + [\kappa+1, n] \in \underline{W}(v)$.
5. If $S \supseteq [1, \kappa]$ and $S \in \underline{W}^m(v)$, then $[1, \kappa] \in \underline{W}^m(\hat{v}^{(\kappa)})$ and $\hat{v}^{(\kappa)}$ is the unanimous game of $[1, \kappa]$.
6. If $S \in \underline{W}^m(v)$, $S \not\subseteq [1, \kappa]$, $S \not\supseteq [1, \kappa]$, and, with $r = l([1, \kappa] \cap S)$, $\varphi(S) \subseteq [1, \kappa-1]$, then $S \cap [1, \kappa] \in \underline{W}^m(\hat{v}^{(\kappa)})$.

Proof:

1. Given v , let $\hat{\ell} \in \mathbb{R}^{\kappa+1}$ be defined via 3.6.2. In view of Chapter 2 of [14], it is known that player i is a step w.r.t. $\hat{v}^{(\kappa)}$ iff $\hat{\ell}_{i,1} = \kappa+1$. In view of Definition 3.4, this leaves two alternatives for $\hat{\ell}_{i,1}$: either $\hat{\ell}_{i,1} = n+1$ - in which case i is a step of v . Or else $\hat{\ell}_{i,1} \geq \kappa$. But then (see (12)) $\pi^{(v)} = \kappa$ and $i \sim \kappa$.

The reverse statement is seen analogously.

2. To prove the second statement, assume that, on the contrary, for some $\bar{S} \subseteq [1, \kappa]$ it turns out that $\bar{S} \in \underline{W}^m(v)$ and $\bar{S} \notin \underline{W}^m(\hat{v}^{(\kappa)})$.

Clearly, since in this case $\ell_1 = \ell$, \bar{S} is not the lex-max min-win coalition of v . Hence, there exists the last dropout of \bar{S} , say $r \notin \bar{S}$, $r < l(S)$.

Now, among all \bar{S} with this property collect those with minimal length $l(\bar{S})$. And, among all those with minimal length, choose the one with maximal last dropout r . Call these now \bar{S} and r again.

Define

$$(14) \quad T := \varphi(\bar{S}).$$

Then, because $l(T) < l(\bar{S})$ holds true, it follows from our choice of \bar{S} that

$$(15) \quad T \in \underline{W}^m(v) \cap \underline{W}^m(\varphi(S)).$$

Next, we know that the procedure indicated by Lemma 3.1 (and Theorem 2.3 of [14]) yields two min-win coalitions of v , say S_0 and S_{r+1} such that player r is expendable in

S_0 and

$$(16) \quad S_{r+1} = \rho_r(S_0), \varphi(S_{r+1}) = S_0.$$

More precisely,

$$(17) \quad l(S_0) = \min \{l(S) \mid S \ni r, S \in \underline{W}^m(v)\} = \tau_r$$

and

$$(18) \quad l(S_{r+1}) = \min \{l(S) \mid S \ni r, S \in r+1, S \in \underline{W}^m(v)\} = \ell_{r+1}$$

$$\text{while } S_0 \cap [1, r-1] = S_{r+1} \cap [1, r-1]$$

is also true. Clearly, $S_0 \in \underline{W}^m(v^{(r)})$ in view of (17); in fact it follows from (17) that $l(S_0) \leq l(T)$. However, $l(S_0) < l(T)$ is impossible in view of our choice of \bar{S} and r . But $l(S_0) = l(T)$ implies via application of φ (cf. (14) and (16)) that $l(S_{r+1}) = l(\bar{S})$ holds true.

In this case, (18) shows that $\bar{S} \in \underline{W}^m(v^{(r)})$, and we have completed our proof of the second statement.

3. The third statement is verified by a sequence of analogous argument.

4. Follows from the definition of ℓ .

5. A trivial consequence.

6. Follows from Definition 3.6 and from 1.

q.e.d.

We are now in the position to tackle the nucleolus of a homogeneous game with steps. To this end, in what follows $\tau = \tau(v)$ denotes the first (largest) step of a homogeneous game v , i.e.,

$$(20) \quad \tau = \min \{i \in \Omega \mid i \in I(v)\}.$$

Similarly, $\tau = \tau(v)$ is the smallest player of the type of τ , i.e.,

$$(21) \quad \tau = \max \{i \in \Omega \mid i \in \tau\}.$$

Note that $[\tau, \tau]$ consists of steps that appear as block in any min-win coalition if they appear at all. Of course $\tau = \tau$ will frequently happen.

Theorem 3.8:

Let v be a homogeneous game on $\Omega = [1, \dots, n]$ and let $\tau = \tau(v)$ be the smallest player of the largest step's type. Let $\nu = \nu(v)$ be the nucleolus of v . Then $\nu_{r+1} = \dots = \nu_n = 0$.

Clearly, we have to treat the case of a game v with steps of different type only. Then, Corollary 2.6 shows at once that $\nu_n = 0$ is necessarily true. The problem is that 2.6 cannot be employed immediately in order to prove that all players behind the first steps get zero at the nucleolus - here we have to fall back on a truncation.

Proof:

Assume that, on the contrary, there is $k \in [\tau+1, n]$ such that

$$\nu_{r+1} \geq \dots \geq \nu_k > 0 = \nu_{k+1} = \dots = \nu_n.$$

The proof proceeds by treating various cases separately.

1st CASE: Assume that there is a coalition $\bar{S} \in \mathcal{K}(\nu, \bar{v})$ such that $\tau \in \bar{S}$, $[\tau+1, \kappa] \notin \bar{S}$.

This case is of course easy: define an imputation $x \in \mathcal{S}^v$ via

$$(22) \quad x_i = \begin{cases} \nu_i & i \in [1, \tau] \cup [\kappa+1, n] \\ \nu_i + \frac{1}{2} \sum_{j=\tau+1}^{\kappa} \nu_j & i = \tau \\ \frac{\nu_i}{2} & i \in [\tau+1, \kappa] \end{cases}$$

and observe that $\mu(x, \nu) \leq \mu(\nu, \nu)$ while $\bar{S} \notin \mathcal{K}(x, \nu)$.

Because "steps rule their followers", no min-win coalition has larger excess at x than at ν and \bar{S} has a smaller one - this contradicts the fact that ν is the nucleolus. This finishes our proof for the 1st CASE at once.

We may now assume that no \bar{S} of the kind treated already exists.

Then $\mathcal{K} = \mathcal{K}(\nu, \nu)$ allows for a partition, say

$$(23) \quad \mathcal{K} = \mathcal{K}_1 + \mathcal{K}_2$$

such that

$$\mathcal{K}_1 := \{s \in \mathcal{K} \mid [\tau, \kappa] \subseteq S\}$$

$$(24)$$

$$\mathcal{K}_2 := \{s \in \mathcal{K} \mid [\tau, \kappa] \subseteq S\}$$

holds true.

Both sets are nonempty since the nucleus of a game is contained in the kernel ([12]).

Now, we turn to the truncation $\bar{v}^{(\kappa)}$ of v at κ which, for short, we abbreviate by $\bar{v} := \bar{v}^{(\kappa)}$.

By Lemma 3.7, we know that τ is a step w.r.t. \bar{v} and, "in \bar{v} ", τ may or may not be of the same type as κ (see 3.7, No.1).

Accordingly, the next two cases treat these two possibilities. The easier one, in which τ and κ are of different type, is considered first.

2nd CASE: Let us treat the case that $\tau \neq \kappa$.

This means in fact that τ is the smallest player of the largest step's type also in \bar{v} (see 3.7, No.1).

First of all, let us define a mapping

$$(25) \quad * : \mathcal{K} \rightarrow \underline{\mathbb{W}}^m(\bar{v})$$

separately for $S \in \mathcal{K}_1$ and $S \in \mathcal{K}_2$.

1. For $S \in \mathcal{K}_1$, define

$$(26) \quad S^* := S \cap [1, \tau].$$

Indeed, $S^* \in \underline{\mathbb{W}}^m(\bar{v})$ is true since "step τ rules his followers" - thus the smaller players to the right of τ (of κ , since $S \in \mathcal{K}_1$) cannot appear in a min-win coalition without τ . Hence a min-win coalition has to be contained in S^* . It cannot be properly contained since $S \in \mathcal{K}$.

2. For $S \in \mathcal{K}_2$, define

$$(27) \quad S^* \text{ is the lexicographically first coalition in } \underline{\mathbb{W}}^m(\bar{v}) \text{ which satisfies}$$

$$S^* \cap [1, \tau] = S \cap [1, \tau].$$

Because of the decomposition (23) and (24), S^* cannot cut into $[\tau, \kappa]$, thus $|(S^*)| \geq \kappa$ and $[\tau, \kappa] \subseteq S^*$.

From our definition of S^* we conclude:

$$(28) \quad S^* \cap [1, \kappa] \in \underline{\mathbb{W}}^m(\bar{v}).$$

Indeed, for $S \in \mathcal{K}_1$ this follows from 3.7, No.2, and for $S \in \mathcal{K}_2$, this follows from 3.7, No.6.

Furthermore, it is seen that

$$(29) \quad S \cap [1, \kappa] = S^* \cap [1, \kappa]$$

for all $S \in \mathcal{K}$.

The final conclusion is straightforward:

By KOHLBERG's result (Remark 2.5) we obtain a set of nonnegative real numbers

$(c_S)_{S \in \mathcal{K}}$ such that

$$(30) \quad \sum_{S \in \mathcal{K}} c_S 1_S \cap [1, \kappa] = 1_{[1, \kappa]}.$$

By (29),

$$(31) \quad \sum_{S \in \mathcal{K}} c_S 1_{S^* \cap [1, \kappa]} = 1_{[1, \kappa]}.$$

This, in view of (28) means that $1_{[1, \kappa]}$ ($[1, \kappa]$ is the grand coalition in \bar{v}) is a linear combination of $(1_S)_{S \in \underline{W}^m(\bar{v})}$. Since \bar{v} has steps (at least τ), this contradicts

Corollary 2.6.

3rd CASE: Now we treat the case that $\tau \sim \kappa$.

Again, we want to construct some contradiction between 2.5 and 2.6 -- however, as we are not in the position to claim that \bar{v} has steps, the procedure of the case has to be modified. We will eventually consider the truncation $v' = \bar{v}^{(\kappa'+1)}$ for some $\kappa' \geq \kappa$ and in this truncation (30) and (31) will have appropriate analogues.

To this end, let us proceed by several steps. The first step is to define the "critical player" κ' .

1st STEP (of the 3rd CASE):

As τ is the smallest player of his type "in v ", there is $\bar{T} \in \underline{W}^m(v)$ such that $\tau \in \bar{T}$, $\tau+1 \notin \bar{T}$.

Since $\tau \sim \kappa$, it follows necessarily that $\bar{T} \cap [1, \kappa] \neq \emptyset$ -- otherwise \bar{T} would be min-win in

\bar{v} (No.2 of 3.7) and separate τ and $\tau+1$.

Now, choose \bar{T} to be lexicographically maximal with the above properties, (i.e., $\tau \in \bar{T}$, $\tau+1 \notin \bar{T}$, $\bar{T} \cap [1, \kappa] \neq \emptyset$, $\bar{T} \in \underline{W}^m(v)$). Then we have, in addition

$$(32) \quad [\tau+2, 1(\bar{T})] \subset \bar{T}$$

and

$$(33) \quad 1(\bar{T}) > \kappa.$$

Again, among all coalitions with these properties, choose the one with minimal length.

Define

$$(34) \quad \kappa' := 1(\varphi_{1,1}(\bar{T})) = \min \{1(S) \mid S \in \underline{W}^m(v), \tau+1 \in S\}.$$

Now, $\kappa' \geq \kappa$ holds true. Indeed, otherwise $\tau \sim \kappa$ would be violated by No.6 of 3.7.

Consider now the truncation of v at $\kappa'+1$, say

$$v' = \bar{v}^{(\kappa'+1)}.$$

By No.6 of 3.7 and (34) it follows that $\bar{T} \cap [1, \kappa'+1] \in \underline{W}^m(v')$, and as $\tau+1 \notin \bar{T}$, τ and $\tau+1$ are of different type in v' .

Thus (see No.1 of 3.7) it turns out that τ is the smallest player of the largest steps type w.r.t. v' , i.e.,

$$\tau(v') = \tau(v) = \tau.$$

Therefore, we shall now concentrate our efforts on v' and try to imitate the procedure of the 2nd CASE.

2nd STEP (of the 3rd CASE):

Consider any coalition $\hat{S} \in \mathcal{K}$, such that $[1, \kappa+1, 1(\hat{S})] \subset \hat{S}$. Such coalitions exist: we may generate them from arbitrary elements of \mathcal{K} , by successively involving the last dropout. We claim:

$$(35) \quad 1(\hat{S}) > \kappa' \text{ for all } \hat{S} \in \mathcal{K}, [1, \kappa+1, 1(\hat{S})] \subset \hat{S}.$$

Indeed, if for some \hat{S} , (35) is violated, then consider

$$\hat{S} := \hat{S} - \{\tau+1\} + \{\tau+1, \dots, 1(\bar{T})\}.$$

This is a winning coalition of v which satisfies

$$v(\hat{S}) = v(\hat{S}) - \nu_{\kappa+1} < \nu(\hat{S}) < 1 - \mu(\nu, v),$$

contradicting the fact that ν is the nucleolus of v .

3rd STEP (of the 3rd CASE):

We can now repeat our argument, as presented in the 2nd CASE, but for v' .

Again define $*$: $\mathcal{K}(\nu, v) \rightarrow \underline{W}^m(v)$:

For $S \in \mathcal{K}$:

$$(36) \quad S^* := S \cap [1, \tau]$$

and $S^* \in \underline{W}^m(v)$ follows exactly as in the 2nd CASE, while $S^* \cap [1, \kappa+1] = S \cap [1, \kappa+1]$ is trivial.

For $S \in \mathcal{K}$, choose $S^* \in \underline{W}^m(v)$ to be lexicographically maximal with

$$(37) \quad S^* \cap [1, \tau] = S \cap [1, \tau].$$

Then $S^* \cap [\kappa+1] = S \cap [\kappa+1]$ follows from (35). Again, in view of No.2 and No.6 of 3.7,

$$(38) \quad \{S^* \cap [1, \kappa+1] \mid S \in \mathcal{K}\} \subseteq \underline{W}^m(v).$$

Next, by Remark 2.5, we find coefficients $(c_S)_{S \in \mathcal{K}}$ such that

$$(39) \quad \sum_{S \in \mathcal{K}} c_S 1_S \cap [1, \kappa] = 1_{[1, \kappa]}.$$

But for $S \in \mathcal{K}$ it is clear that $S^* \cap [1, \kappa+1] = S \cap [1, \kappa]$ (cf. (36)).

Fortunately, for $S \in \mathcal{K}$, (35) and (37) yield $S^* \cap [1, \kappa+1] \supseteq [1, \kappa+1, \kappa+1]$. Thus, from (39) it follows that

$$(40) \quad \sum_{S \in \mathcal{K}} c_S 1_{S^* \cap [1, \kappa+1]} = 1_{[1, \kappa+1]}.$$

But, in view of (38), (40) contradicts Corollary 2.6.

q.e.d.

Corollary 3.9:

With the notation of Theorem 3.8 the vector $(\nu(\hat{v}^{\tau(v)})), 0, \dots, 0)$ is the nucleolus of v . $b^{-\tau(v)}$ times

Proof:

By Theorem 3.8 $\nu(v)_i = 0$ for all $i > \tau(v)$. Thus $\nu(v)$ is the nucleolus of the game $v_{\lambda}^M((\tau(v)+1, b])$ (where $v = v_{\lambda}^M$). This game obviously coincides with the truncated game $\hat{v}^{\tau(v)}$ with $b^{-\tau(v)}$ additional dummies.

Example 3.10:

q.e.d.

Consider the pair

$$(M, \lambda) = (12, 10, 5, 3, 2, 2, 1, 1, 2, 2)$$

and the game $v = v_{\lambda}^M$. Put $a = 1, b = 8$, i.e., $\Omega = \{1, \dots, 8\}$.

Then it is easy to see that players 3 and 8 are the steps of v . Thus, according to Theorem 3.8, we have $\tau(v) = 3$. Writing coalitions as 0-1 vectors is instructive, thus from the following sequence of min-win coalitions (cf. Lemma 3.1)

$$\begin{array}{l} S_1 = (1100000000) \\ S_2 = (0111111000) \\ S_3 = (1011110000) \\ S_4 = (1001111111) \\ S_5 = (1010111100) \\ S_6 = (1011010000) \\ S_7 = (0111110110) \\ S_8 = (1010111010) \\ S_9 = (0111110101) \end{array}$$

The coalition S_3 at the origin of this arrow is used to construct the coalition S_6 at the top of the arrow via $\rho_5(S_3)$.

we obtain the incidence vector \hat{l} of $\hat{v} = v_{\lambda}^M$ (cf. Definition 3.6):

$$\hat{l} = (2, 6, 5, 9, 7, 6, 8, 8, 9).$$

Now the truncation of \hat{l} at 3 is $(2, 3, 3, 4)$, which generates $(1, 1, 1, 2)$, thus the truncated game $v^{(3)}$ can be represented by $(1, 1, 1, 2)$, showing that all players are of the same type. Corollary 3.9 at once enables us to write down the nucleolus of v :

$$\left\{ \frac{1}{3}, \frac{1}{3}, 0, 0, 0, 0, 0 \right\}.$$

LITERATURE

- [1] Isbell, J.R.:
A class of majority games.
Quarterly Journal Math. 7 (1956), pp.183-187
- [2] Isbell, J.R.:
A class of simple games.
Duke Math. Journal 25 (1958), pp.423-439
- [3] Isbell, J.R.:
On the enumeration of majority games.
Math Tables Aids Comput. 13 (1959) pp.21-28
- [4] Kohlberg, E.:
On the nucleolus of a characteristic function game.
SIAM Journal Appl. Math. 20 (1971), pp.62-66
- [5] Mascher, M., Peleg, B., and Shapley, L.S.:
Geometric properties of the kernel, nucleolus, and related solution concepts.
Math. of Operations Research 4 (1979), pp.303-338
- [6] Ostmann A.:
On the minimal representation of homogeneous games.
Int. Journal of Game Theory 16 (1987), pp.69-81
- [7] Peleg, B.:
On the kernel of constant-sum simple games with homogeneous weights.
Ill. Journal Math. 10 (1966), pp.39-48
- [8] Peleg, B.:
On weights of constant-sum majority games.
SIAM Journal Appl. Math. 16 (1968), pp.527-532
- [9] Peleg, B. and Rosenmüller, J.:
The least-core, nucleolus, and kernel of homogeneous weighted majority games.
Working paper No. 193, Institute of Mathematical Economics, University of Bielefeld (1990), 26 pp.
- [10] Rosenmüller, J.:
Homogeneous games : recursive structure and computation.
Math. of Operations Research 12 (1987), pp.309-330
- [11] Rosenmüller, J.:
Homogeneous games with countably many players.
Math. Social Sciences 17 (1989), pp.131-159
- [12] Schneider, D.:
The nucleolus of a characteristic function game.
SIAM Journal of Applied Mathematics 17 (1969), pp.1163-1170
- [13] Shapley L.S.:
On balanced sets and cores.
Naval Res. Logist. Quarterly 14 (1967), pp.453-460
- [14] Sudhaker, P.:
Homogeneous games as anti step functions.
Intern. Journal of Game Theory 18 (1989), pp.433-469
- [15] von Neumann, J. and Morgenstern, O.:
Theory of games and economic behavior.
Princeton Univ. Press, NJ, (1944)